



DISCUSSION CLASS WEEK 7

CS 141 F20

Dynamic Programming

Graphs Basics

INTRO TO PALINDROME

- A word, phrase, or sequence that reads the same backward as forward
 - madam, 1991, 11, 5, a, cc
 - All strings of length 1 (one) is a palindrome

FUNCTION

```
bool isPalindrome(string str)
1  n = str.length
2  i = 1, j = n
3  while i < j
4      if str[i] != str[j]
5          return false
6      i++, j--
    return true
```

- Time complexity??

GENERALIZED FUNCTION

```
bool isPalindrome(string str, int start, int end)
1  i = start, j = end
2  while i < j
3      if str[i] != str[j]
4          return false
5      i++, j--
    return true
```

- Time complexity??

BUILDING PALINDROME

- **madam** is palindrome, then **ada** would be a palindrome too
- Opposite is true too.
- Adding same chars to both ends of a palindrome would still be a palindrome
- **P** is a palindrome, then **aPa** would be a palindrome too. (P is palindrome of length ≥ 0)

PROBLEM DEFINITION

- Given a string, find the length of the longest palindrome
 - madam = 5
 - babad = 3 (bab, aba)
 - dbabad = 3 (bab, aba)
 - cbbd = 2 (bb)
 - a = 1
 - ac = 1 (a, c)

NAÏVE SOLUTION

- Iterate over starting and ending position and check if it's a palindrome

```
int longestPalindrome(String str)
1  n = str.length, ans = 1
2  for i = 1 to n
3      for j = i+1 to n
4          if (isPalindrome(str, i, j))
5              ans = max(ans, j-i+1)
    return ans
```

- Time complexity??
- Can we make things better?

A BETTER NAÏVE SOLUTION

- If we already found a palindrome of length x , we only would need to check for palindrome for length $> x$

```
int longestPalindrome(String str)
1  n = str.length, ans = 1
2  for i = 1 to n
3      for j = i+ans to n
4          if (isPalindrome(str, i, j))
5              ans = max(ans, j-i+1)
    return ans
```


BETTER IDEA??

- If we **know** all the palindrome of size x , we can easily check for palindromes of $x+2$ with $O(1)$
 - Know = Memorization
- Finding answers of bigger palindrome easier with smaller palindrome. Bottom-up approach

FORMING DP

Base Cases

- All strings of length 1 (one) is a palindrome
- Let's find all palindromes of length 2

Recurrence

- Let's find palindromes of length x using palindromes of length $x-2$

MEMORIZATION

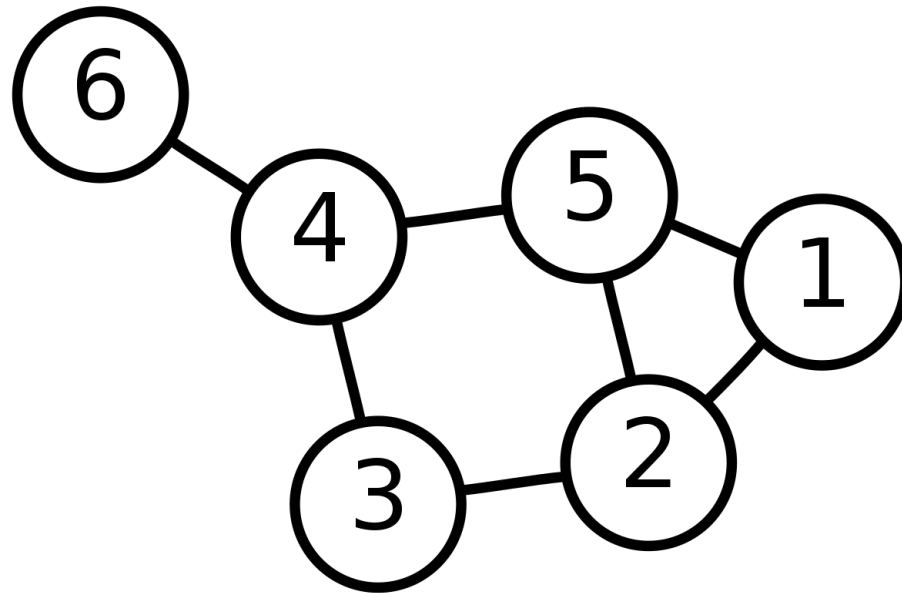
- How do we store the palindromes?
- 2d array signifying starting and ending points
- `if (isPalindrome(str, i, j) == true) => mem[i][j] = 1`
- `mem[x][y] = mem[x+1][y-1] & str[x] == str[y]`

DP ALGORITHM

```
int longestPalindrome(string str)
1  n = str.length, ans = 1
2  mem[n][n] = {0} // 2d array initialized
3  for i = 1 to n
4      mem[i][i] = 1
5  for i = 2 to n
6      if str[i] == str[i-1]
7          mem[i-1][i] = 1
8  for len = 3 to n
9      for i = 1 to n
10         j = i + len - 1
11         if (str[i] == str[j] && mem[i+1][j-1])
12             mem[i][j] = 1
13             ans = len
    return ans
```

GRAPHS

- A graph $G=(V,E)$ consists of vertices ($V, |V|=n$) and edges ($E, |E|=m$) that connect vertices together



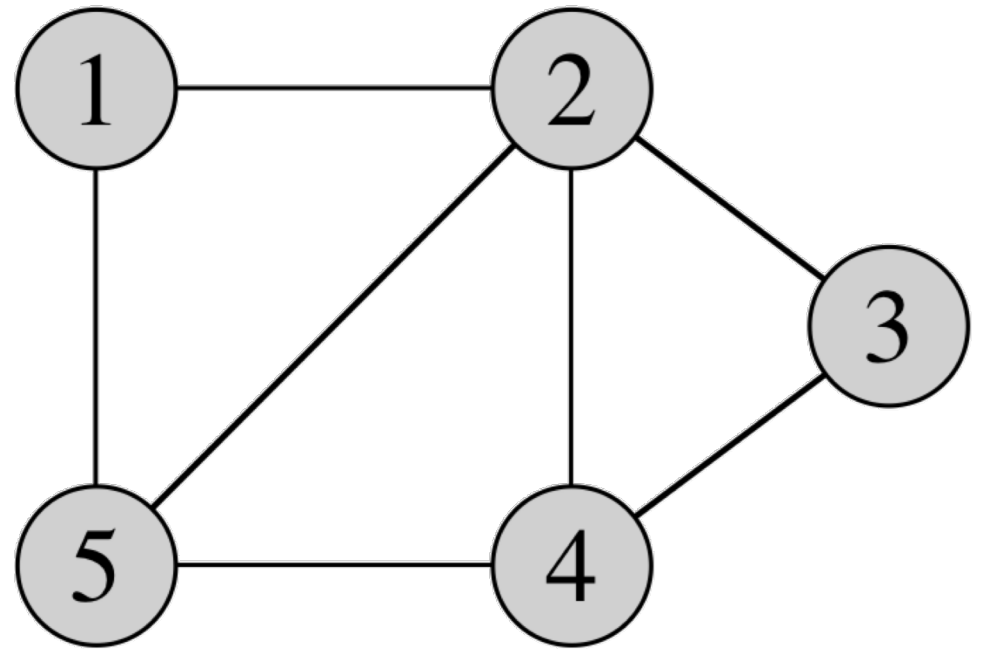
$n=6, m=7$

TYPES OF GRAPHS

- Directed and Undirected graphs
- Weighted and Unweighted graphs
- Connected graphs
- Bipartite graphs
- Acyclic graphs

