

Efficient Concolic Testing of MPI Applications

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A hierarchical diagram showing the relationship between Concolic Testing and its two subtypes: Concrete Execution and Symbolic Execution. The top node is "Concolic Testing", which is connected by a vertical line to a horizontal line. This horizontal line then branches into two vertical lines, each leading to a node below: "Concrete Execution" on the left and "Symbolic Execution" on the right. Each node is contained within a light blue rounded rectangle, which is itself set against a darker blue rectangular background.

Concolic Testing

Concrete Execution

Symbolic Execution

It Is Popular



- ▶ Programming languages:
 - ▶ Binary machine code, C, Java, and JavaScript.

- ▶ Application types:
 - ▶ web applications, sensor network applications, Unix utilities, database applications, and embedded software, GPU programs, image processing software, and so on

- ▶ Various tools:
 - ▶ KLEE, DART, SAGE, PEX, jCute, CREST, Jalangi, etc.

COMPI [IPDPS 2018]



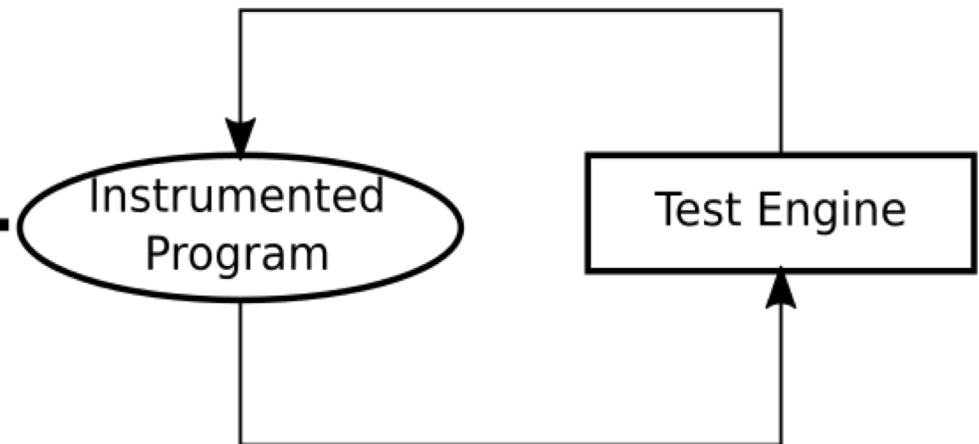
- COMPI is a concolic testing tool for MPI programs with following major features:
 - Deals with basic MPI semantics
 - Deals with high testing cost caused by input values, parallelism, and loops
- COMPI achieves 69-86% branch coverage within a few hours

Concolic Testing

```
int x;  
mark_symbolic(x);  
...  
__load_symbol(x);  
__load_value(0);  
__apply('>');  
if (x > 0) {  
    __log_constraint("x > 0");  
...  
} else {  
    __log_constraint("x <= 0");  
...  
}
```

(1) Instrumentation

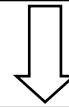
test inputs: # procs, focus, marked_variables



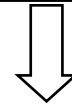
constraints & branch coverage

(2) Iterative testing

Issues of COMPI

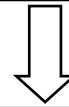


Our Solutions

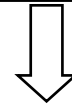


Evaluation

Issues of COMPI



Our Solutions



Evaluation

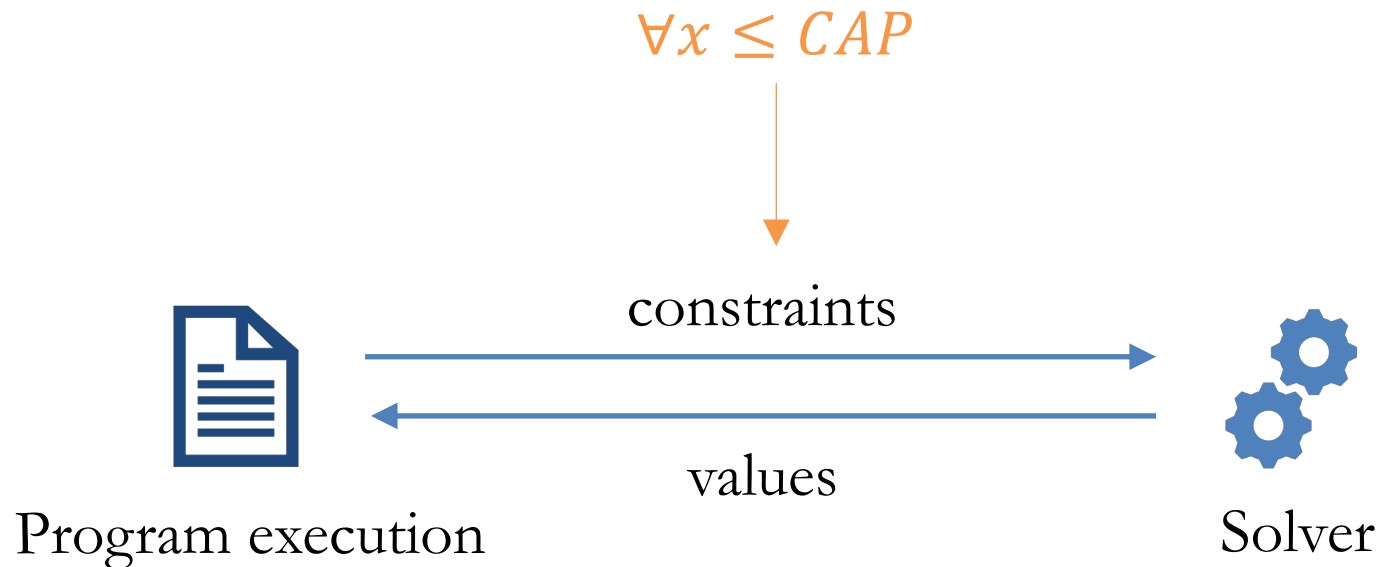
Issues of COMPI



- › **Input generation** does not guarantee cost-effective testing
- › **Floating point** data types and operations are not supported

