CSE 153
Design of Operating Systems

Fall 2018

Lecture 07: Scheduling

Scheduling Overview

- Scheduler runs when we context switching among processes/threads to pick who runs next
  - Under what situation does this occur?
  - What should it do? Does it matter?
- Making this decision is called scheduling
- Now, we'll look at:
  - The goals of scheduling
  - Starvation
  - Various well-known scheduling algorithms
  - Standard Unix scheduling algorithm

Multiprogramming

- Increase CPU utilization and job throughput by overlapping I/O and CPU activities
- Mechanisms vs. policy
- We have covered the mechanisms
  - Context switching, how and when it happens
  - Process queues and process states
- Now we'll look at the policies
  - Which process (thread) to run, for how long, etc.
- We'll refer to schedulable entities as jobs (standard usage) – could be processes, threads, people, etc.

Scheduling Goals

- Scheduling works at two levels in an operating system

1. Control multiprogramming level –number of jobs loaded into memory
   » Moving jobs to/from memory is often called swapping
   » Long term scheduler: infrequent

2. To decide what job to run next
   » Does it matter? What criteria?
   » Short term scheduler: frequent
   » We are concerned with this level of scheduling

Scheduling

- The scheduler is the OS module that manipulates the process queues, moving jobs to and from
- The scheduling algorithm determines which jobs are chosen to run next and what queues they wait on
- In general, the scheduler runs:
  - When a job switches from running to waiting
  - When an interrupt occurs
  - When a job is created or terminated

Preemptive vs. Non-preemptive scheduling

- We'll discuss scheduling algorithms in two contexts
  - In preemptive systems the scheduler can interrupt a running job (involuntary context switch)
  - In non-preemptive systems, the scheduler waits for a running job to explicitly block (voluntary context switch)

Scheduling Goals

- What are some reasonable goals for a scheduler?
- Scheduling algorithms can have many different goals:
  - CPU utilization
  - Job throughput (# jobs/unit time)
  - Turnaround time ($T_{finish} – T_{start}$)
    - » Normalized turnaround time = Turnaround time/process length
  - Avg Waiting time (Avg(T_{wait}): avg time spent on wait queues)
  - Avg Response time (Avg(T_{ready}): avg time spent on ready queue)
- Batch systems
  - Strive for job throughput, turnaround time (supercomputers)
- Interactive systems
  - Strive to minimize response time for interactive jobs (PC)

Starvation

Starvation is a scheduling "non-goal":

- Starvation: process prevented from making progress because other processes have the resource it requires
  - Resource could be the CPU, or a lock (recall readers/writers)
- Starvation usually a side effect of the sched. Algorithm
  - E.g., a high priority process always prevents a low priority process from running on the CPU
  - E.g., one thread always beats another when acquiring a lock
- Starvation can be a side effect of synchronization
  - E.g., constant supply of readers always blocks out writers

First In First Out (FIFO)