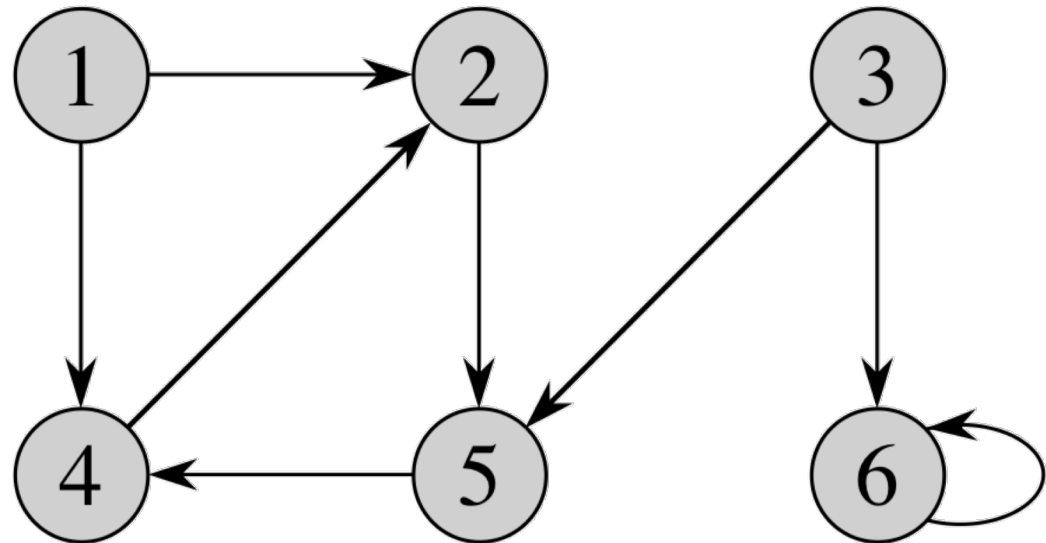
UNDIRECTED GRAPHS

- No direction in edges
- An edge can be traversed in both ways
- E.g., Facebook friends, most roads, most flights



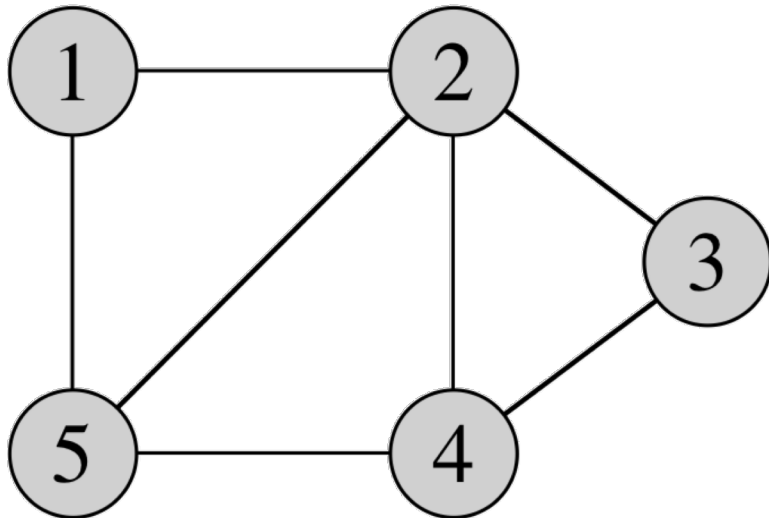
DIRECTED GRAPH

- Direction on edges
- An edge can be traversed in one direction
- E.g., Twitter follows, code flow analysis

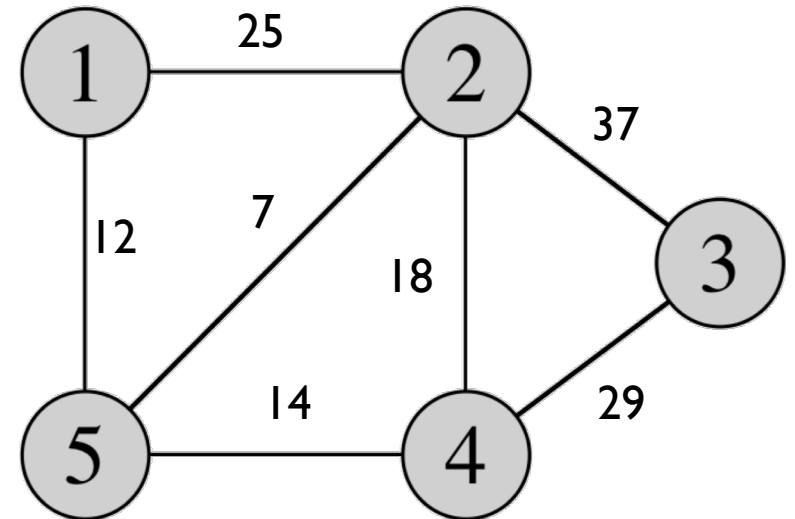


WEIGHTED GRAPH

- Vertices and/or edges can be assigned weights
- Weights can be cost, capacity, etc.
- E.g., road network, computer network



