Synergy: Quality of Service Support for Distributed Stream Processing Systems

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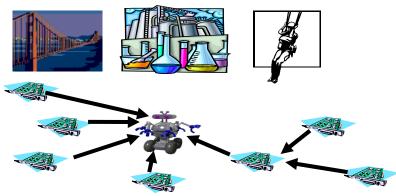


Research Contributions

- Distributed Stream Processing Systems
 - Sharing-Aware Component Composition [Middleware'06, TPDS'08 (rev.)]
 - Load Prediction and Hot-Spot Alleviation [DSN'08, DBISP2P'07]
 - Replica Placement for High Availability [DEBS'08]
- Management of Large-Scale, Distributed, Real-Time Applications
 - Adaptation to Resource Availability [IPDPS'05]
 - Fair Resource Allocation [ISORC'06, WPDRTS'05]
- Peer-to-Peer Systems
 - Adaptive Data Dissemination and Routing [MDM'05]
 - Decentralized Trust Management [MPAC'06]
- Software Distributed Shared Memory Systems
 - Data Migration [Cluster'05, Cluster'04]
- Replication in Distributed Multi-Tier Architectures [IBM'07]
- Collaborative Spam Filtering [Intel'06]
- Distributed Logging for Asynchronous Replication [HP'05]



On-Line Data Stream Processing Network traffic monitoring for intrusion detection



Analysis of readings coming from sensors or mobile robots

Click stream analysis for purchase recommendations or advertisements





Customization of multimedia or news feeds







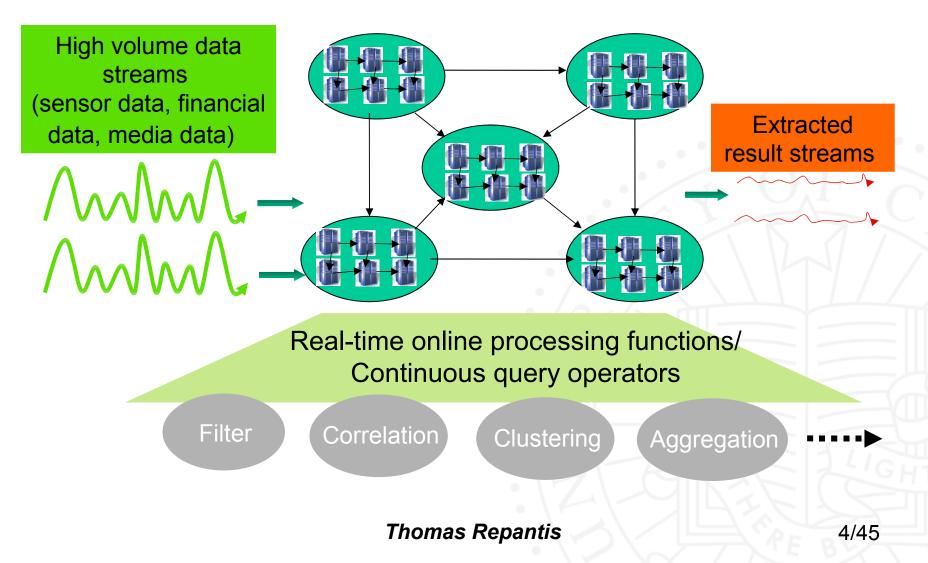






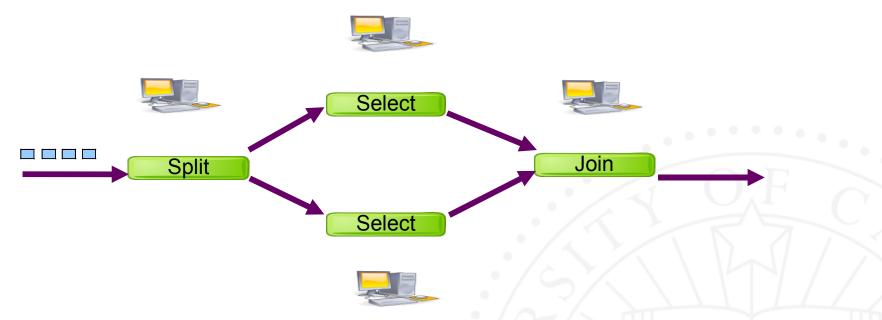


Distributed Stream Processing System





Stream Processing Environment



- Streams are processed online by components distributed across hosts
- Data arrive in large volumes and high rates, while workload spikes are not known in advance
- Stream processing applications have QoS requirements, e.g., e2e delay

QoS for Distributed Stream Processing Applications

- Our goal: How to run stream processing applications with QoS requirements, while efficiently managing system resources
 - Share existing result streams
 - Share existing stream processing components
 - Predict QoS violations
 - Alleviate hot-spots
 - > Maximize availability

Benefits

- Enhanced QoS provision
- Reduced resource load

> Challenges

- Concurrent component sharing
- Highly dynamic environment
- On-demand stream application requests
- Scale that dictates decentralization Thomas Repantis



Roadmap

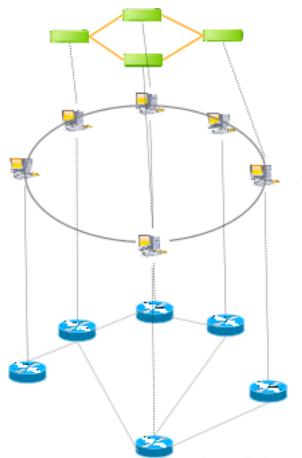
Motivation and Background

- Synergy Architecture
- Design and Algorithms
 - Component Composition
 - Composition Protocol
 - Component and Stream Sharing
 - Load Balancing
 - Hot-Spot Prediction
 - Hot-Spot Alleviation
 - High Availability
 - Replica Placement
- Conclusion
- > Demo