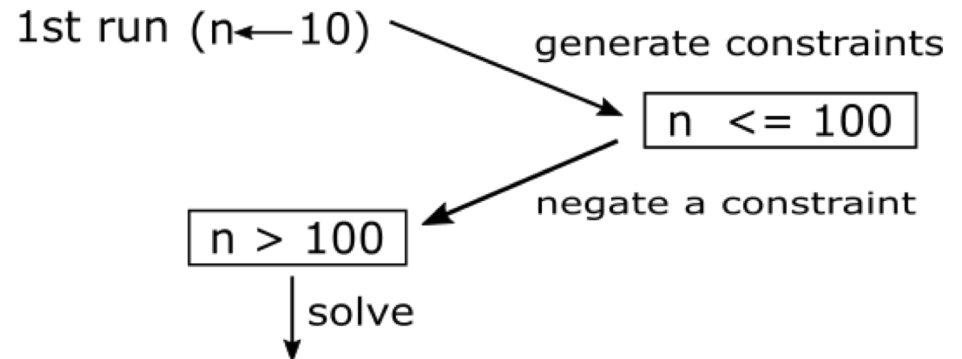
Issue I

- › Larger input values \rightarrow Longer execution
- › Solution of COMPI: **Input Capping**



Example

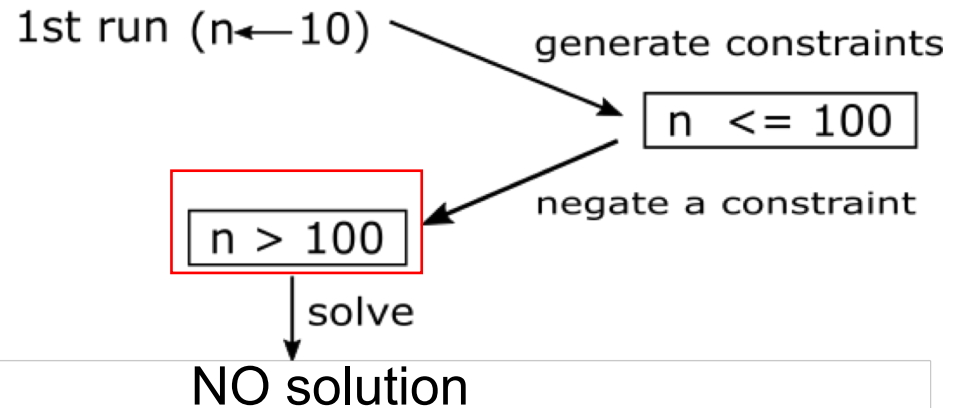
```
main () {  
    MPI_Init();  
    // n denotes the width  
    // of square matrices  
    int n;  
    if (n <= 100)  
        small_matrix_multi();  
    else  
        large_matrix_multi();  
    MPI_Finalize();  
}
```



An MPI program performing matrix multiplication.

Small Cap

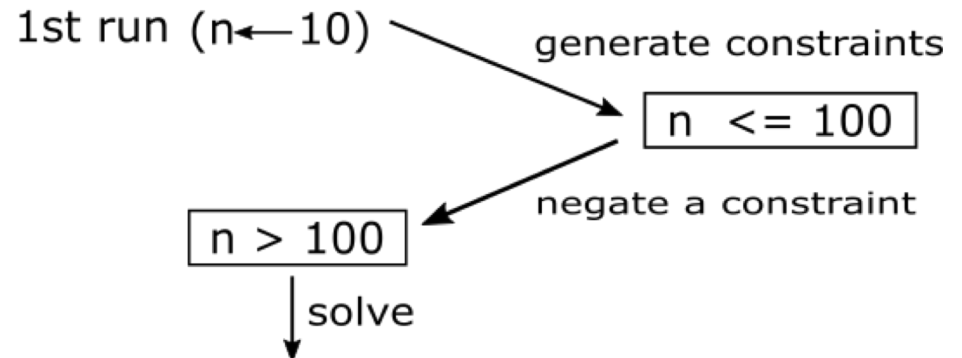
```
main () {  
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    int n;  
    if (n <= 100)  
        small_matrix_multi();  
    else  
        large_matrix_multi();  
    MPI_Finalize();  
}
```



$CAP = 50 \rightarrow$ Fail to cover the *else* branch

Big Cap

```
main () {  
    MPI_Init();  
    // n denotes the width  
    // of square matrices  
    int n;  
    if (n <= 100)  
        small_matrix_multi();  
    else  
        large_matrix_multi();  
    MPI_Finalize();  
}
```