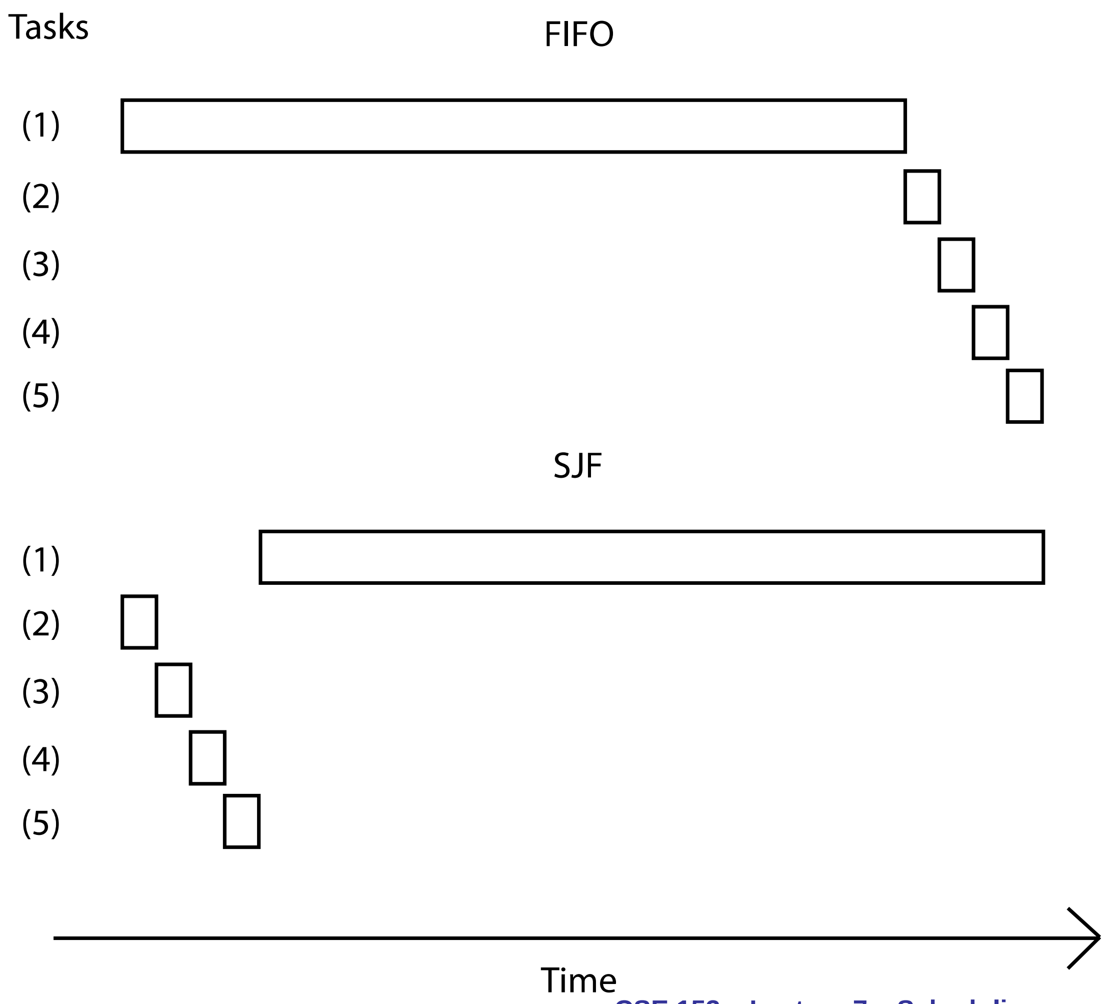
- Schedule tasks in the order they arrive
  - Continue running them until they complete or give up the processor
- Example: memcached
  - Facebook cache of friend lists, …
- On what workloads is FIFO particularly bad?
  - Imagine being at supermarket to buy a drink of water, but get stuck behind someone with a huge cart (or two!)
    - » …and who pays in pennies!
  - Can we do better?

Shortest Job First (SJF)

- Always do the task that has the shortest remaining amount of work to do
  - Often called Shortest Remaining Time First (SRTF)

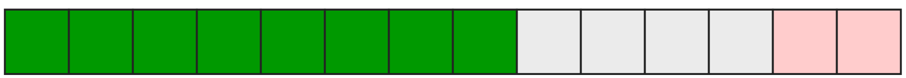
- Suppose we have five tasks arrive one right after each other, but the first one is much longer than the others
  - Which completes first in FIFO? Next?
  - Which completes first in SJF? Next?

FIFO vs. SJF

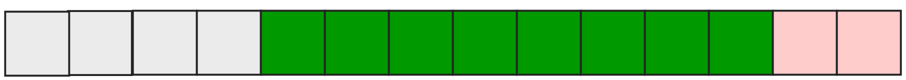


Whats the big deal? Don’t they finish at the same time?