Synergy Middleware

- A middleware managing the mappings:
 - From application layer to stream processing overlay layer
 - From stream processing overlay layer to physical resource layer



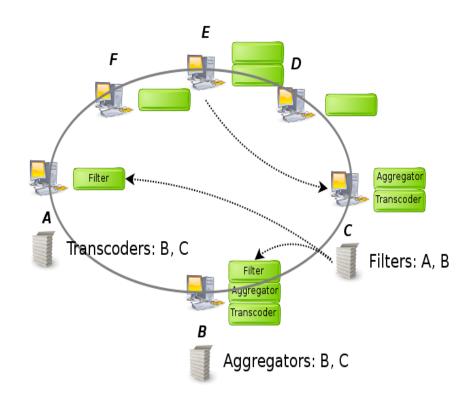
Distributed Stream Processing Application

Synergy Distributed Stream Processing Middleware

IP Network

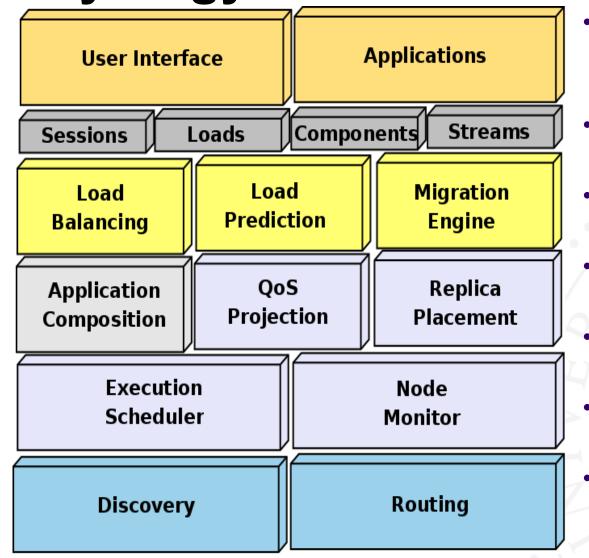


Metadata Layer Over a DHT



- Decouples stream and component placement from their discovery
- Stream and component names are hashed in a DHT
- DHT maps the hashed names to nodes currently offering the specified stream or component

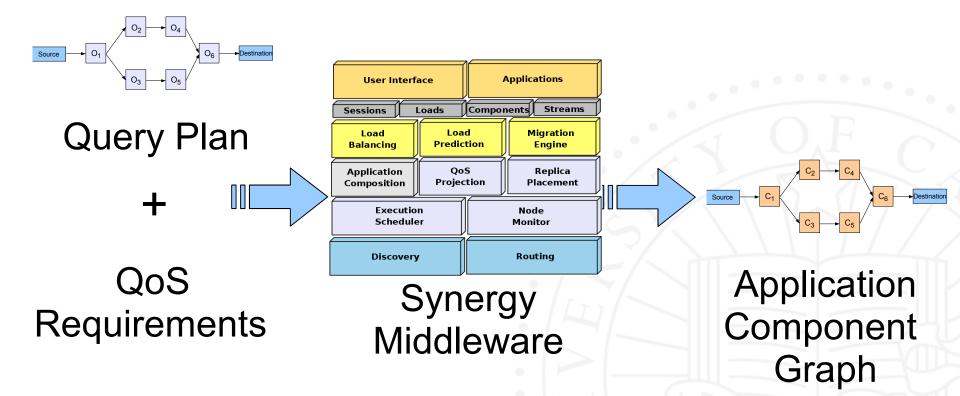
Synergy Node Architecture



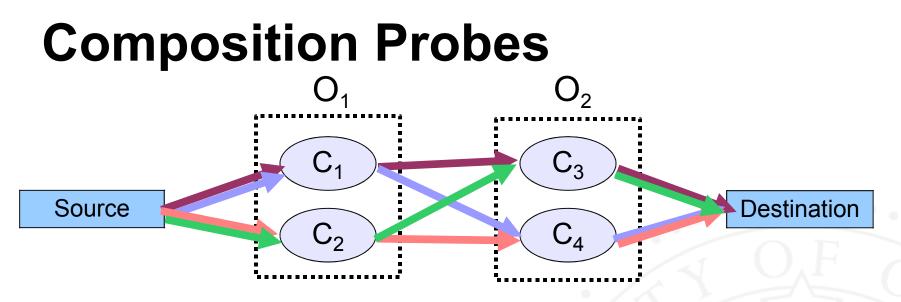
- Application Composition and QoS Projection instantiate applications
- Replica Placement places components
- Load Balancing and Load Prediction detect hot-spots
- Migration Engine alleviates hot-spots
- Monitor measures processor and bandwidth
- Discovery locates streams and components
- Routing transfers streaming data