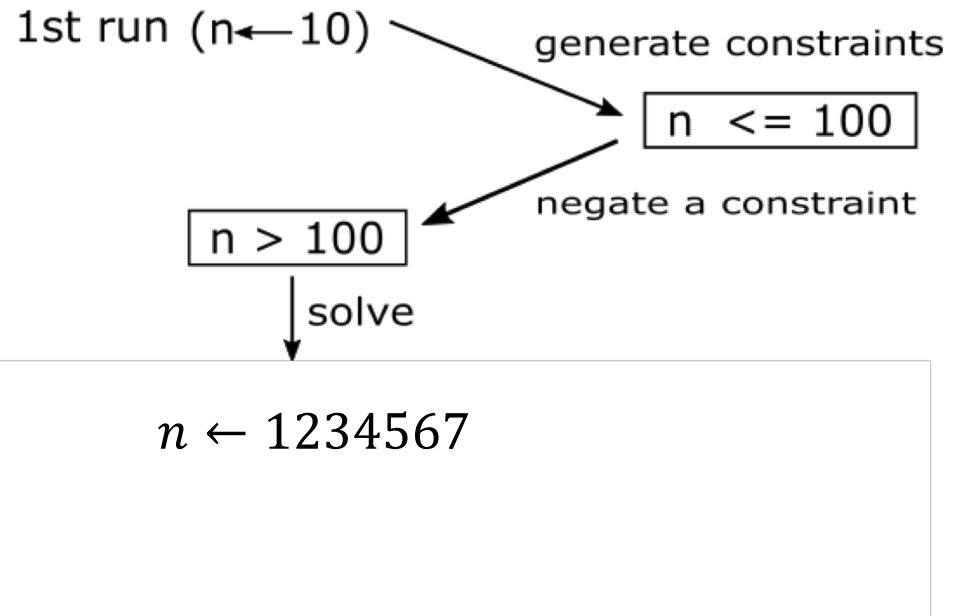


$n \leftarrow C$, where $100 < C \leq 500$

$CAP = 500 \rightarrow$ High testing cost

No Input Capping

```
main () {  
    MPI_Init();  
    // n denotes the width  
    // of square matrices  
    int n;  
    if (n <= 100)  
        small_matrix_multi();  
    else  
        large_matrix_multi();  
    MPI_Finalize();  
}
```



No Capping → Execution failure

Issue II

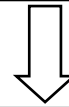


- ▶ Floating-point data types and operations are not supported

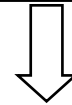
```
main () {  
    ...  
    int a; float b;  
    COMPI_int(a);  
    ...  
    if (b > 1.1) f1();  
    else f2();  
    float c = a * 1.1;  
    if (c > 2) f3();  
    else f4();  
}
```

- Unable to mark *b*
 - Fixing it to a value →
either *f1* or *f2* could not be explored
- Unable to record *a * 1.1*
 - Symbolic representation of *c* is not existing →
either *f3* or *f4* could not be explored

Issues of COMPI



Our Solutions



Evaluation

Our Solutions



- › **Input tuning** → cost effective testing
- › **Floating-point extension** → exploration of branches related to the use of floating-point arithmetic

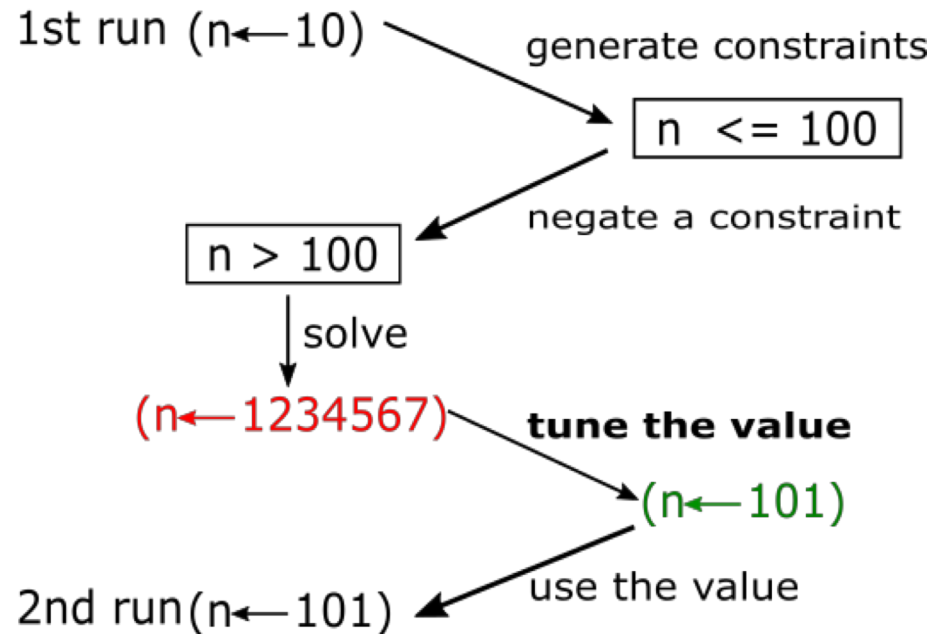
Our Solutions



- **Input tuning** → cost effective testing
- **Floating-point extension** → exploration of branches related to the use of floating-point calculations
 - Constraint solving using reals instead of floating-point numbers → faster constraint solving

Input Tuning

```
main () {  
    MPI_Init();  
    // n denotes the width  
    // of square matrices  
    int n;  
    if (n <= 100)  
        small_matrix_multi();  
    else  
        large_matrix_multi();  
    MPI_Finalize();  
}
```



Input Tuning

Binary search of *upper* in $(0, 1234567)$ satisfying:
 $\{n > 100\} \cup \{n \leq \textit{upper}\}$ is solvable
AND $\{n > 100\} \cup \{n \leq \textit{upper} - 1\}$ is unsolvable