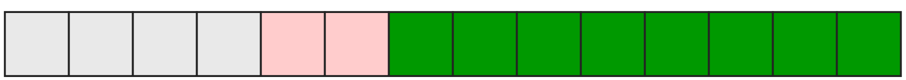
SJF Example



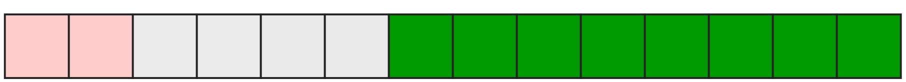
$AWT = (8 + (8+4)+(8+4+2))/3 = 11.33$



$AWT = (4 + (4+8)+(4+8+2))/3 = 10$



$AWT = (4+ (4+2)+(4+2+8))/3 = 8$



$AWT = (2 + (2+4)+(2+4+8))/3 = 7.33$

SJF

- Claim: SJF is optimal for average response time
  - Why?
- For what workloads is FIFO optimal?
  - For what is it pessimal (i.e., worst)?
- Does SJF have any downsides?

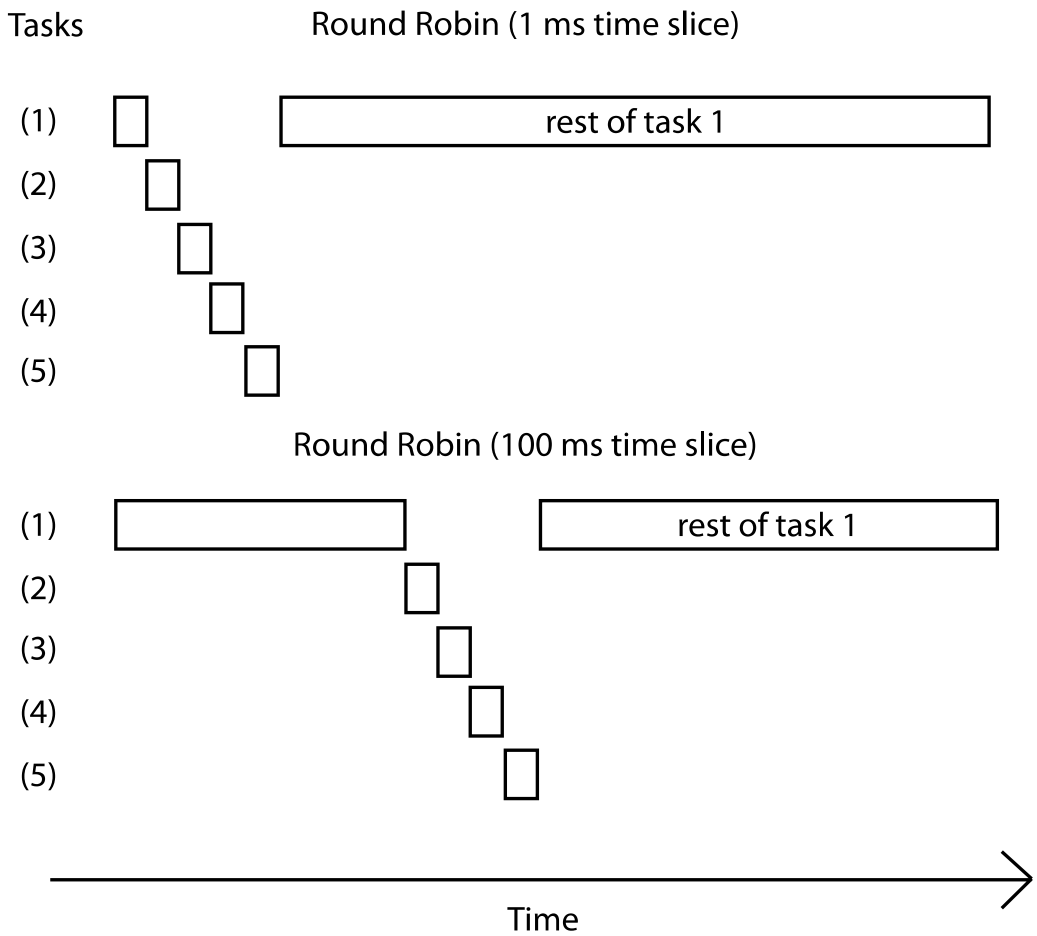
Shortest Job First (SJF)

- Problems?
  - Impossible to know size of CPU burst
    - » Like choosing person in line without looking inside basket/cart
  - How can you make a reasonable guess?
  - Can potentially starve
- Flavors
  - Can be either preemptive or non-preemptive
  - Preemptive SJF is called shortest remaining time first (SRTF)

Round Robin

- Each task gets resource for a fixed period of time (time quantum)
  - If task doesn't complete, it goes back in line
- Need to pick a time quantum
  - What if time quantum is too long?
    - » Infinite?
  - What if time quantum is too short?
    - » One instruction?