Component Composition



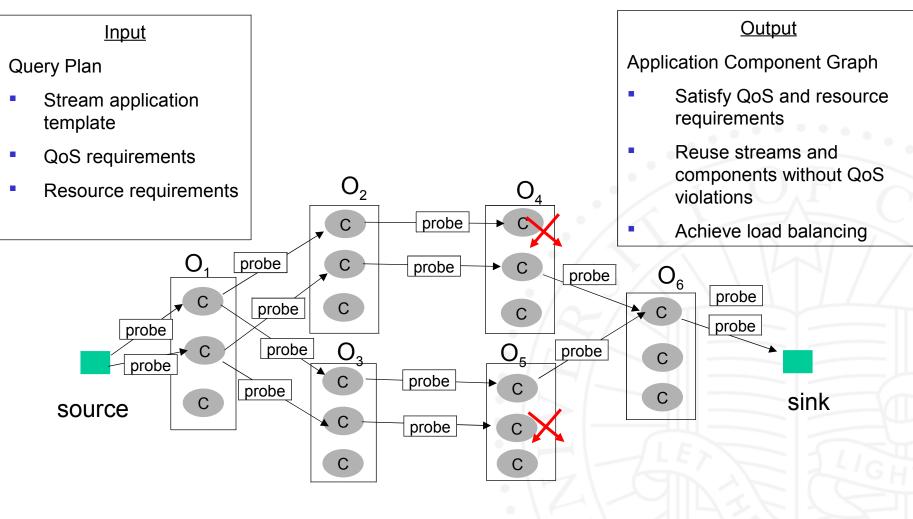




- Carry query plan, resource, and QoS requirements
- Collect information about:
 - Resource availability
 - End-to-end QoS
 - QoS impact on existing applications



Composition Protocol





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Composition Selection

- All successful probes returning to source have been checked against constraints on:
 - Operator functions
 - Processing capacity
 - > Bandwidth
 - > QoS
- The most load balanced one is selected among all qualified compositions by minimizing:

$$\phi(\lambda) = \sum_{v_i \in V_\lambda, o_i \in \xi} \frac{p_{o_i}}{r p_{v_i} + p_{o_i}} + \sum_{l_j \in \lambda, s_j \in \xi} \frac{b_{s_j}}{r b_{l_j} + b_{s_j}}$$



Component Sharing

QoS Impact Projection Algorithm

All existing and the new application should not exceed requested execution time:

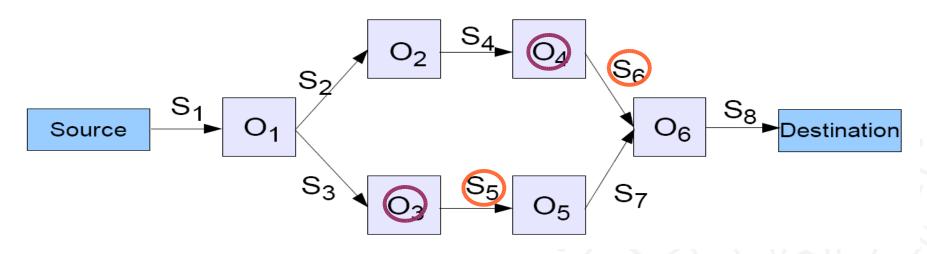
$$\delta + \hat{t} \le q_t$$

Impact estimated using a queueing model for the execution time:

$$\delta = \hat{t}' - \hat{t} = \frac{\tau_{c_i}}{1 - (p_{v_i} + p_{c_i})} - \frac{\tau_{c_i}}{1 - p_{v_i}}$$



Stream Sharing



Maximum Sharing Discovery Algorithm

- Breadth first search on query plan to identify latest possible existing output streams
- Backtracking hop-by-hop, querying the metadata layer

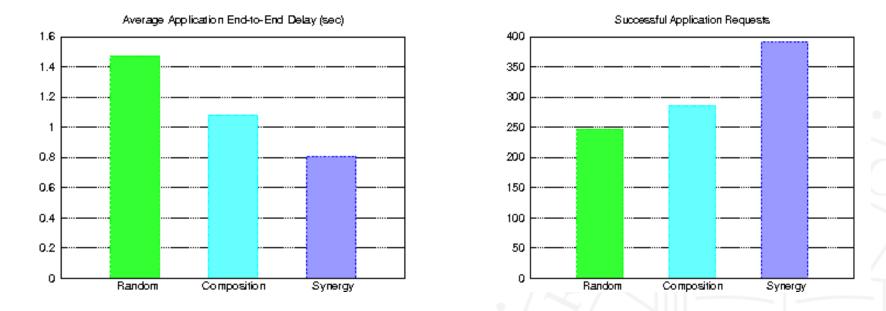


Experimental Setup

- PlanetLab multi-threaded prototype of about 35000 lines of Java running on 88 PlanetLab nodes
- Simulator of about 8500 lines of C++ for 500 random nodes of a GT-ITM topology of 1500 routers
- > 5 replicas of each component
- Synergy vs Random, Greedy, and Composition



Composition Performance

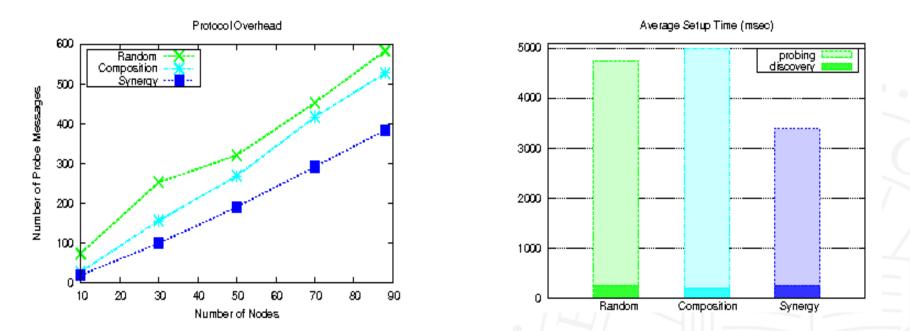


Stream reuse improves end-to-end delay by saving processing time and increases system capacity