Tuning

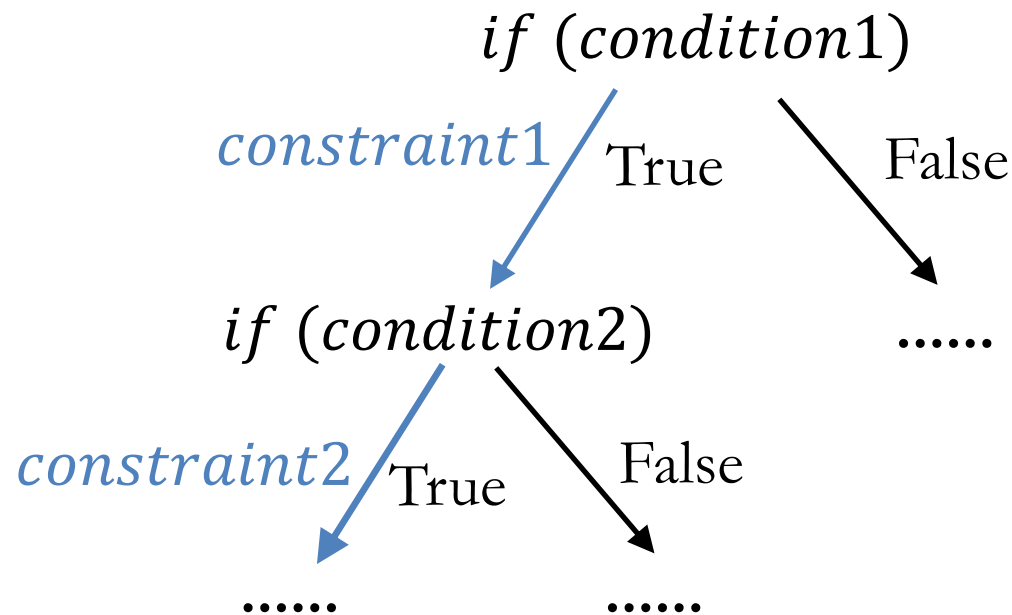
$n \leftarrow 1234567$

Solution by solver for $\{n > 100\}$

$n \leftarrow 101$

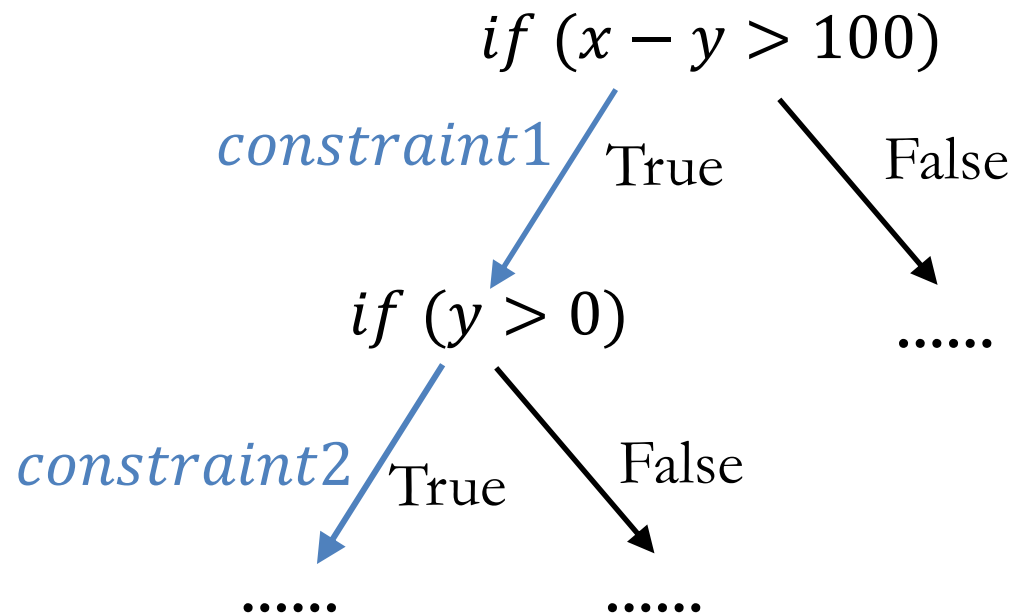
Tuned solution

Input Tuning for Multi-variable Multi-constraint Case



Need to solve {*constraint1*, *constraint2*}

Input Tuning for Multi-variable Multi-constraint Case



Need to solve $\{x - y > 100, y > 0\}$

Input Tuning for Multi-variable Multi-constraint Case



$\{x = 4321, y = 1234\}$

Stage I Tuning

$\{x = 4321, y = 1234, x \leq upper1, y \leq upper1\}$
 $\min\{upper\} = 102$

Stage II Tuning

$\{x = 4321, y = 1234, x \leq upper1, y \leq upper2\}$
 $\min\{upper2\} = 1$

Tuning for $\{x - y > 100, y > 0\}$

$\{x = 102, y = 1\}$

Input Tuning -- Summary



- ▶ Stage I avoids too large values being generated for **ALL** variables appearing in dependent constraints
- ▶ Stage II ensures the smallest value is generated for the **SINGLE** variable appearing in the target constraint based on Stage I

Our Solutions



- › Input tuning → cost effective testing
- › **Floating-point extension** → exploration of branches related to the use of floating-point arithmetic

Floating-Point Extension



- ▶ Two floating-point data types supported: *float*, *double*
- ▶ The extension adopts the design methodology of symbolic reasoning for integers
 - ▶ Instrument floating-point operations
 - ▶ Records only *linear constraints*
 - ▶ Non-linear constraints are simplified using concrete values, e.g., $x * y$ is recorded as $C * x$ with C being the concrete value of y

Floating-Point Extension

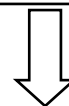
- ▶ Two solvers: Real v.s. Float
 - ▶ Accuracy: Real < Float
 - ▶ Solving speed: Real > Float

Real v.s. Float based on 100 iterative tests of a synthetic program that compares expression e with constant 0.

