Dynamic Taint Analysis

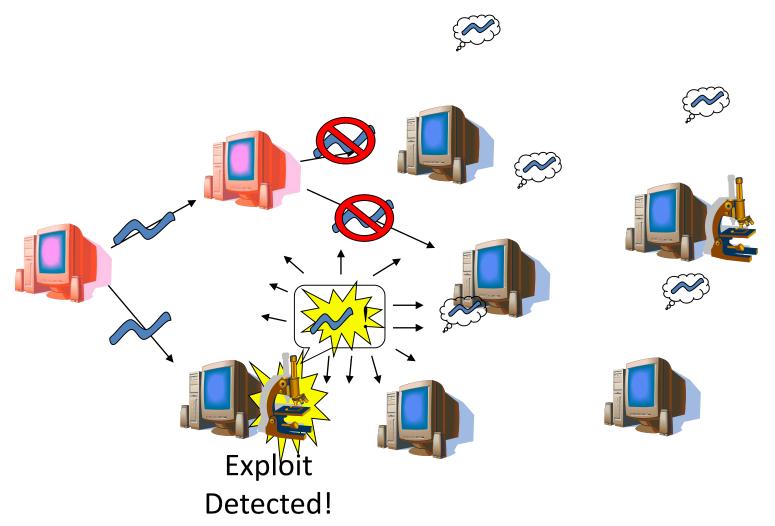
Dynamic Taint Analysis for Automatic Detection, Analysis and Signature Generation of Exploits on Commodity Software

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Appeared in NDSS'06

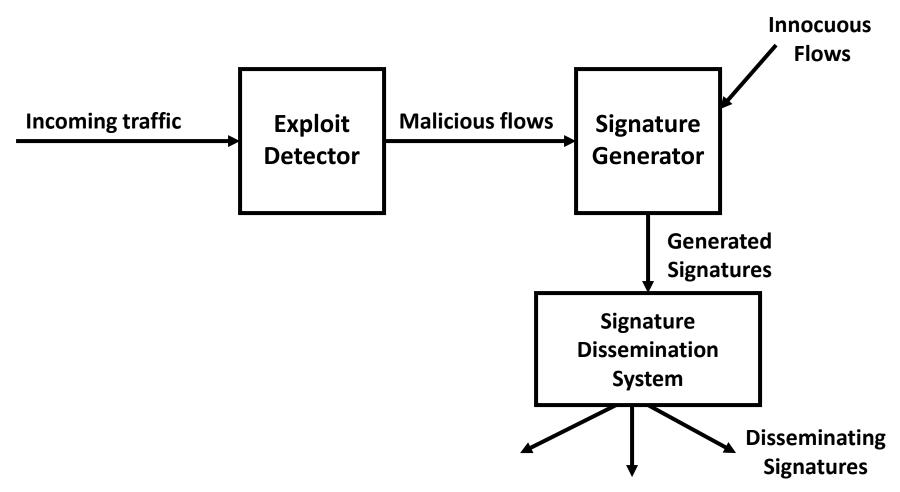
Problem: Internet Worms

- Propagate by exploiting vulnerable software
- No human interaction needed to spread
- Able to rapidly infect vulnerable hosts
 - Slammer scanned 90% of Internet in 10 minutes
- Need automatic defense against new worms

Automatic Worm Defense



Architecture



Common Traits of Software Exploits

- Most known exploits are overwrite attacks
- Attacker's data overwrites sensitive data
- Common overwrite vulnerabilities:
 - Buffer overflows
 - Format string
 - Double-free
- Common overwrite targets:
 - Return address
 - Function pointer