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Composition Overhead

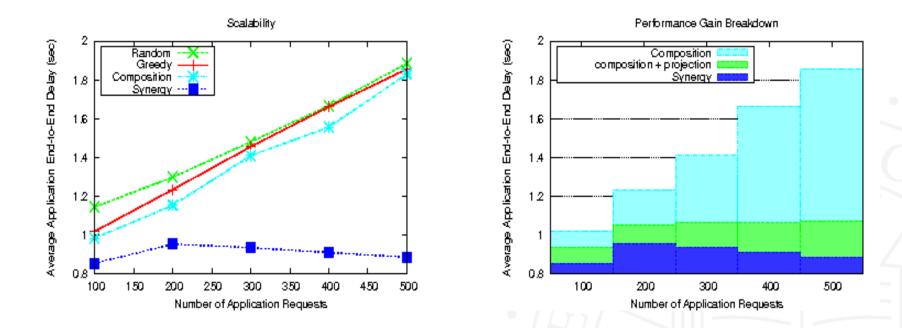


Stream reuse decreases probing overhead and setup time

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Performance on Simulator

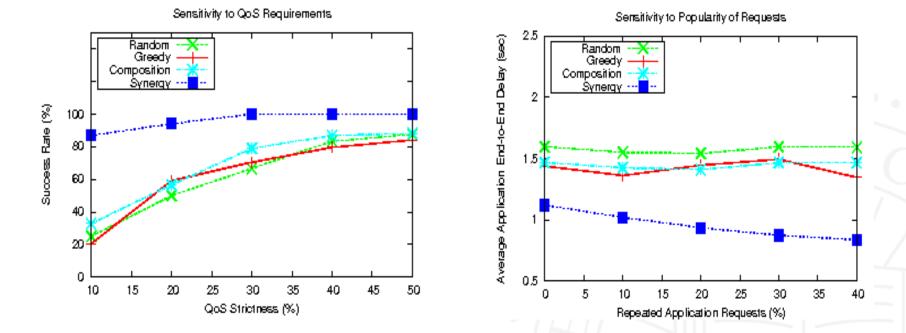


End-to-end delay scales due to stream reuse and QoS impact projection

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Sensitivity on Simulator

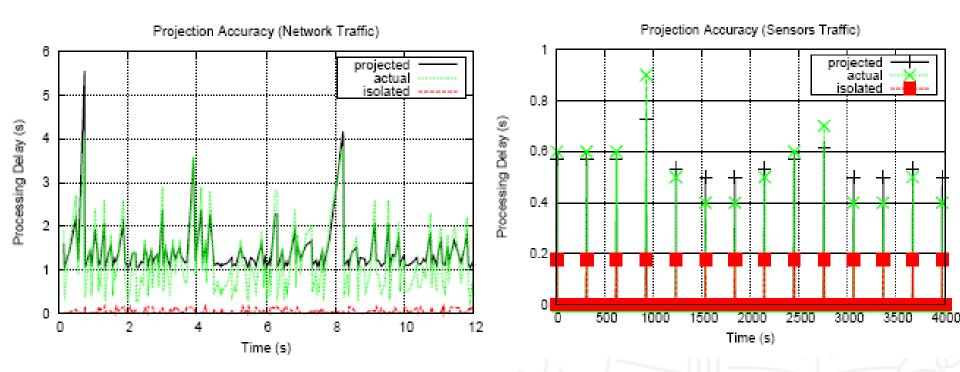


Synergy performs consistently better, regardless of QoS strictness or query popularity

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Projection Accuracy



Pessimistic projections for low rate segments may cause conservative compositions but no QoS violations

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Roadmap

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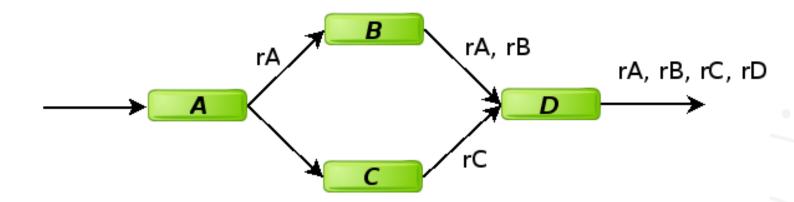


Application-Oriented Load Management

- System hot-spots: Overloaded nodes
- Application hot-spots: QoS violations
 - Sensitive hot-spot detection
 - Triggered even when underloaded, if stringent QoS
 - Fine-grained hot-spot alleviation
 - Only suffering applications migrate
 - Proactively prevent QoS degradation



Predicting QoS Violations



Calculate slack time t_s on every component based on execution time t_e and communication time t_c