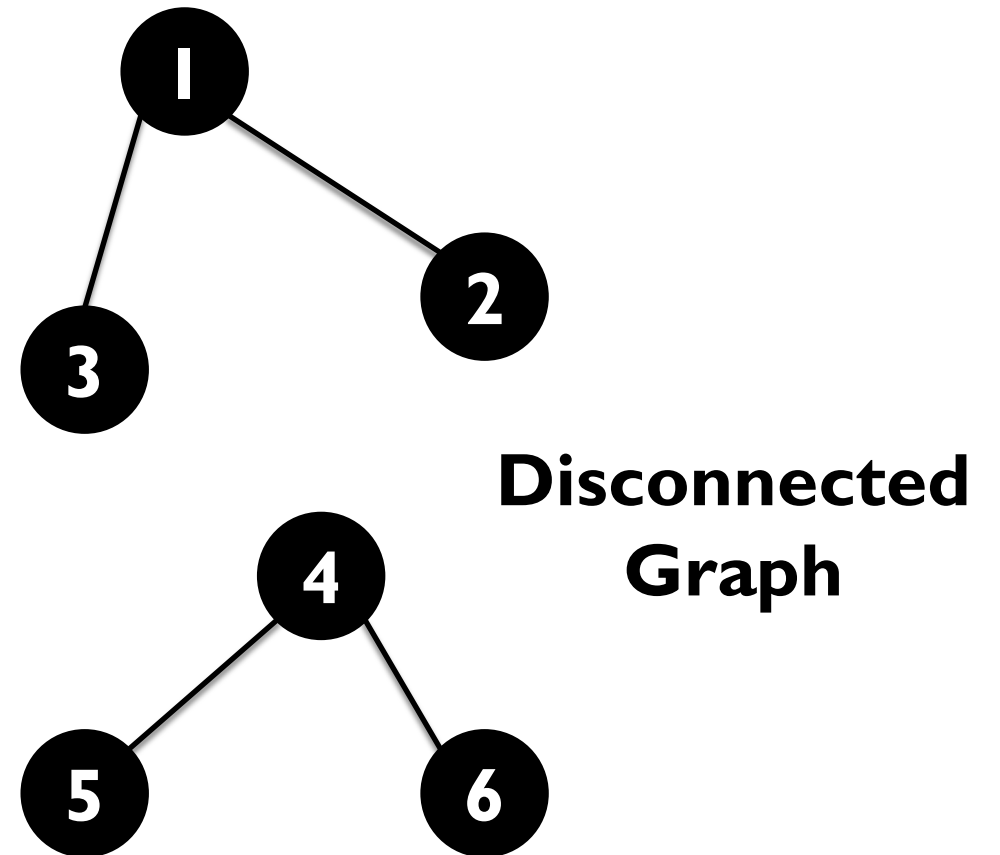
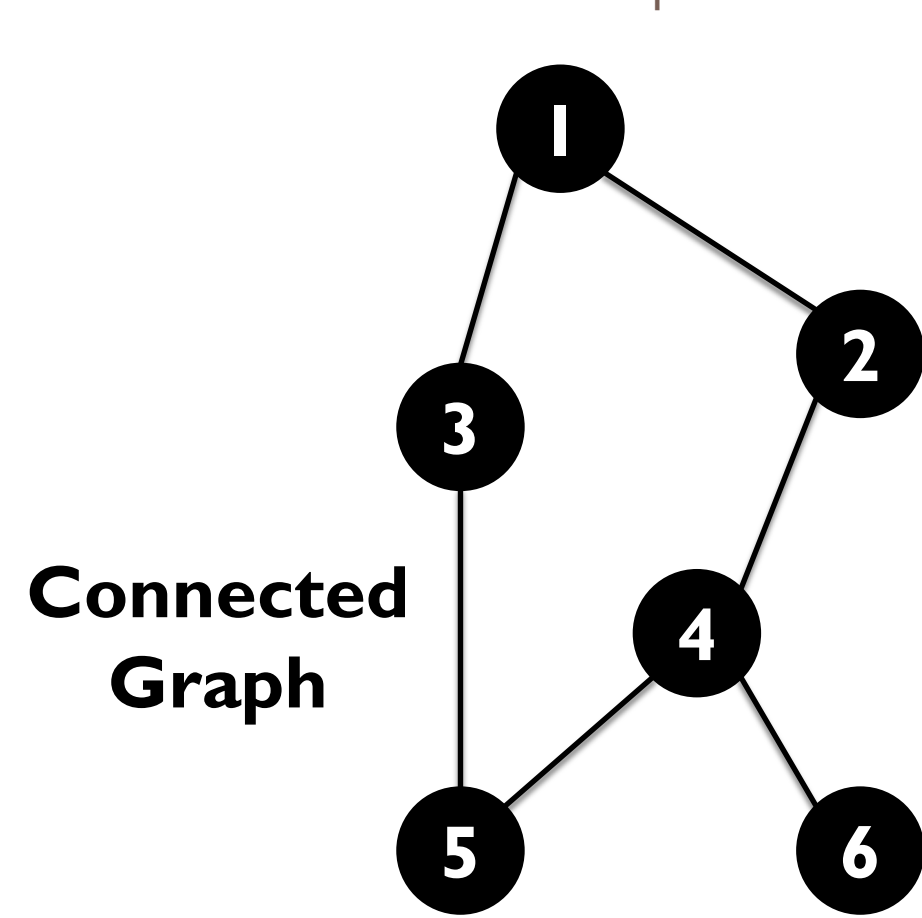
Unweighted Graph



Weighted Graph

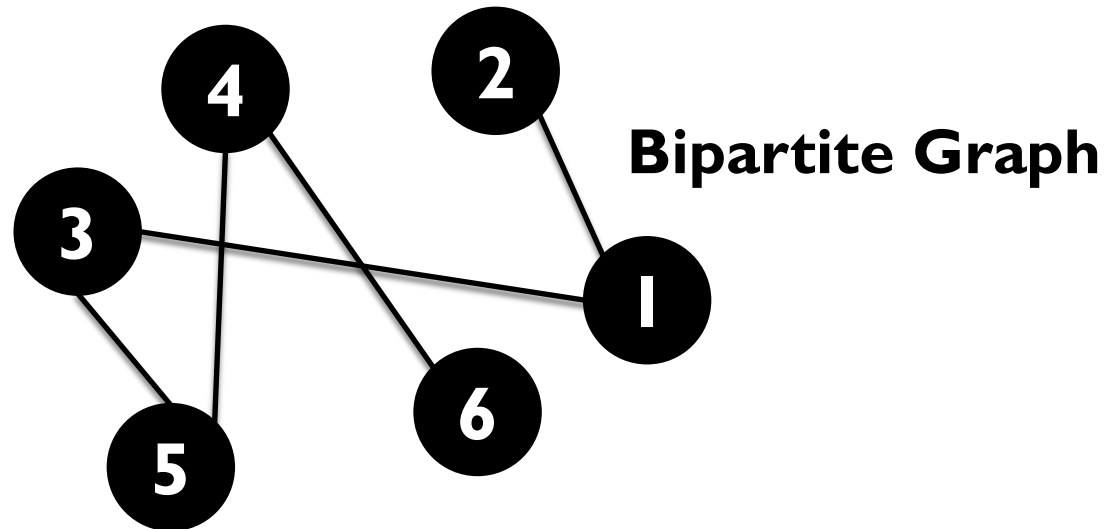
CONNECTED GRAPHS

- For simplicity, most graph algorithms assume the graph is connected. Otherwise, we can run connectivity first, and work on each component.