Round Robin

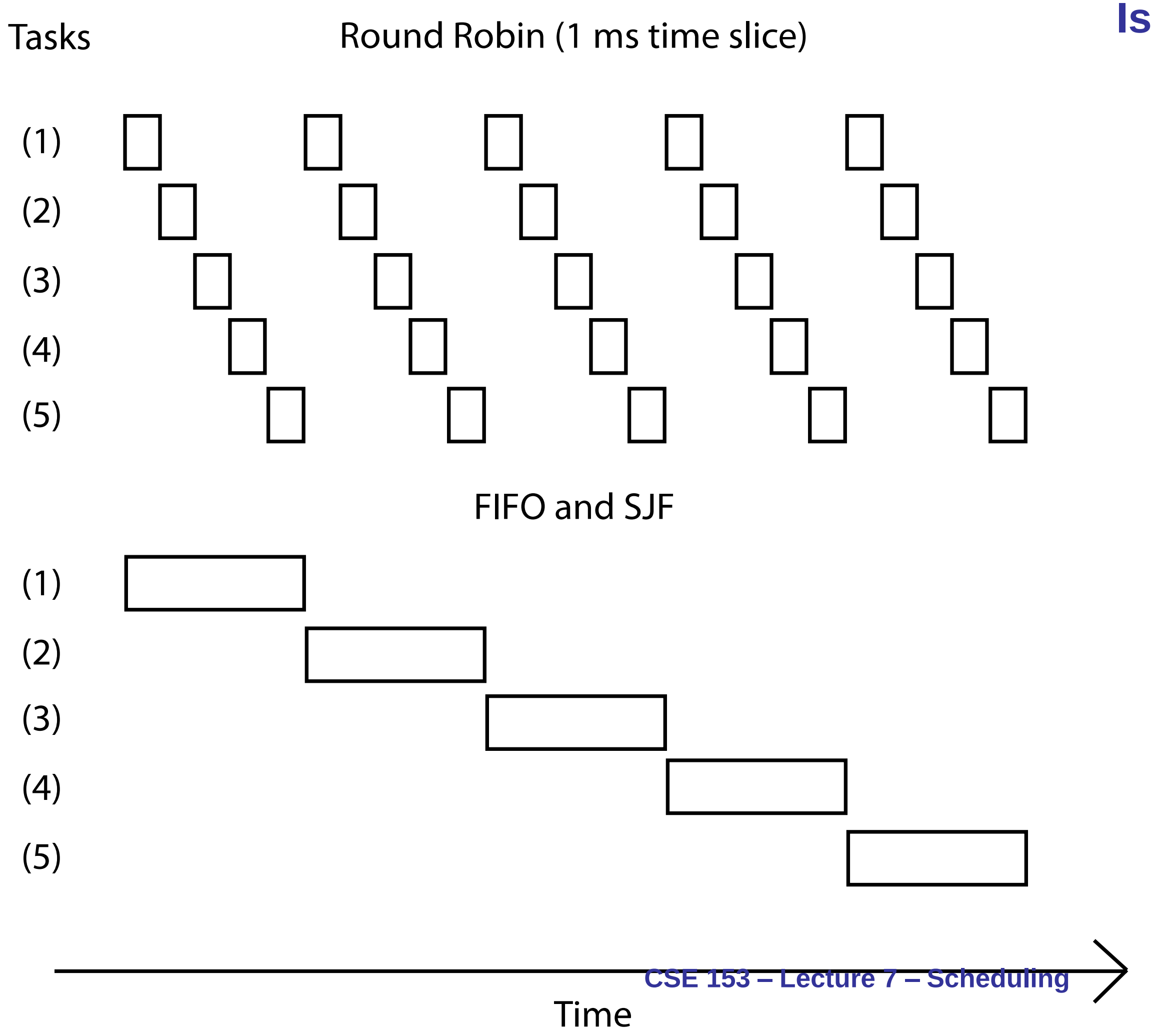

Tasks
Round Robin (1 ms time slice)
(1)
rest of task 1
(2)
(3)
(4)
(5)
Round Robin (100 ms time slice)
(1)
rest of task 1
(2)