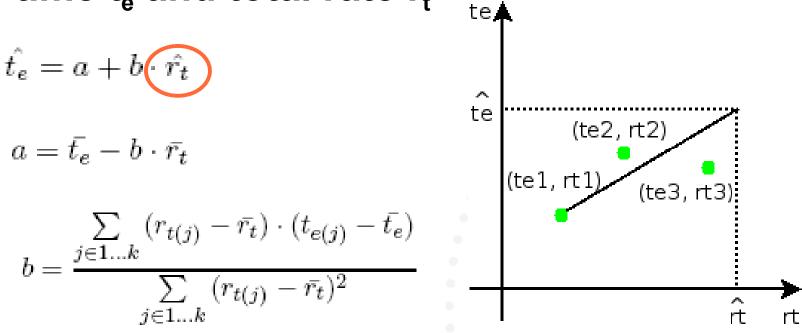
$$t_{s(i)} = q_t - \left(\sum_{j \in 1...i-1} t_{c(j \to j+1)} + \sum_{j \in 1...i-1} t_{e(j)} + \sum_{i...v-1} t_{c(j \to j+1)} + \sum_{j \in i...v} (t_{e(j)}) > 0\right)$$

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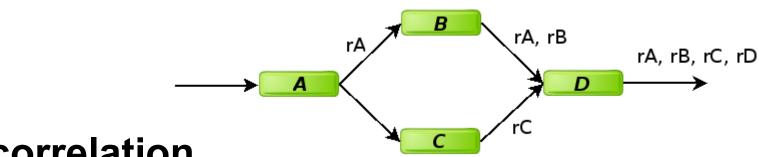
Execution Time Prediction

Linear regression to bind execution time t_e and total rate r_t









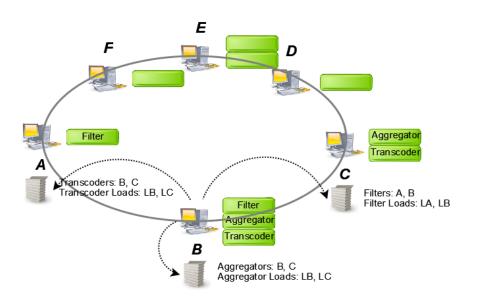
Auto-correlation

$$\hat{r_k} = \frac{\arg_m \max R_{(k)}}{\arg_m \max R_{(k-1)}} \cdot r_{k-1} = \frac{r_{k(m)}}{r_{k-1(m)}} \cdot r_{k-1}$$

Cross-correlation (Pearson Product Moment)

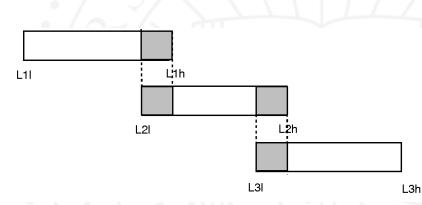
$$R_{i} = \frac{\sum_{j \in 1...(k-1)} (r_{j(i)} - \bar{r_{(i)}})(r_{j} - \bar{r})}{\sqrt{\sum_{j \in 1...(k-1)} (r_{j(i)} - \bar{r_{(i)}})^{2} \sum_{j \in 1...(k-1)} (r_{j} - \bar{r})^{2}}} (r_{j(i)} - \bar{r_{(i)}})^{2} \sum_{j \in 1...(k-1)} (r_{j} - \bar{r})^{2}} (r_{j(i)} - \bar{r_{(i)}})^{2} \sum_{j \in 1...(k-1)} (r_{j(i)} - \bar{r_{(i)}})^{2}} (r_{j(i)} - \bar{r_{(i)}})^{2} (r_{j(i)} - \bar{r_{(i)}})^{2}} (r_{j(i)} - \bar{r_{(i)}})^{2}} (r_{j(i)} - \bar{r_{(i)}})^{2}} (r_{j(i)} - \bar{r_{(i)}})^{2} (r_{j(i)} - \bar{r_{(i)}})^{2} (r_{j(i)} - \bar{r_{(i)}})^{2}} (r_{j(i)} - \bar{r_{(i)}})^{2} (r_{j(i)} - \bar{r_{(i)}})^{2}} (r_{j(i)} - \bar{r_{(i)}})^{2} (r_{j(i)} - \bar{r_{(i)}})^{2}} (r_{j(i)} - \bar{r_{(i)}})^{2}} (r_{j(i)} - \bar{r_{(i)}})^{2} (r_{j(i)} - \bar{r_{(i)}})^{2}} (r_{j(i)} - \bar{r_{(i)}})^{2} (r_{j(i)} - \bar{r_{(i)}})^{2} (r_{j(i)} - \bar{r_{(i)}})^{2}} (r_{j(i)} - \bar{r_{(i)}})^{2} (r_{j(i)} - \bar{r_{(i)}})^{$$

Decentralized Load Monitoring



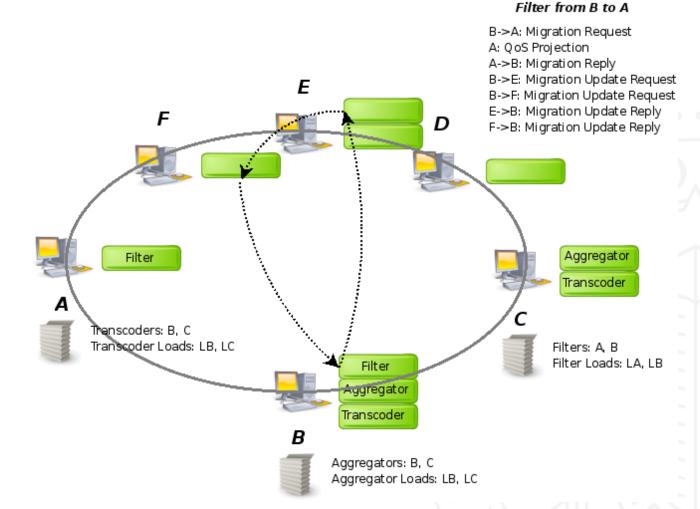
- Load updates pushed when intervals change
- Overlapping intervals absorb frequent changes