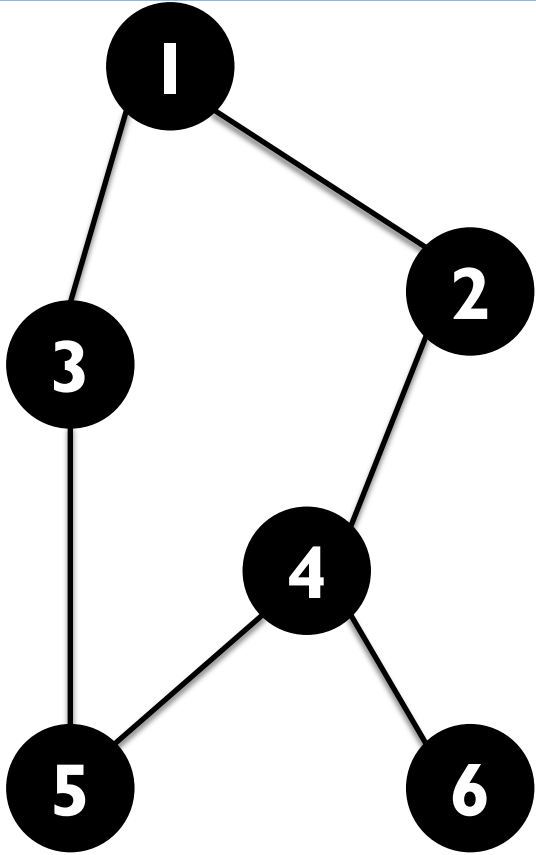


BIPARTITE GRAPH

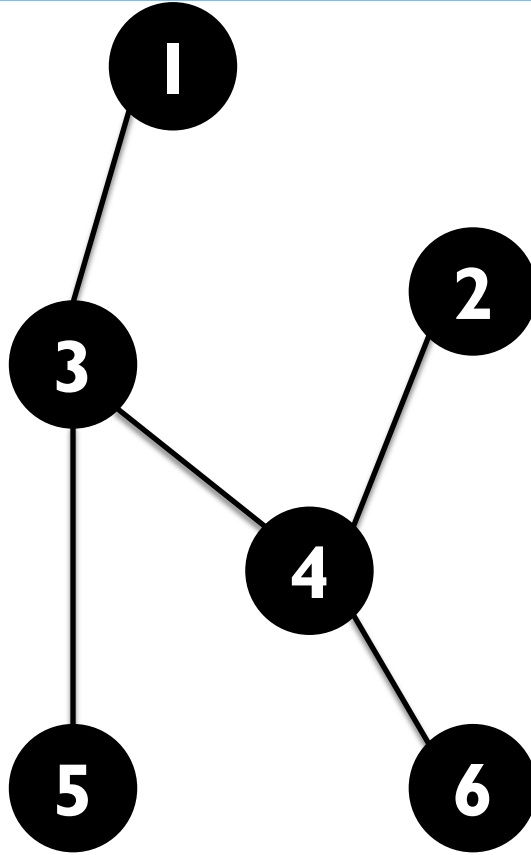
- A graph where the vertices can be partitioned into two subsets:
 - No edges within a subset and all the edges are between two subsets
- Usually, vertices in two subsets have different meanings
 - E.g., students and courses, courses and classrooms, jobs and applicants



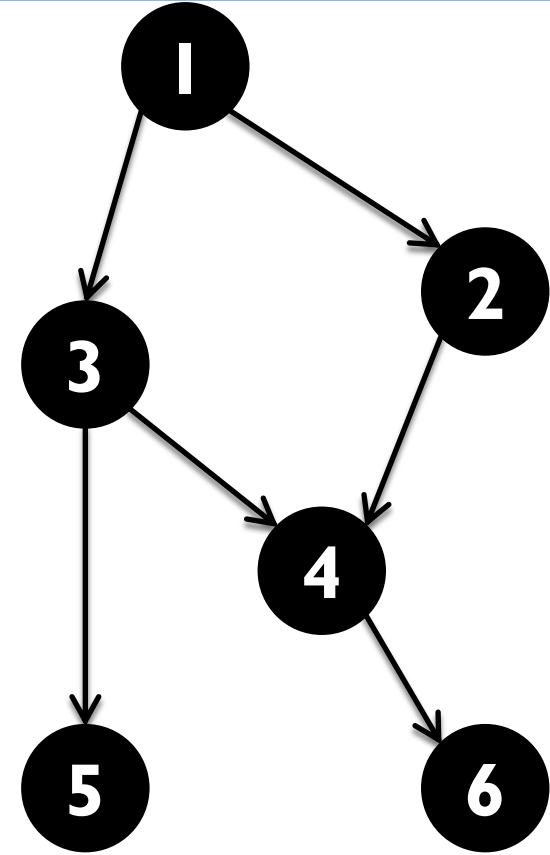
CYCLIC GRAPH



Cyclic Graph



Acyclic Graph



Directed Acyclic Graph (DAG)

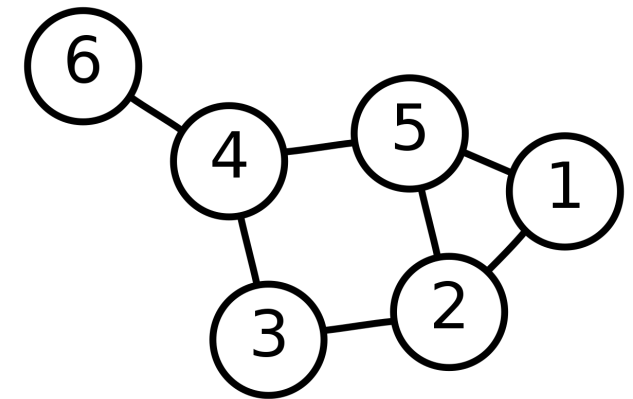
GRAPH REPRESENTATIONS

- Adjacency matrix:
 - Use a 2D matrix ADJ of size $n \times n$
 - If there is an edge between vertices a and b, $ADJ[a][b] = w_{a,b}$ (for unweighted graphs $ADJ[a][b] = 1$)
 - If there is not an edge between a and b, $ADJ[a][b] = 0$
 - Takes too much space ($O(n^2)$)
- Adjacency list:
 - Create n singly linked lists, whose root nodes correspond to vertices in the graph
 - Each linked list holds all neighboring vertices of the vertex represented by the root node