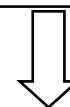
$e =$	x	$x + y$	$x + y + z$
Cost (float)	31.4s	75.0s	91.2s
Cost (real)	8.2s	8.1s	8.2s

3.8-11.1X faster

Limitations of COMPI



Our Solutions



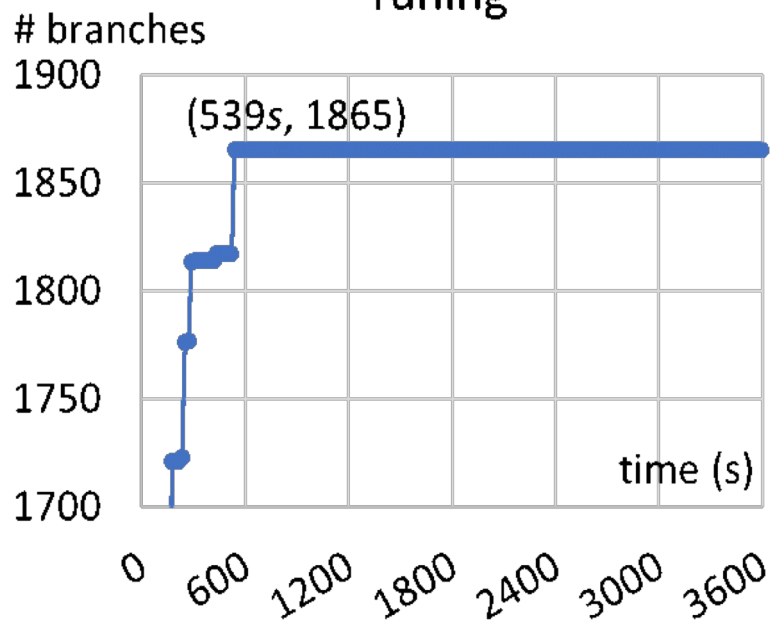
Evaluation

Evaluation

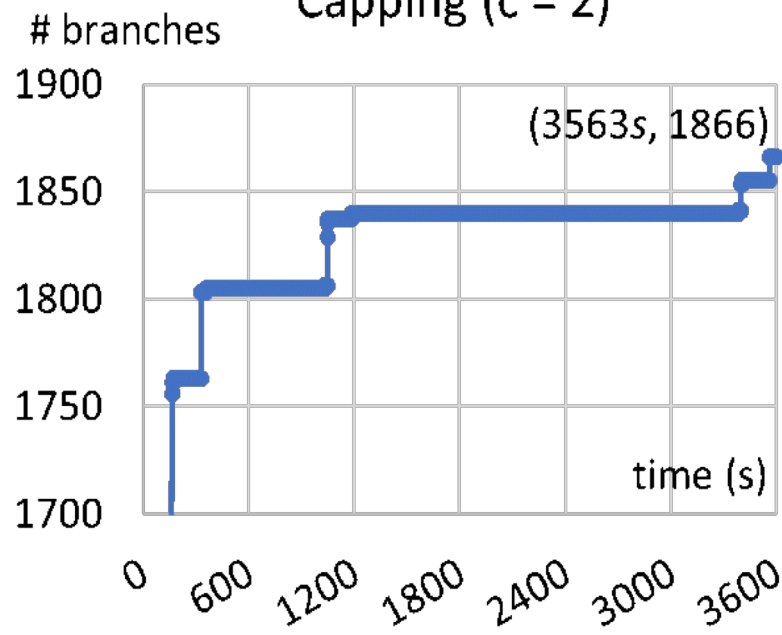


- › Input tuning is evaluated using HPL, IMB-MPI1, and SUSY-HMC
- › Floating-point extension is evaluated using SUSY-HMC
- › One hour testing at each configuration
- › Initial input values are 1 for all variable in the first test
- › In the evaluation of input capping, we selects the same cap for all variables