- DHT maps component names to the loads of peers hosting them
- Peers detect overloads and imbalances between all hosts of a component





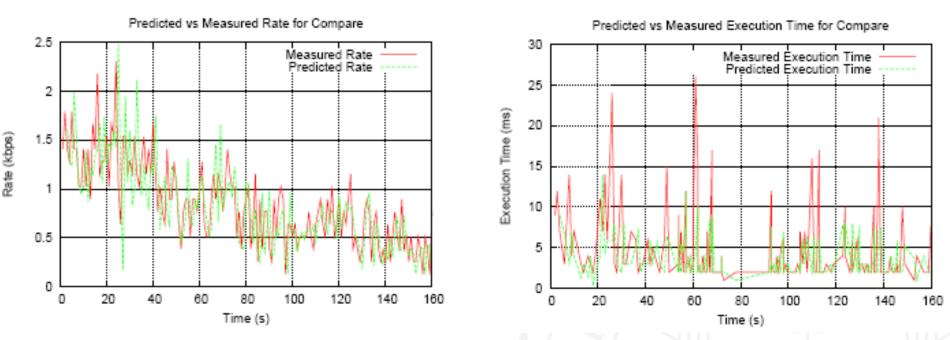
Alleviating Hot-Spots via Migration



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Hot-Spot Prediction and Alleviation

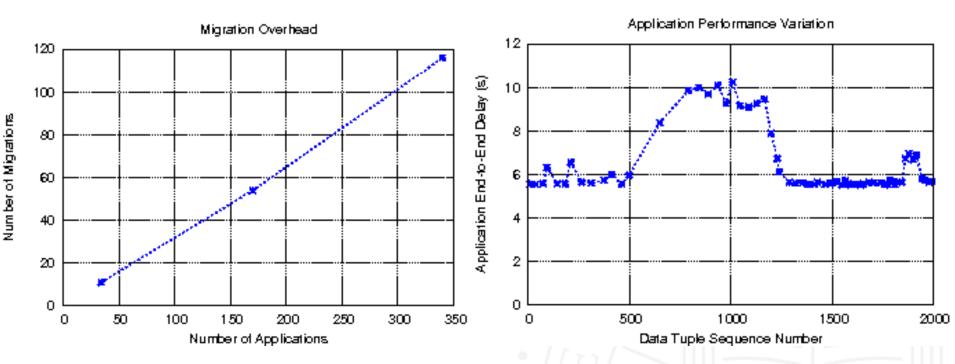


Average prediction error 3.7016% Average prediction overhead 0.5984ms

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Hot-Spot Prediction and Alleviation

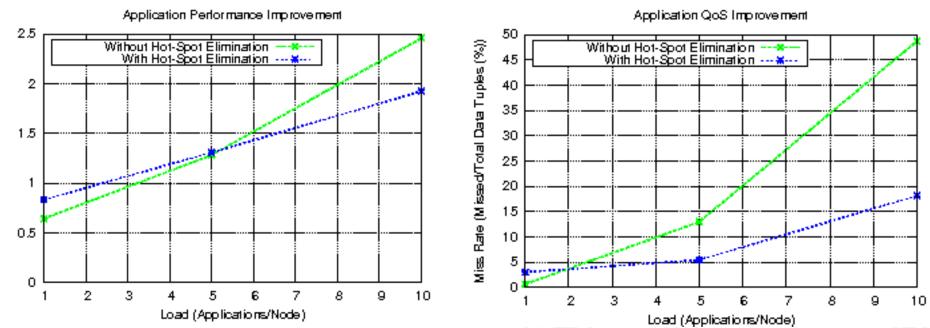


Average one migration every three applications Average migration time 1144ms

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QoS Improvement



As load increases the benefits of hot-spot elimination become evident

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Average Application End-to-End Delay (s)



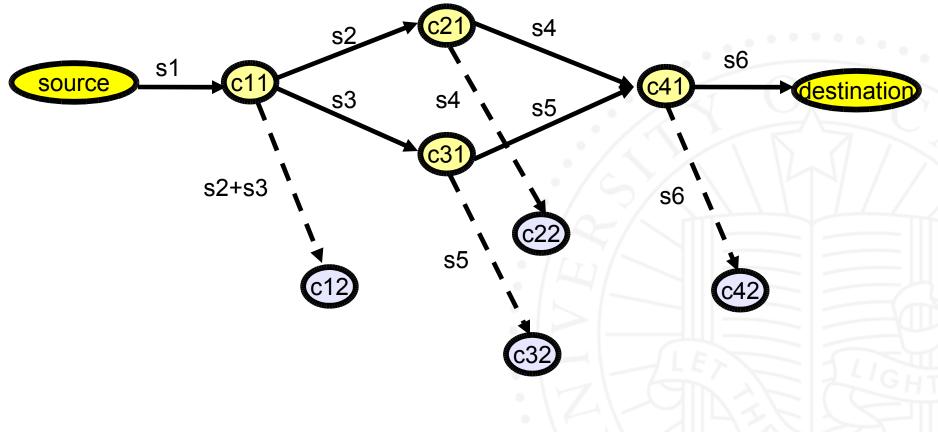
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Component Replication