(3)
(4)
(5)
Time

Round Robin vs. FIFO

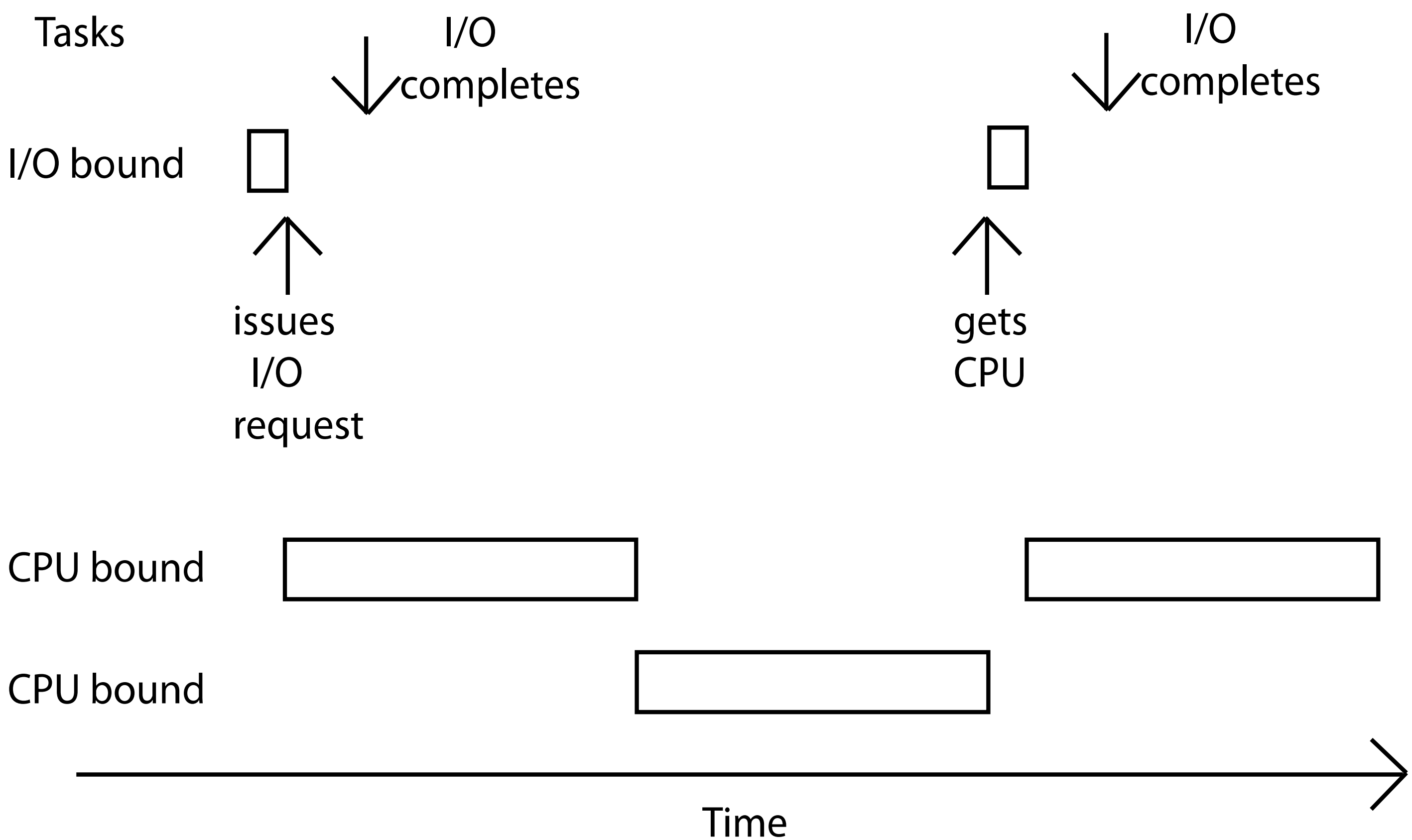
- Many context switches can be costly
- Other than that, is Round Robin always better than FIFO?