COMPRESSED SPARSE ROW

- Adjacency matrix stores unnecessary information: too sparse!!
- We only need to know when there is an edge, we can infer the other case from non-existence of an edge
- Idea:
 - Only store the existing edge information in an array Edges
 - Keep an extra array, Offset, which indicates which location to look in the Edges array for searching a specific vertex's neighbors

Vertex IDs	1		2		3		4		5		6	
Offset	0		2		5		7		10		13	



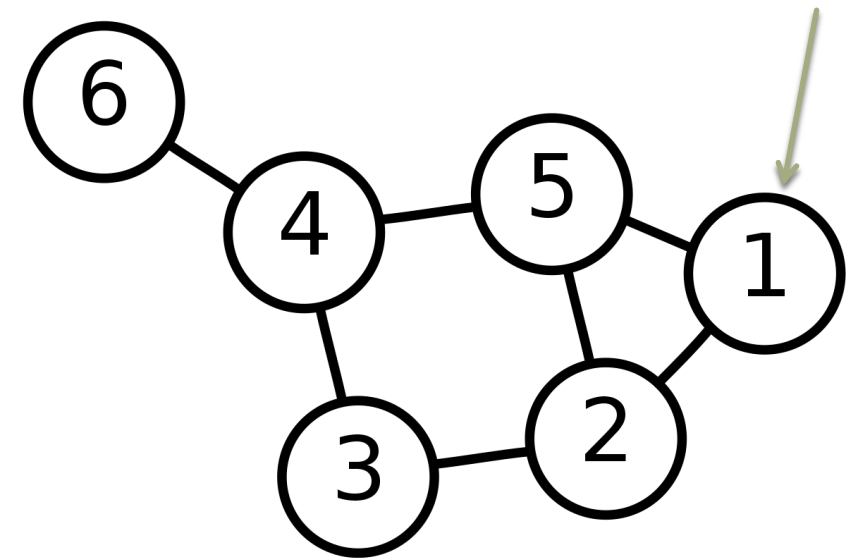
Space: $O(n+m)$

GRAPH TRAVERSALS: BREADTH FIRST SEARCH (BFS)

1	2	3	4	5	6
0	0	0	0	0	0

- Start from 1
- Mark 1 as visited
- Next, add 2 and 5 to the queue (visit them next)

1	2	3	4	5	6
1	0	0	0	0	0



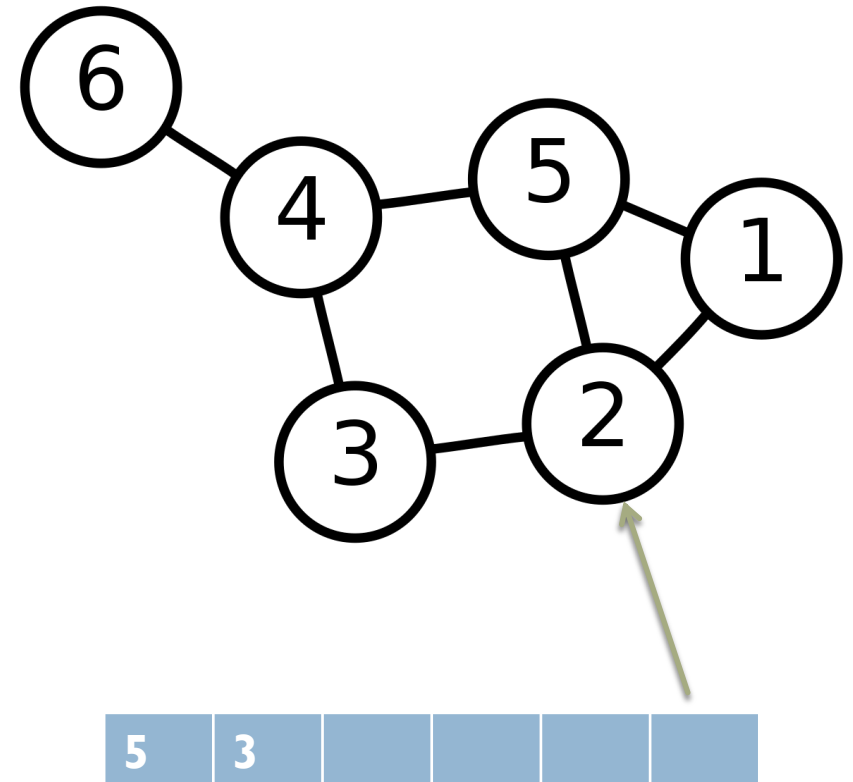
2	5				
---	---	--	--	--	--

GRAPH TRAVERSALS: BREADTH FIRST SEARCH (BFS)

1	2	3	4	5	6
1	0	0	0	0	0

- Pick 2 (Pop 2 from queue)
- Mark 2 as visited
- Add 3 to the queue

1	2	3	4	5	6
1	1	0	0	0	0

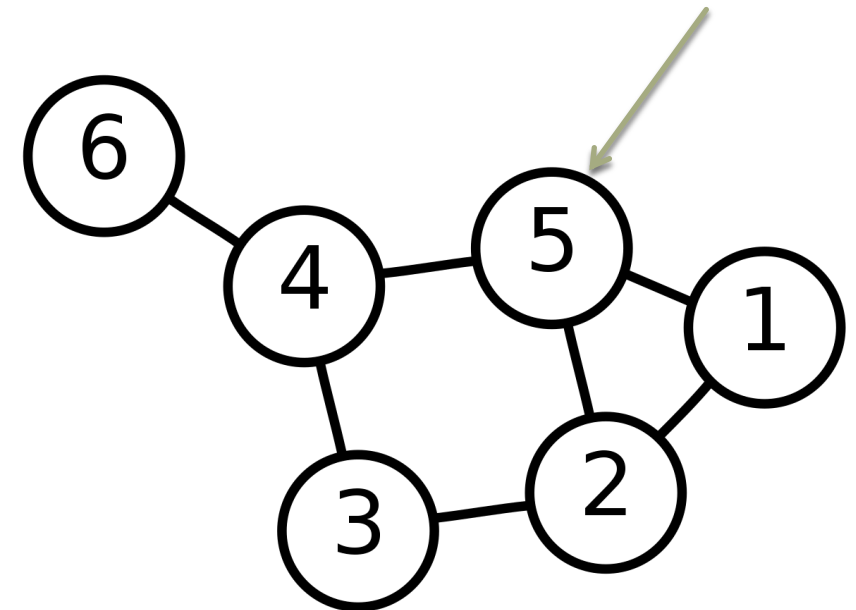


GRAPH TRAVERSALS: BREADTH FIRST SEARCH (BFS)