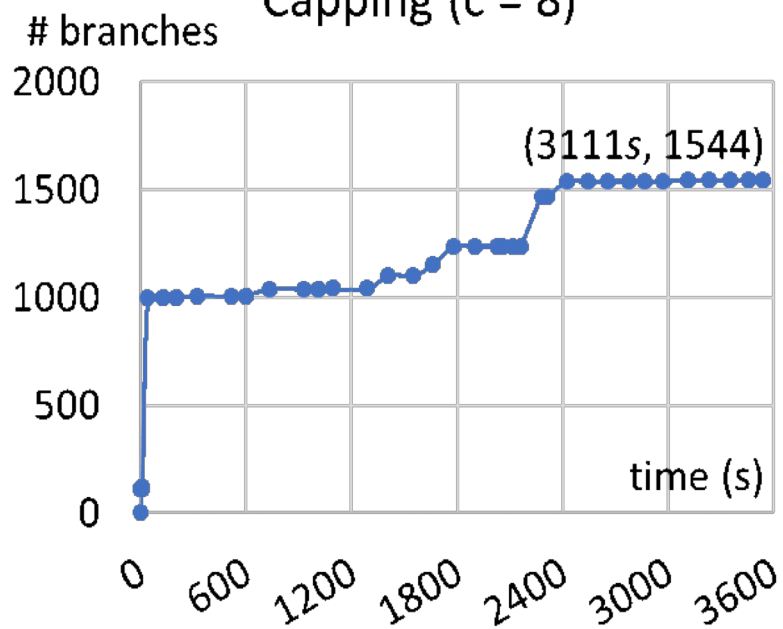
Tuning



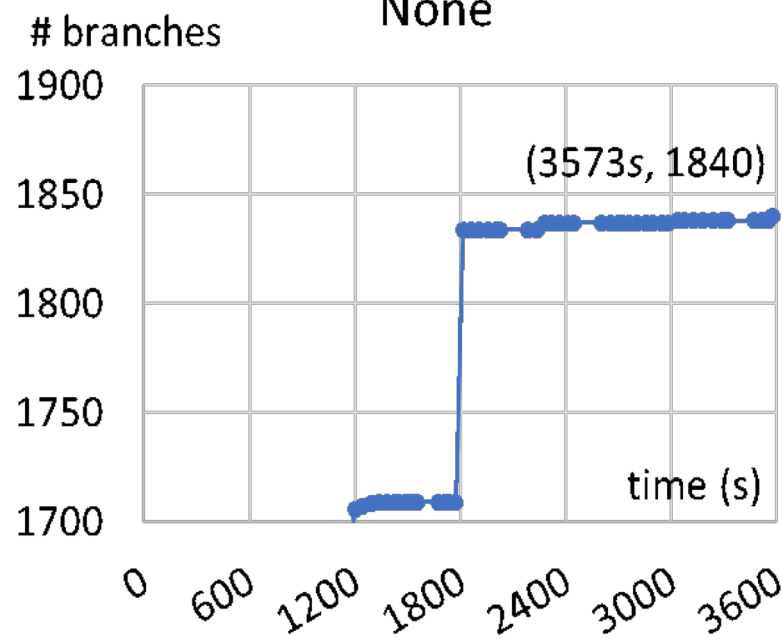
Capping (c = 2)



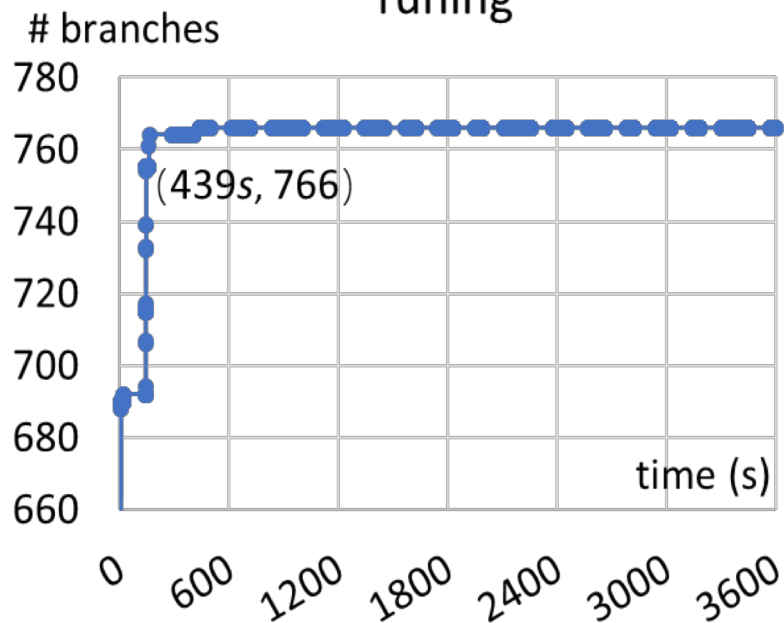
Capping (c = 8)