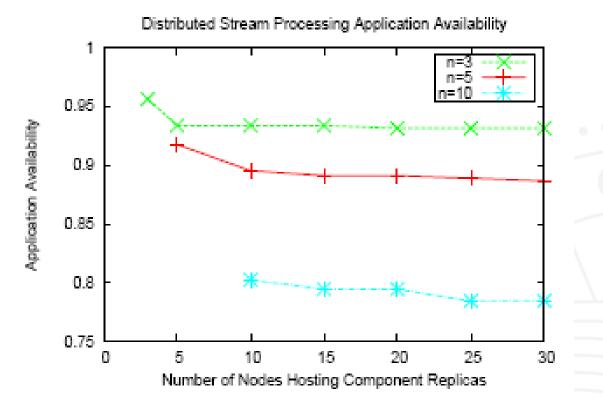


Component Replica Placement

- Maximize availability of composite applications
 - Optimal: Place complete graph on each node
- Respect node resource availability
 - Processing capacity
 - Network bandwidth
- Maximize application performance
 - Inter-operator communication cost (between primaries)
 - Intra-operator communication cost (between primaries and backups)



Placement for High Availability

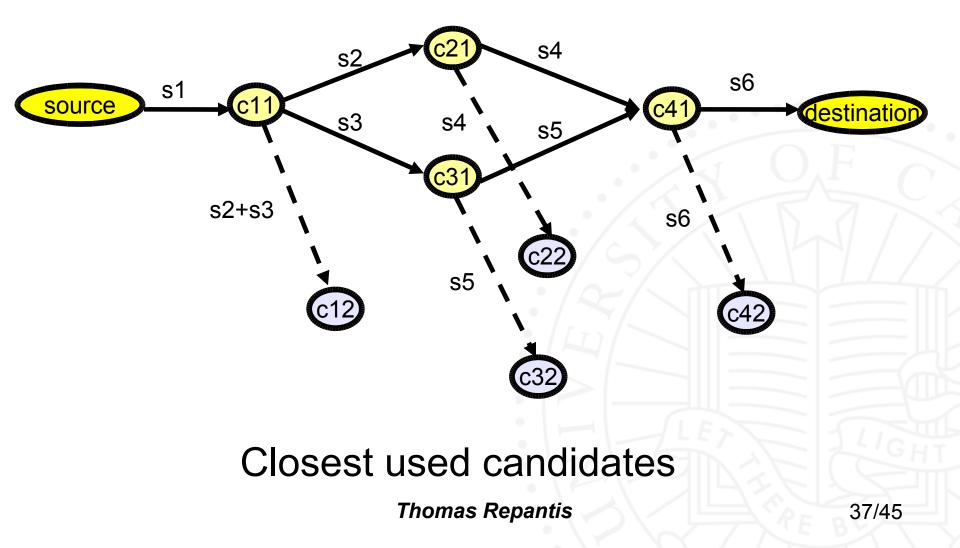


Availability decreases with larger graphs and increases with higher concentration

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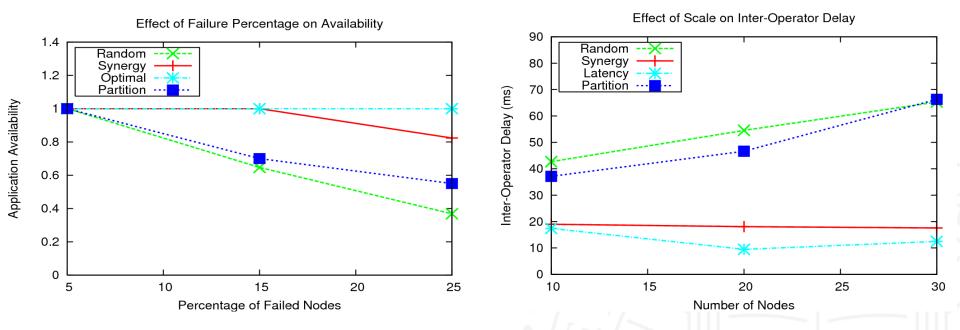


Distributed Placement Protocol





Replica Placement



Increase availability and performance 5539ms to gather latencies for 30 nodes

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Related Work

- System S: IBM stream processing middleware
- SBON, SAND, IFLOW: Component placement
- Borealis, Flux, PeerCQ: Load balancing
- Borealis, TelegraphCQ: Load shedding
- Borealis, Flux: Fault tolerance
- SpiderNet, sFlow: Component composition



Conclusion

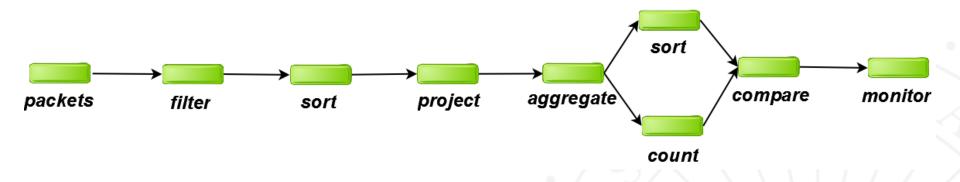
Synergy: QoS-Enabled Distributed Stream Processing System

- Component Composition
 - Fully distributed composition protocol
 - Reuse existing streams and components
- Load Balancing
 - Predict QoS violations
 - Alleviate hot-spots using migration
- High Availability
 - Place component replicas
- Future work
 - Efficient and consistent replication
 - Adaptive topology management
 - Secure composite applications



Demo

Monitor source-destination pairs in top 5% of total traffic over last 20 minutes [Stream Query Repository]