1	2	3	4	5	6
1	1	0	0	0	0

- Pick 5
- Mark 5 as visited
- Add 4 to the queue

1	2	3	4	5	6
1	1	0	0	1	0



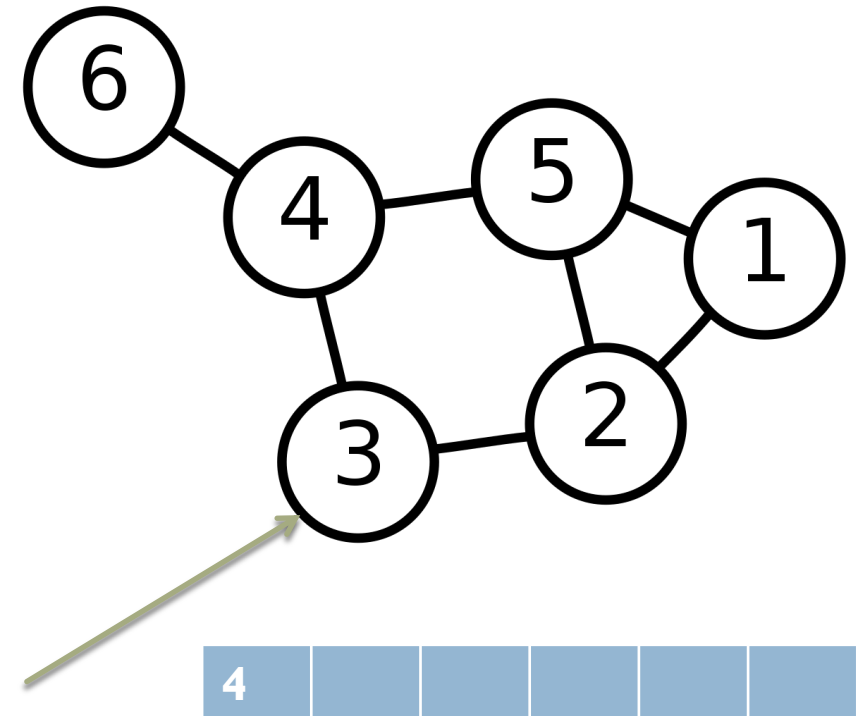
3	4				
---	---	--	--	--	--

GRAPH TRAVERSALS: BREADTH FIRST SEARCH (BFS)

1	2	3	4	5	6
1	1	0	0	1	0

- Pick 3
- Mark 3 as visited

1	2	3	4	5	6
1	1	1	0	1	0

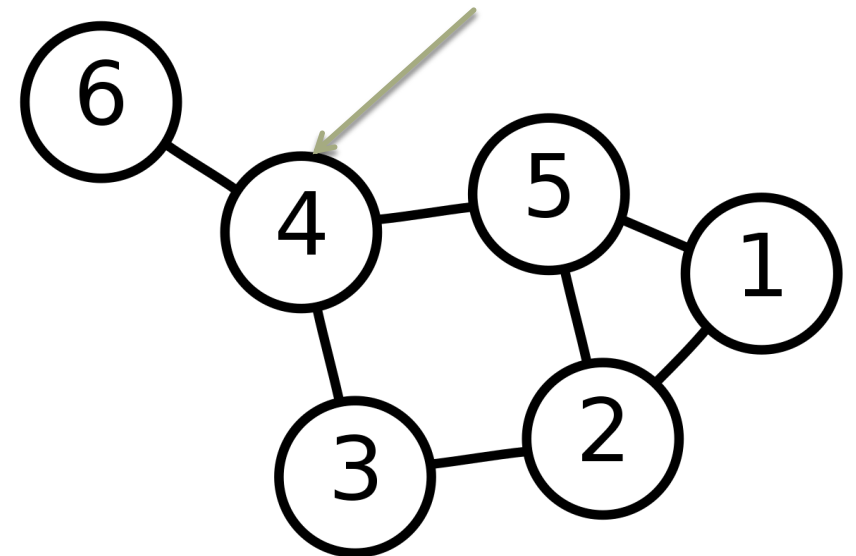


GRAPH TRAVERSALS: BREADTH FIRST SEARCH (BFS)

1	2	3	4	5	6
1	1	1	0	1	0

- Pick 4
- Mark 4 as visited
- Add 6 to the queue

1	2	3	4	5	6
1	1	1	1	1	0



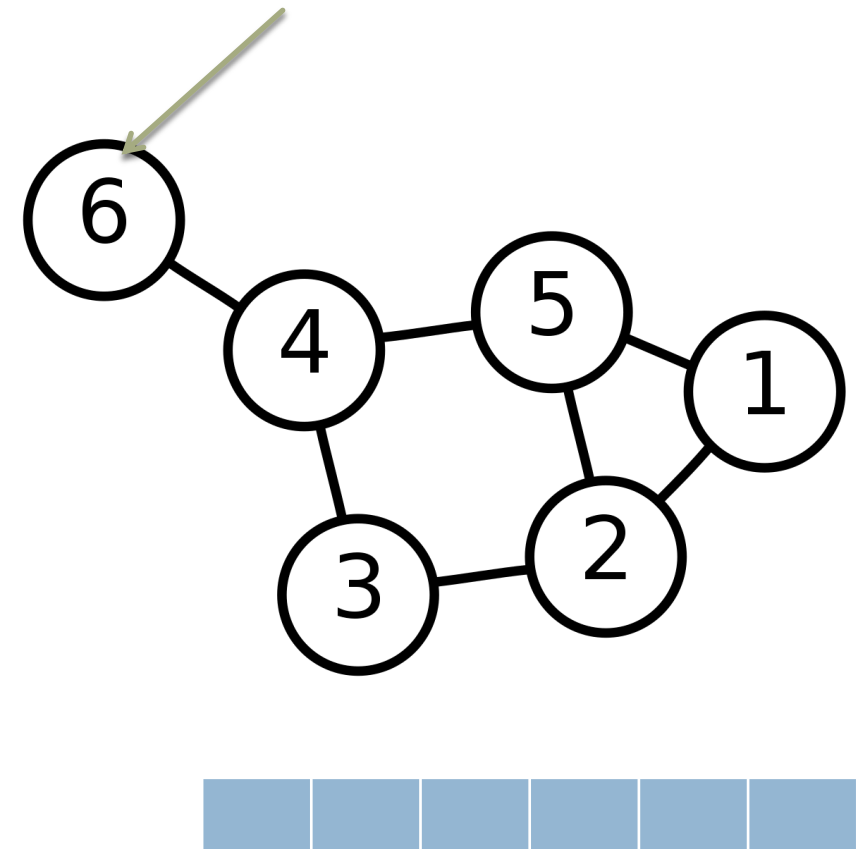
6					
---	--	--	--	--	--

GRAPH TRAVERSALS: BREADTH FIRST SEARCH (BFS)