Round Robin vs. FIFO

Is Round Robin always fair?



Mixed Workload



Max-Min Fairness

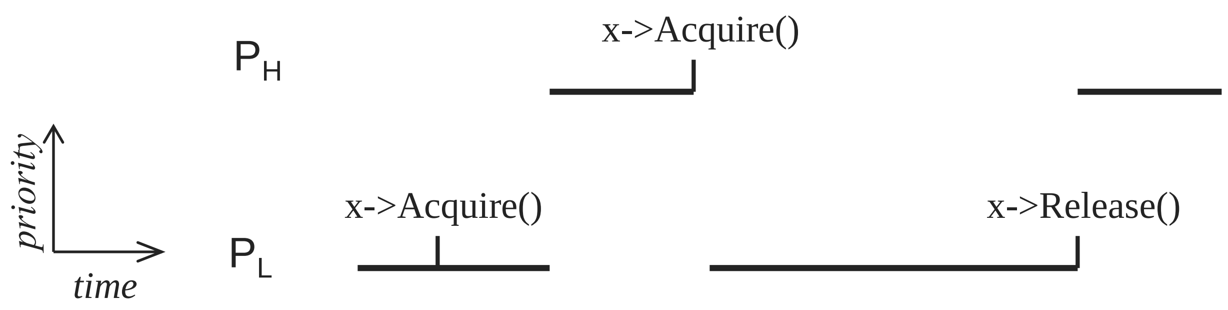
- How do we balance a mixture of repeating tasks:
  - Some I/O bound, need only a little CPU
  - Some compute bound, can use as much CPU as they are assigned
- One approach: maximize the minimum allocation given to a task
  - Schedule the smallest task first, then split the remaining time using max-min

Priority Scheduling

- Priority Scheduling
  - Choose next job based on priority
    - » Airline checkin for first class passengers
  - Can implement SJF, priority = 1/(expected CPU burst)
  - Also can be either preemptive or non-preemptive
- Problem?
  - Starvation – low priority jobs can wait indefinitely
- Solution
  - “Age” processes
    - » Increase priority as a function of waiting time
    - » Decrease priority as a function of CPU consumption

More on Priority Scheduling

- For real-time (predictable) systems, priority is often used to isolate a process from those with lower priority. *Priority inversion* is a risk unless all resources are jointly scheduled.



Priority inheritance

- If lower priority process is being waited on by a higher priority process it inherits its priority
  - How does this help?
  - Does it prevent the previous problem?
- Priority inversion is a big problem for real-time systems
  - Mars pathfinder bug (link)

Problems of basic algorithms

- FIFO: Good: fairness; bad: turnaround time, response time
- SJF: good: turnaround time; bad: fairness, response time, need to estimate run-time
- RR: good: fairness, response time; bad: turnaround time
- Is there a scheduler that balances these issues better?
  - Challenge: limited information about a process in the beginning
  - Challenge: how to prevent gaming the scheduler to get more run-time

MLQ: combining algorithms

- Scheduling algorithms can be combined
  - Have multiple queues
  - Use a different algorithm for each queue
  - Move processes among queues
- Example: Multiple-level feedback queues (MLFQ)
  - Multiple queues representing different job types
    - » Interactive, CPU-bound, batch, system, etc.
  - Queues have priorities, jobs on same queue scheduled RR
  - Jobs can move among queues based upon execution history
    - » Feedback: Switch from interactive to CPU-bound behavior