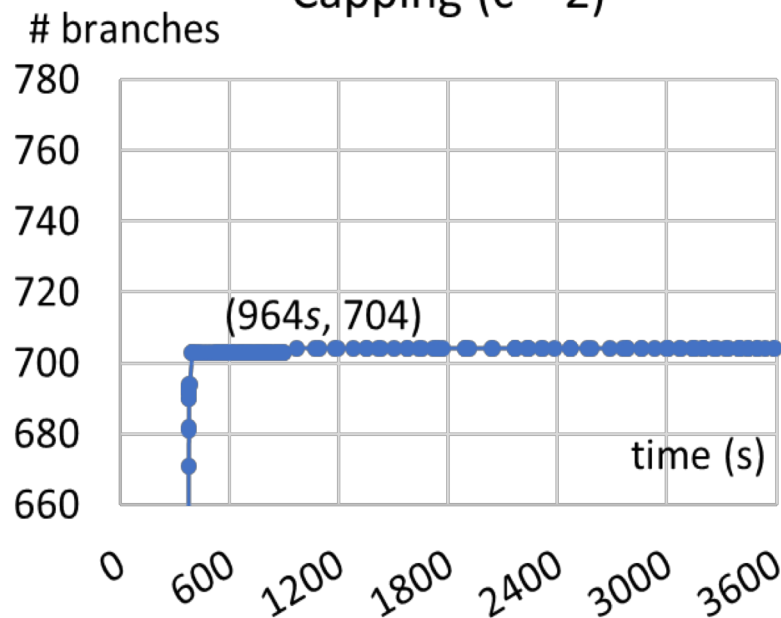
None



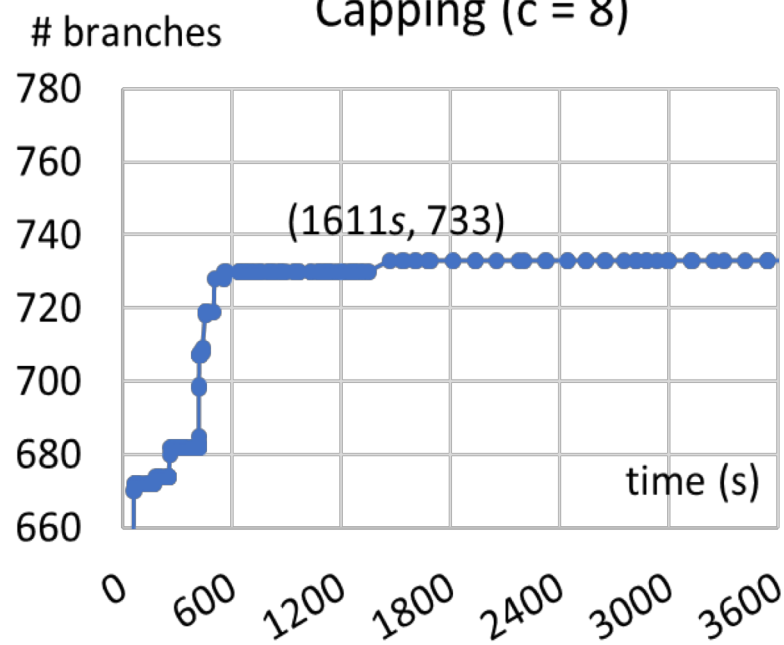
Tuning



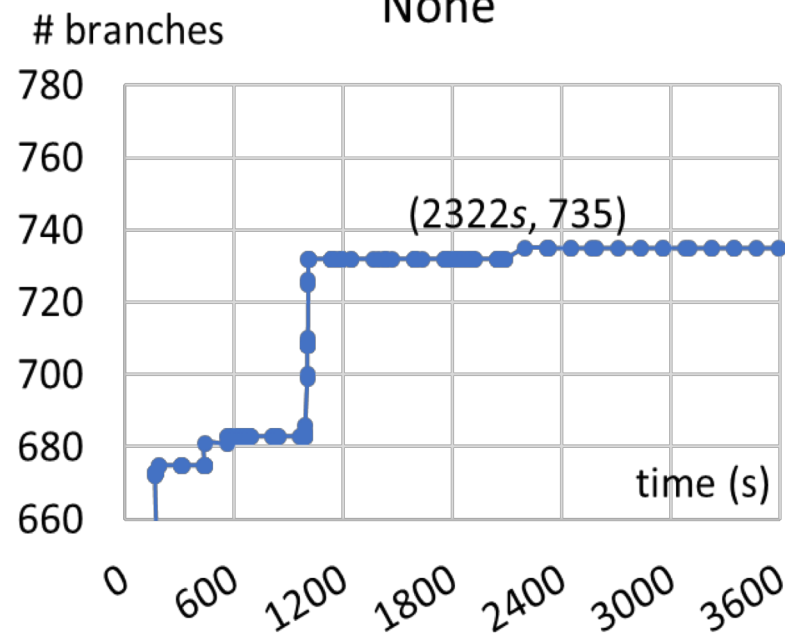
Capping (c = 2)



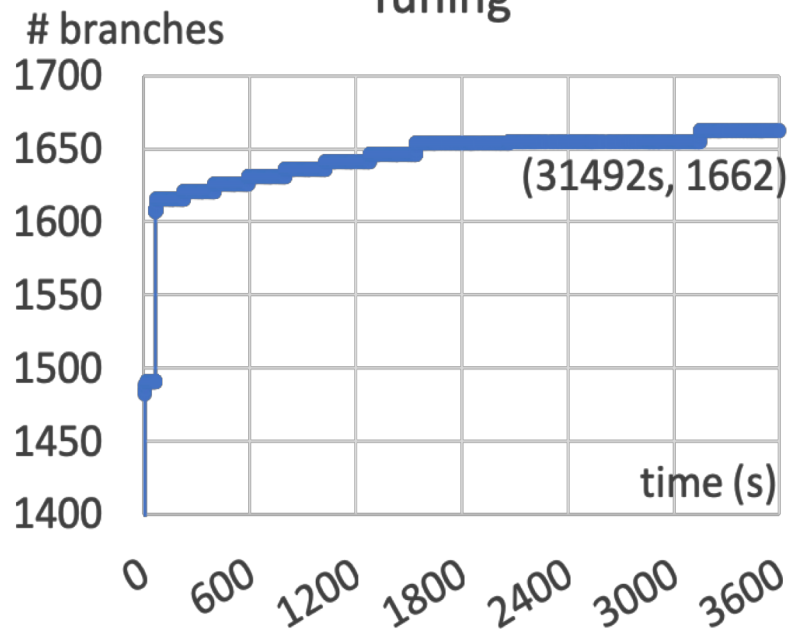
Capping (c = 8)