1	2	3	4	5	6
1	1	1	1	1	0

- Pick 6
- Mark 6 as visited
- Queue is empty, we are done

1	2	3	4	5	6
1	1	1	1	1	1

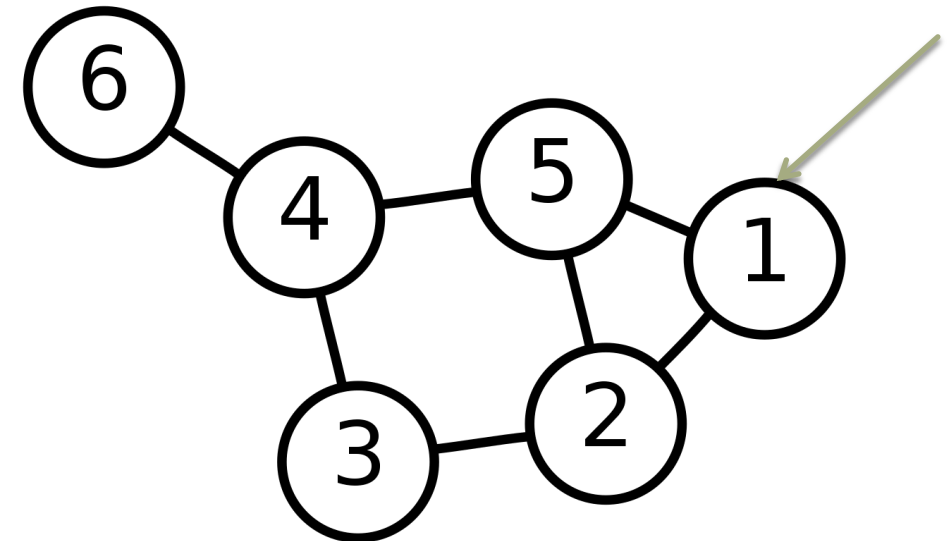


GRAPH TRAVERSALS: DEPTH FIRST SEARCH (DFS)

1	2	3	4	5	6
0	0	0	0	0	0

- Start from 1
- Mark it as visited
- Put 5, 2 in stack

1	2	3	4	5	6
1	0	0	0	0	0



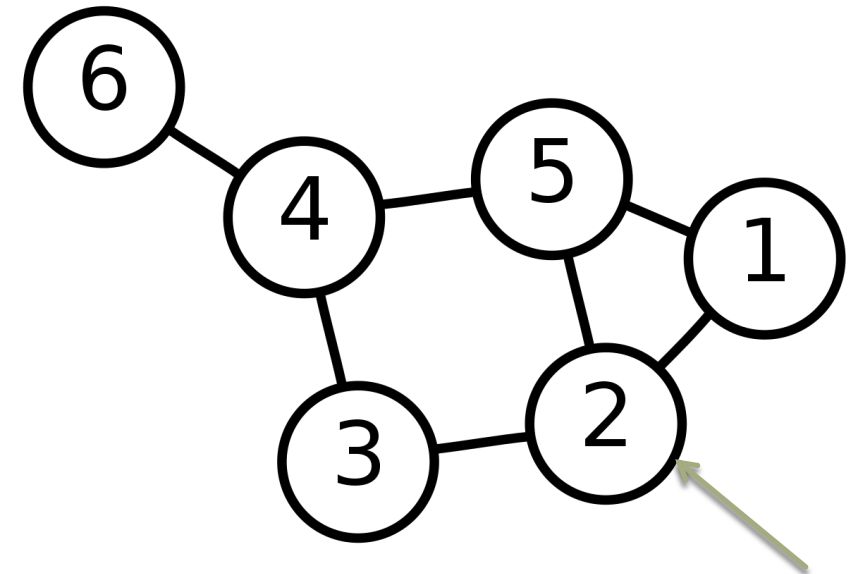
2	5				
---	---	--	--	--	--

GRAPH TRAVERSALS: DEPTH FIRST SEARCH (DFS)

1	2	3	4	5	6
1	0	0	0	0	0

- Visit 2 (Pop 2 from stack)
- Mark it as visited
- Put 3 in stack

1	2	3	4	5	6
1	1	0	0	0	0



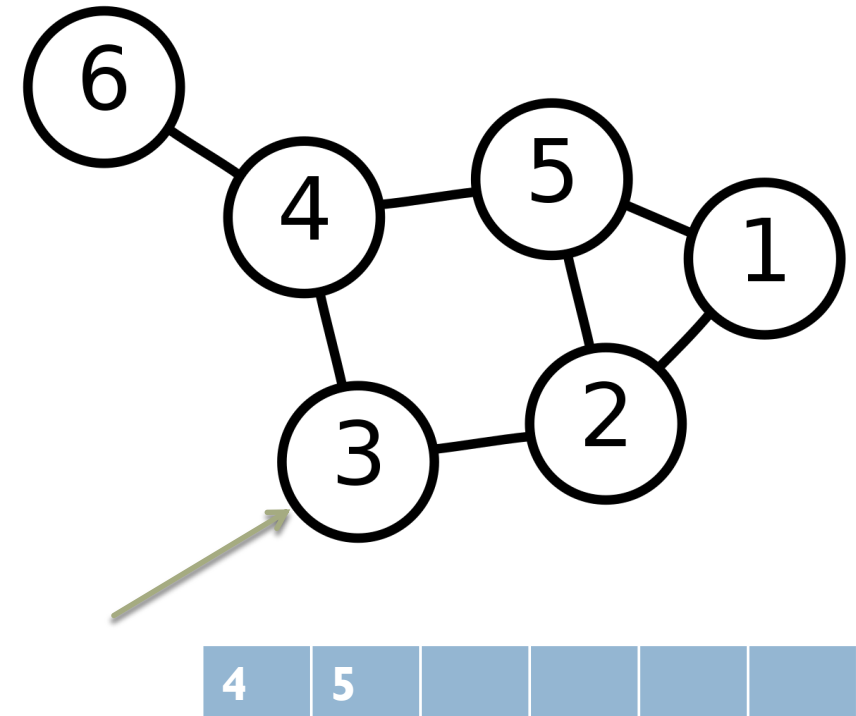
3	5				
---	---	--	--	--	--

GRAPH TRAVERSALS: DEPTH FIRST SEARCH (DFS)