Multi-level Feedback Queue (MFQ)

- Goals:
  - Responsiveness
  - Low overhead
  - Starvation freedom
  - Some tasks are high/low priority
  - Fairness (among equal priority tasks)
- Not perfect at any of them!
  - Used in Linux (and probably Windows, MacOS)

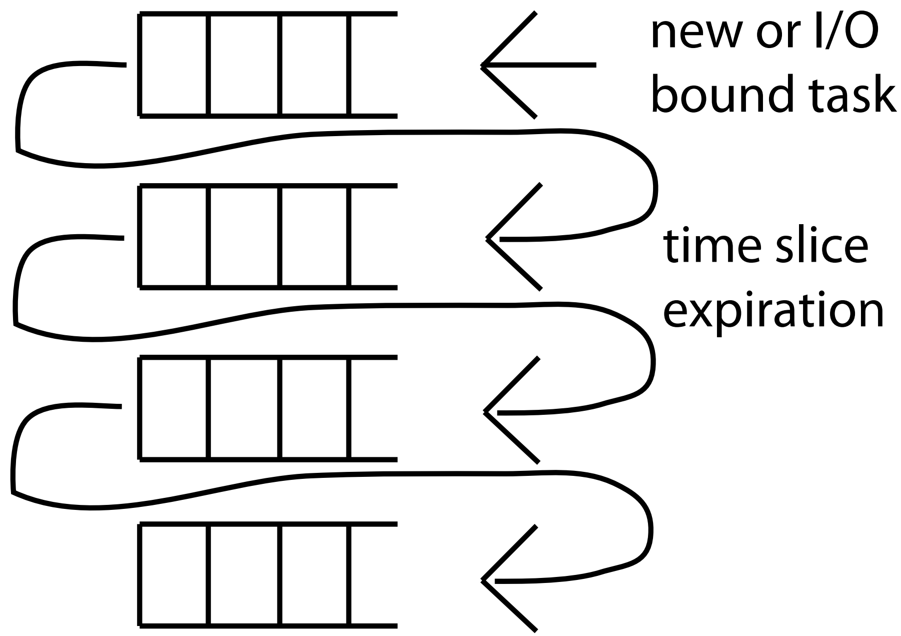
MFQ

- Set of Round Robin queues
  - Each queue has a separate priority
- High priority queues have short time slices
  - Low priority queues have long time slices
- Scheduler picks first thread in highest priority queue
- Tasks start in highest priority queue
  - If time slice expires, task drops one level

MFQ

Priority	Time Slice (ms)
1	10
2	20
3	40
4	80

Round Robin Queues



Unix Scheduler

- The canonical Unix scheduler uses a MLFQ
  - 3-4 classes spanning ~170 priority levels
    - Timesharing: first 60 priorities
    - System: next 40 priorities
    - Real-time: next 60 priorities
    - Interrupt: next 10 (Solaris)
- Priority scheduling across queues, RR within a queue
  - The process with the highest priority always runs
  - Processes with the same priority are scheduled RR
- Processes dynamically change priority
  - Increases over time if process blocks before end of quantum
  - Decreases over time if process uses entire quantum

Motivation of Unix Scheduler

- The idea behind the Unix scheduler is to reward interactive processes over CPU hogs
- Interactive processes (shell, editor, etc.) typically run using short CPU bursts
  - They do not finish quantum before waiting for more input
- Want to minimize response time
  - Time from keystroke (putting process on ready queue) to executing keystroke handler (process running)
  - Don't want editor to wait until CPU hog finishes quantum
- This policy delays execution of CPU-bound jobs
  - But that's ok

Scheduling Summary

- Scheduler (dispatcher) is the module that gets invoked when a context switch needs to happen
- Scheduling algorithm determines which process runs, where processes are placed on queues
- Many potential goals of scheduling algorithms
  - Utilization, throughput, wait time, response time, etc.
- Various algorithms to meet these goals
  - FCFS/FIFO, SJF, Priority, RR
- Can combine algorithms
  - Multiple-level feedback queues
  - Unix example