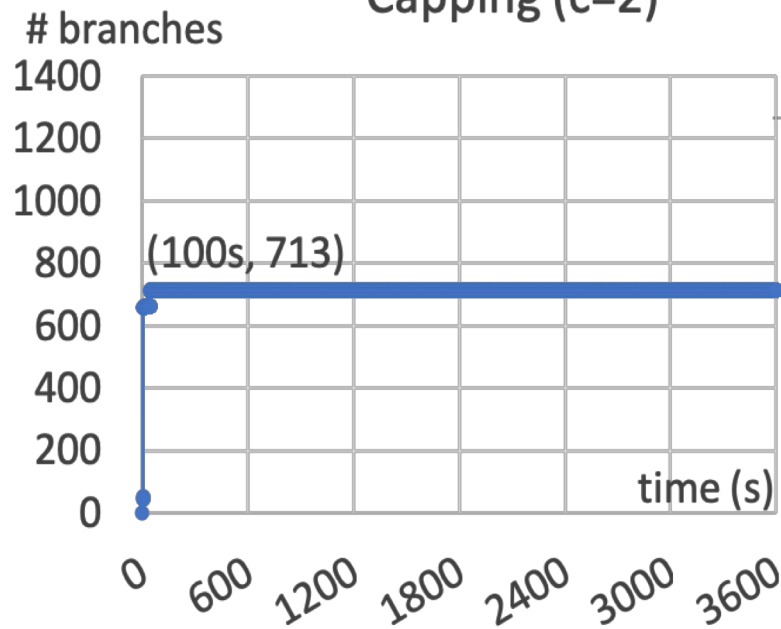
None



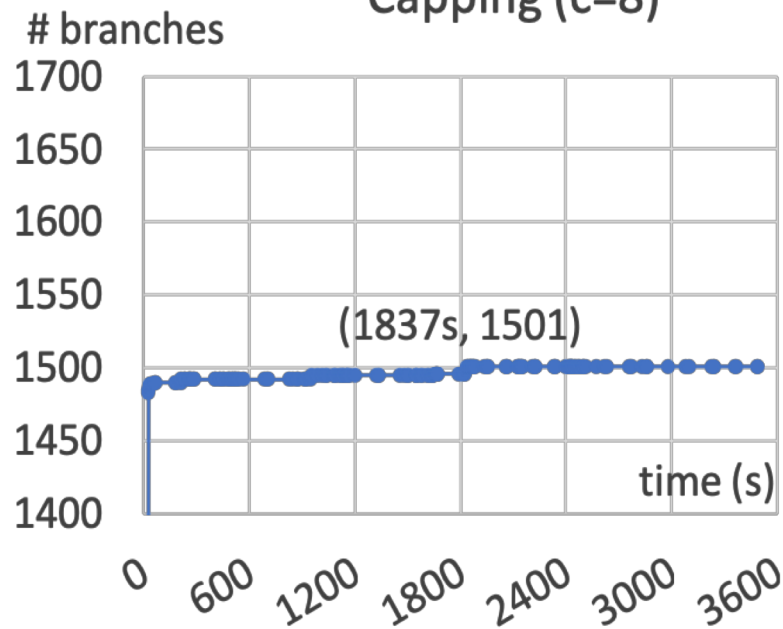
Tuning



Capping (c=2)



Capping (c=8)

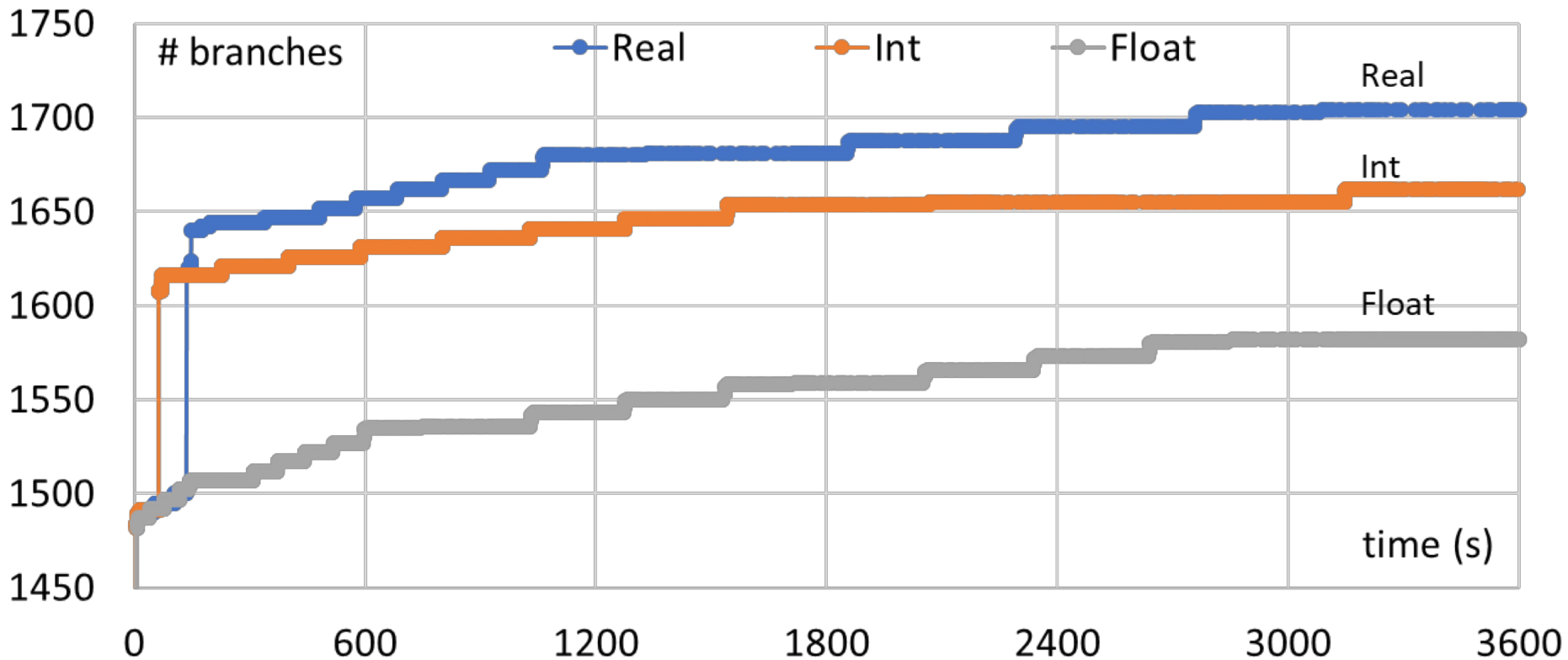


Input Tuning



- ▶ **10-minute** coverage (input tuning) \geq **1-hour** coverage (other methods)
- ▶ SUSY-HMC: 1-hour coverage (input tuning) is about **1.2-2.3X** higher than 1-hour coverage (other methods)

Floating-point Extension



Floating-point Extension



- ▶ $\text{Real (1704)} > \text{Int (1662)} > \text{Float (1582)}$
- ▶ Constraint solving time of Real (1.7%) < Constraint solving time of Float (10.9%)

Conclusion



- Input tuning
 - **10-minute** coverage (input tuning) \geq **1-hour** coverage (other methods)
 - SUSY-HMC: 1-hour coverage (input tuning) is about **1.2-2.3X** higher than 1-hour coverage (other methods)

- Floating-point Extension
 - Floating-point extension using reals achieve 42-122 more branches than the other two

Thank you!