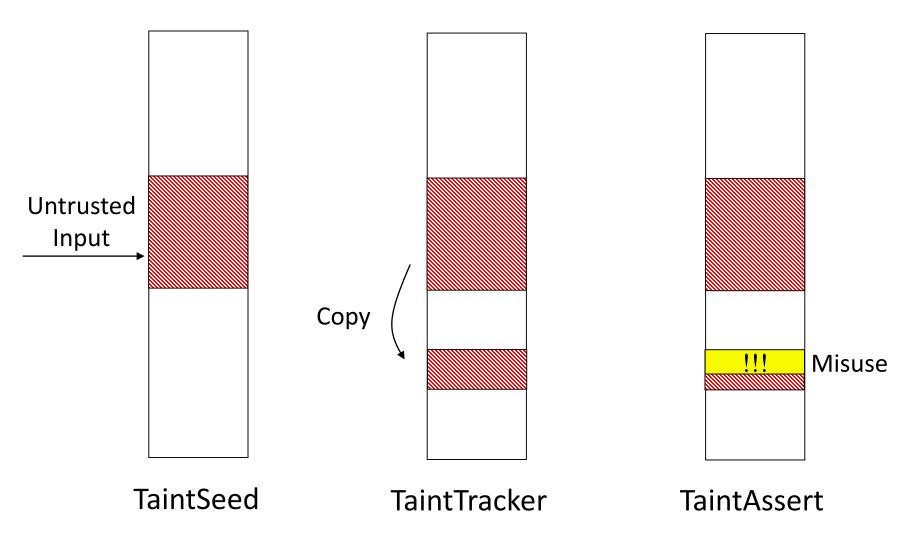
Approach: Dynamic Taint Analysis

- Hard to tell if data is sensitive when it is written
 - Binary has no type information
- Easy to tell it is sensitive when it is used
- Approach: Dynamic Taint Analysis:
 - Keep track of tainted data from untrusted sources
 - Detect when tainted data is used in a sensitive way
 - e.g., as return address or function pointer

Design & Implementation: TaintCheck

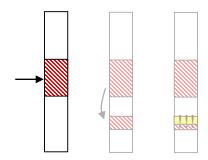
- Use Valgrind to monitor execution
 - Instrument program binary at run-time
 - No source code required
- Track a taint value for each location:
 - Each byte of tainted memory
 - Each register

TaintCheck Components



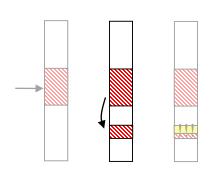
TaintSeed

- Monitors input via system calls
- Marks data from untrusted inputs as tainted
 - Network sockets (default)
 - Standard input
 - File input
 - (except files owned by root, such as system libraries)



TaintTracker

- Propagates taint
- Data movement instructions:
 - e.g., move, load, store, etc.
 - Destination tainted iff source is tainted
 - Taint data loaded via tainted index
 - e.g., unicode = translation_table[tainted_ascii]
- Arithmetic instructions:
 - e.g., add, xor, mult, etc.
 - Destination tainted iff any operand is tainted
- Untaint result of constant functions
 - xor eax, eax



TaintAssert

- Detects when tainted data is misused
 - Destination address for control flow (default)
 - Format string (default)
 - Argument to particular system calls (e.g., execve)
- Invoke Exploit Analyzer when exploit detected

Coverage: Attack Classes Detected

Return Address

Function Pointer

Fn Ptr Offset (GOT)

Jump Address

Format	Strings Stack	Mertlow O	heillon C	orruption or high property of the property of
√	√	N/A	√	
✓	√	✓	✓	
✓	√	✓	✓	
√	by CSRG-Yin Lab	✓	√	1

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Experimental Results: Detects Many Attacks

Vulnerable Program	Overwrite Method	Overwrite Target	Detected		
ATPhttpd	Buffer overflow	Return address	✓		
Synthetic	Buffer overflow	Function pointer	✓		
Synthetic	Buffer overflow	Format string	√		
cfingerd	syslog format string	GOT entry	✓		
wu-ftpd	vsnprintf format string	Return address	√		

Comparison to Previous Mechanisms

- Used Wilander testbed [NDSS03]
 - 20 exploit tests
 - Overwrite Targets: return address, base pointer, function pointer, longjmp buffer
 - Overwrite Techniques: overflow to target, overflow to pointer to target
 - Evaluate previous run-time detection mechanisms

Comparison Results

Mechanisms	Attacks Prevented or Halted					
StackGuard	15%					
Stack Shield	30%					
ProPolice	50%					
Libsafe & Libverify	20%					
TaintCheck	100%					

Other Applications

- Information leakage detection/analysis
- Malware analysis
- Fuzzing
- A base for symbolic execution/concolic testing
- ...

Pointer Tainting

mov eax, [ebx + 4]
 When ebx is tainted, shall eax be tainted?