TCP traffic trace, LBL, 2 hours, 1.8 million packets [Internet Traffic Archive]

timestamp	sourceIP	destinationIP	sourcePort	destinationPort	size
	:				15/61

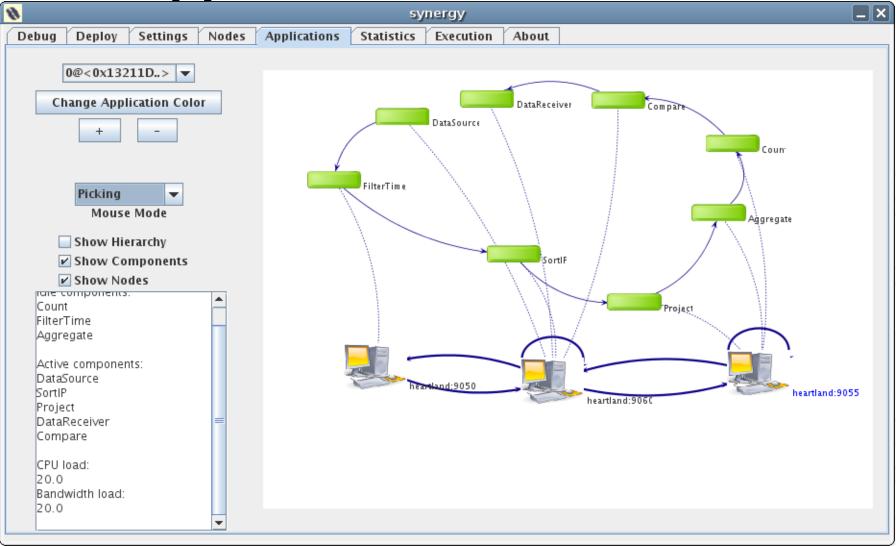


GUI Settings

synergy							
Debug Deploy Settings Nodes Applications Statistics Execution About							
Initialization Connection Node Service Application Algorithms Load Balancing Fault To	lerance Replication						
The application to run.	netmon 💌						
Number of services in an application.	3						
Maximum requested rate from the source to the destination of the application in kbps.	50						
Number of data units to send before we exit.	1000000						
The size of the payload of every data unit.	250						
	Restore Save						
Thomas Repantis	42/45						



GUI Application



GUI Execution

1						synergy			
Debug	Deploy	Settin	gs Nodes	Applications	Statistic	s Execution	About		
		[
Applica	ation List		1.348691 44	29 1032	1.627910 4	23 1024	1.933559 4 53 1024	1.313309 3 4 135	1.3
0@<0x	7039AC>	_	2.945475 54	\$ 55 1136	2.492097 4	23 1024	2.579007 49 48 945	2.710139 2 5 664	3.1
DC 10A			3.313864 4	23 1066	4.200897 2	9 58 399	3.329686 23 4 318	4.052514 16 17 273	4.1
			5.234655 4	23 1818	4.499406 5	54 55 784	5.273786 41 40 512	4.362552 1 2 369	5.1
			5.707645 40		6.317355 2		5.372508 54 55 616	5.514410 3 4 419	6.2 💻
			6.912369 54	\$ 55 1968	6.363782 2	20 19 552	6.790638 58 29 525	6.365901 29 63 515	7.1
			7.63305140) 41 1024	7.6680012	29 58 793	7.518405 4 23 556	7.953093 3 4 249	8.0
			8.861617 40		8.372301 4		9.130298 44 29 1024	8.455924 16 17 828	9.1
			9.503625 40		9.999841 4		9.866365 44 4 1013	10.141051 48 49 685	9.5
			10.783522 2			49 48 2048	11.140872 29 30 1080	10.457738 14 15 952	11.
			11.843829 4		11.480383		12.147954 3 4 404	12.047270 4 3 75	12.:
			13.074514 4			40 41 1051	13.107023 3 4 165	12.479074 10 11 160	13.(
			14.175230 5		13.875018		14.296178 3 4 294	13.966418 11 10 100	14.:
			15.232693 4		15.417207		14.861595 10 11 128	15.178908 31 13 103	14.
			16.314291 4		15.923578		16.171491 44 29 1583	15.707918 40 41 1536	16.
			17.105083 4			29 30 1402	17.178927 29 58 984	17.031698 4 23 949	17.:
			18.462050 4			44 29 2219	18.450281 4 53 2083	17.874455 40 41 2048	18.
			18.7160512		18.950916		19.556264 4 53 2320	19.433072 4 23 1024	19.
			20.350927 4			40 41 1536	19.841882 4 53 1484	20.107164 4 52 1156	20.
			20.6475184		21.039337	· · ·	21.195188 40 41 3160	21.238316 29 30 2684	21.
			22.177906 4			29 58 1050	21.614500 40 41 1024	22.540584 4 53 396	21.
			23.372508 2		22.992167		23.298510 4 53 754	23.305994 44 29 737	23.1
						40 41 2560	24.508596 44 4 1541	23.635492 95 94 768	24.
						40 41 4096	25.151524 44 4 2093	25.148330 95 94 1408	25.
			26.034145 9			40 41 6144	25.671303 4 53 2833	26.131268 44 29 2605	25.0
			27.631830 4		27.388336		26.818782 40 41 2560	27.080234 44 29 2121	27.
			27 646943 4	10 41 4100	27 769050	29 58 2598	28 195778 29 30 2323	28 594887 4 52 2079	28 🔍



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- > Bilson Campana, UC Riverside



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