1	2	3	4	5	6
1	1	0	0	0	0

- Visit 3
- Mark it as visited
- Put 4 in stack

1	2	3	4	5	6
1	1	1	0	0	0

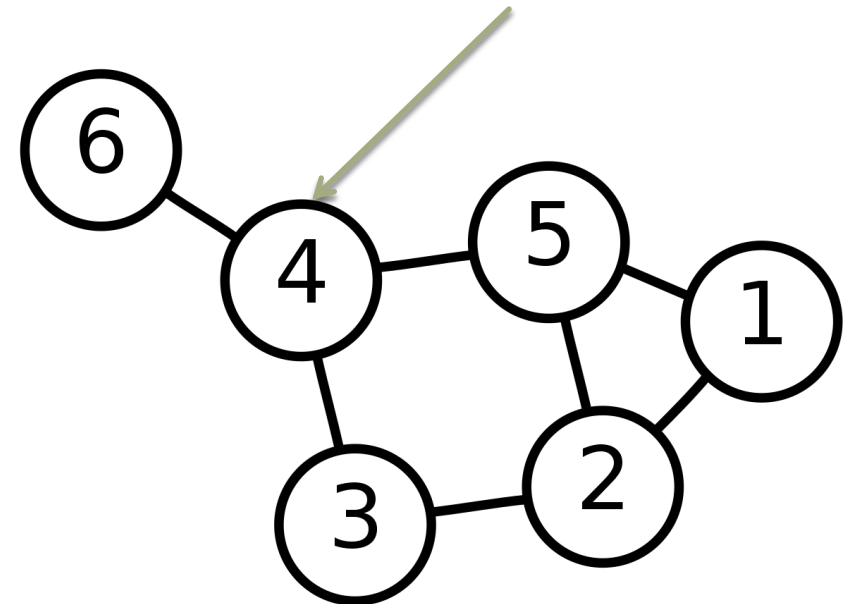


GRAPH TRAVERSALS: DEPTH FIRST SEARCH (DFS)

1	2	3	4	5	6
1	1	1	0	0	0

- Visit 4
- Mark it as visited
- Put 6 in stack

1	2	3	4	5	6
1	1	1	1	0	0



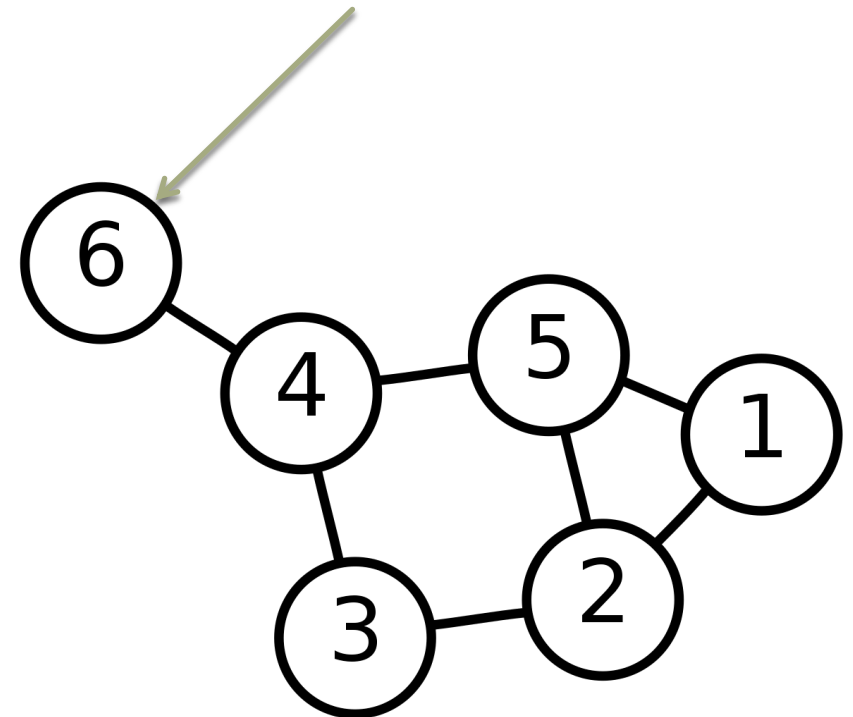
6	5				
---	---	--	--	--	--

GRAPH TRAVERSALS: DEPTH FIRST SEARCH (DFS)

1	2	3	4	5	6
1	1	1	1	0	0

- Visit 6
- Mark it as visited

1	2	3	4	5	6
1	1	1	1	0	1



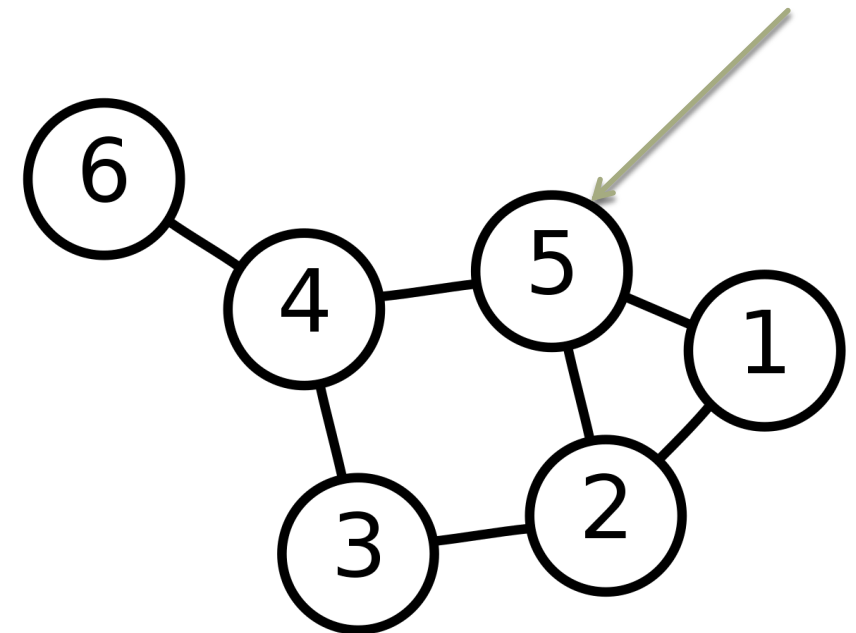
5				
---	--	--	--	--

GRAPH TRAVERSALS: DEPTH FIRST SEARCH (DFS)

1	2	3	4	5	6
1	1	1	1	0	1

- Visit 5
- Mark it as visited
- Stack is empty, we are done

1	2	3	4	5	6
1	1	1	1	1	1





QUESTIONS??