- Often used for table lookup, e.g.,
 - Convert from ascii to Unicode
 - Convert a date from one format to another

It may cause taint explosion

Over tainting & Under tainting

- xor eax, eax
- sub eax, eax

- Taint granularity is important (bit, byte, word, etc.)
 - Coarser granularity may cause over tainting

Examples of bit-level tainting rules

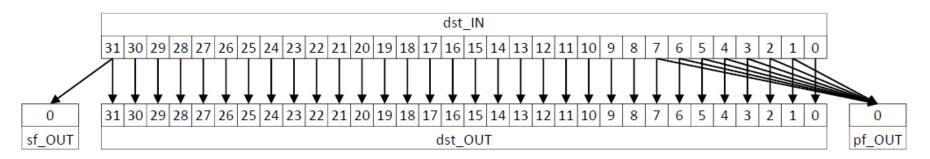


Figure 2: Information flows of *dst* in the or instruction

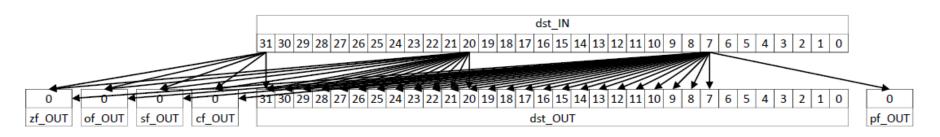


Figure 3: Information flow of bits 7, 20 and 31 of dst in sbb

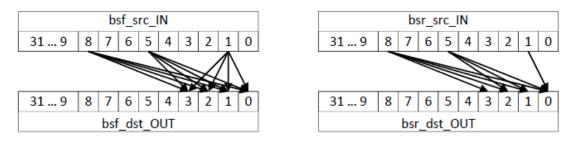


Figure 4: Comparison between *bsf* and *bsr*

Rules for x86 Instructions

	Previously Published Taint Trackers													
Instruction	Inputs	Outputs	# Cases	Runtime	Flow Type	DroidScope[16]	Cat1	libdft[9]	Minemu[10]	Cat2	TEMU[15]	Cat3	Memcheck[32]	SPITA
adc dst, src	dst,src,cf	dsr,src,zf,of,sf,af,cf,pf	4550	1m19s	U	A	A	I	A	A	S	S	U	S
add dst, src	dst,src	dst,src,zf,of,sf,af,cf,pf	4480	1m13s	U	A	A	I	A	A	A	A	S	S
and dst, src	dst,src	dst,src,zf,sf,pf	4288	1m05s	I	A	I	I	A	I	I	S	S	S
dec dst	dst	dst,zf,of,sf,af,pf	1184	20s	U	A	A	I	A	A	A	A	U	S
div rm32	edx,eax,rm32	edx,eax,rm32	9216	95m48s	D	A	A	I	N	A	A	A	A	D
idiv rm32	edx,eax,rm32	edx,eax,rm32	9216	307m04	Α	A	Α	I	N	A	A	Α	A	A
imul1 rm32	eax,rm32	edx,eax,rm32,of,cf	6272	289m51s	U	A	A	I	N	A	A	A	U	U
imul2 dst, rm32	dst,rm32	dst,rm32,of,cf	4224	52m37s	U	A	A	I	N	A	A	A	U	U
imul3 dst, rm32, imm32	rm32,imm32	dst,rm32,imm32,of,cf	6272	53m56s	U	A	A	I	N	A	A	A	U	U
inc dst	dst	dst,zf,of,sf,af,pf	1184	19s	U	A	A	I	A	A	A	A	U	S
mul <i>rm32</i>	eax,rm32	edx,eax,rm32,of,cf	6272	16m02s	U	A	A	I	N	A	A	A	U	U
not dst	dst	dst	1024	15s	I	A	I	I	A	I	I	I	I	I
or dst, src	dst,src	dst,src,zf,sf,pf	4288	1m05s	I	A	I	I	A	I	I	S	S	S
rcl dst, imm8	dst,imm8,cf	dst,imm8,of,cf	1722	42s	Α	Α	Α	N	A	A	Α	Α	Α	S
rer dst, imm8	dst,imm8,cf	dst,imm8,of,cf	1722	42s	A	A	Α	N	A	A	A	A	A	S
rol dst, imm8	dst,imm8	dst,imm8,of,cf	1680	41s	A	Α	Α	N	A	A	A	Α	S	S
ror dst, imm8	dst,imm8	dst,imm8,of,cf	1680	41s	Α	Α	Α	N	A	A	A	Α	S	S
sal dst, imm8	dst,imm8	dst,imm8,zf,of,sf,af,cf,pf	1840	35s	U	A	A	N	A	A	S	S	S	S
sar dst, imm8	dst,imm8	dst,imm8,zf,of,sf,af,cf,pf	1840	34s	D	A	A	N	A	A	S	S	S	S
sbb dst, src	dst,src,cf	dst,src,zf,of,sf,af,cf,pf	4550	1m21s	U	A	A	I*	A*	A	A	A*	A	S
shr dst, imm8	dst,imm8	dst,imm8,zf,of,sf,af,cf,pf	1840	35s	D	A	A	N	A	A	S	S	S	S
sub dst, src	dst,src	dst,src,zf,of,sf,af,cf,pf	4480	1m17s	U	A	A	I*	A*	A*	A*	A*	S	S
xor dst, src	dst,src	dsr,src,zf,sf,pf	4288	1m05s	I	A	I	I*	A*	A*	A*	A*	I	I
bsf dst, src	src	dst,src,zf	2080	31s	A	N	A	I	N	A	Α	A	A	S
bsr dst, src	src	dst,src,zf	2080	31s	S	N	A	I	N	A	A	A	A	S
cmpxchg rm32, r32	eax,rm32,r32	eax,rm32,r32,zf,of,sf,af,cf,pf	9792	2m39s	S	N	E	E	N	E	E	E	E	S
TOTAL			102064	13h52m48s										

How to verify tainting rules

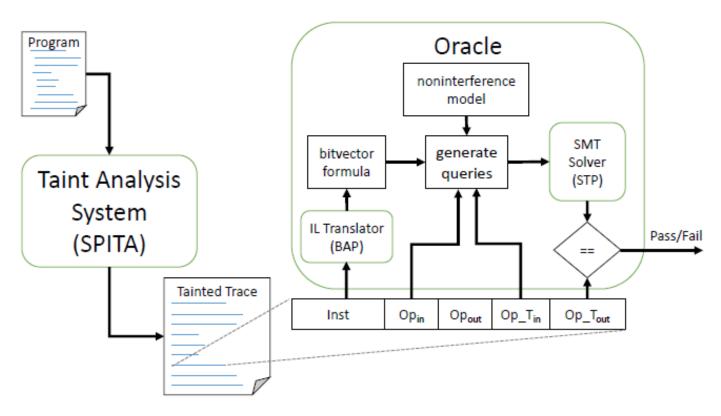


Figure 8: Per-Trace Verification Overview

```
1. // Query for bit [31] of R_EBX:u32
2. R EBX C:u32 = 0x46018902:u32
3. R_EBX_T:u32 = 0x56718e20:u32
4. //Concretization of flags
5. goal:bool = false
 6. R_{EBX:u32} = (R_{EBX}_{01:u32} \& R_{EBX}_{T:u32})
              | (R_EBX_C:u32 & ~R_EBX_T:u32)
7. R EBX:u32 = 0:u32 // sets R EBX to 0
8. //BAP IR for calculating the flags for xor ebx, ebx
9. goal1:u32 = R_EBX:u32 & 0x80000000:u32

    R_EBX:u32 = (R_EBX_O2:u32 & R_EBX_T:u32)

              | (R_EBX_C:u32 & ~R_EBX_T:u32)
11. //Same BAP IR for emulating xor
12. goal2:u32 = R_EBX:u32 & 0x80000000:u32
13. goal:bool = goal1:u32 <> goal2:u32
```

Figure 9: Query to determine whether bit 31 of EBX should be tainted

Table IV: Comparing SPITA with TEMU on tainted shell commands. "n / m" indicates that "n" bytes are tainted, and "m" tainted EIPs are observed.

Command	SPITA	TEMU
dir	207 / 0	639 / 0
cd	146 / 0	616 / 0
cipher c:	929 / 0	3617 / 0
echo hello	660 / 0	3808 / 0
find "jone" a.txt	967 / 0	5684 / 0
findstr /s /i jone ./*	945 / 0	1333 / 0
ls	350 / 3	34923 / 0
cd	306 / 3	301 / 0
cat ./readme	545 / 31	26619 / 0
echo hello	744 / 9	704 / 0
ln -s a.txt nbench	1122 / 35	24707 / 0
mkdir test	551 / 9	23766 / 0