

# Advanced Systems Security: Security-Enhanced Linux

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#### Reference Monitor for Linux



- LSM provides a reference monitor interface for Linux
  - Complete Mediation
- You need a module and infrastructure to achieve the other two goals
  - Tamperproofing
  - Verifiability
- SELinux is a comprehensive reference validation mechanism aiming at reference monitor guarantees

## **SELinux History**



- Origins go back to the Mach microkernel retrofitting projects of the 1980s
  - ▶ DTMach (1992)
  - DTOS (USENIX Security 1995)
  - Flask (USENIX Security 1999)
  - **▶** SELinux (2000-...)
- Motivated by the security kernel design philosophy
  - But, practical considerations were made

## Inevitability of Failure



- Philosophy of the approach
- Flawed Assumption:
  - That security can be managed by the application space without OS security support (protection is not sufficient)
- Paraphrase: Can't build a secure system without a reference monitor and MPS
  - And a secure operating system needs an entire ecosystem
- Come back to this later...

## The Rest of the SELinux Story



- Tamperproof
  - Protect the kernel
  - Protect the trusted computing base
  - Use MPS to provide tamperproofing of TCB?
- Verifiability
  - Code correctness
  - Policy satisfy a security goal
  - Use MPS to express secrecy and integrity requirements

## Design MPS



- Do not believe that classical integrity is achievable in practice
  - Too many exceptions
  - Commercial systems will not accept constraints of classical integrity
- Instead, focus on providing comprehensive control of access aiming for
  - Confining root processes (tamperproof)
  - Least privilege in general (verifiability)
- How does 'least privilege' affect security?

## SELinux Policy Model



• See slides in 07-TypeEnforcement...

## SELinux Policy Rules



- SELinux Rules express an MPS
  - Protection state ALLOW subject-label object-label ops
  - Labeling state TYPE\_TRANSITION subject-label objectlabel new-label (at create — objects)
    - Default is to label to same state as creator
  - Transition state TYPE\_TRANSITION subject-label objectlabel new-label (at exec – processes)
- Tens of thousands of rules are necessary for a standard Linux distribution
  - Protect system processes from user processes
  - User data can be protected by MLS

## SELinux "Setuid"



• How does SELinux enable a normal user to run a privileged (setuid) process, such as passwd?

#### **SELinux Transition State**



- For user to run passwd program
  - Only passwd should have permission to modify /etc/shadow
- Need permission to execute the passwd program
  - allow user\_t passwd\_exec\_t:file execute (user can exec /usr/bin/passwd)
  - allow user\_t passwd\_t:process transition (user gets passwd perms)
- Must transition to passwd\_t from user\_t
  - allow passwd\_t passwd\_exec\_t:file entrypoint (run w/ passwd perms)
  - type\_transition user\_t passwd\_exec\_t:process passwd\_t
- Passwd can the perform the operation
  - allow passwd\_t shadow\_t:file {read write} (can edit passwd file)

## SELinux Deployment



- You've configured your SELinux policy
  - Now what is left?
- Surprisingly, a lot
  - Many services must be aware of SELinux
  - Got to get the policy installed in the kernel
  - Got to manage all this policy
- And then there is the question of getting the policy to do what you want



 What kind of security decisions are made by user-space services?





- What kind of security decisions are made by user-space services?
  - Authentication (e.g., sshd)
  - Access control (e.g., X windows, DBs (servers), browsers (middleware), etc.)
  - Configuration (e.g., policy build and installation)
- Also, many services need to be aware of SELinux to enable usability
  - E.g., Listing files/processes with SELinux contexts (ls/ps)

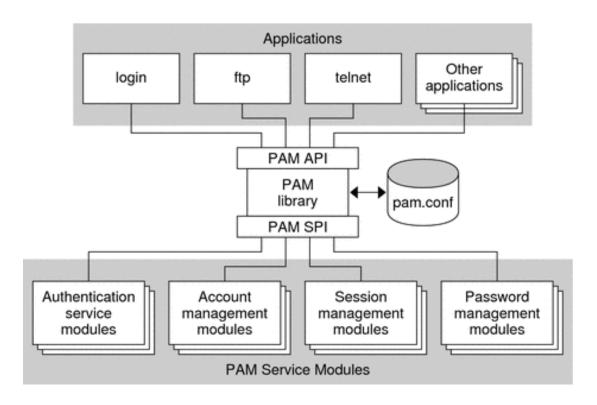


- Authentication
  - Various authentication services need to create a "SELinux subject context" on a user login
  - Like login in general, except we set an SELinux context and a UID for the generated shell
- How do you get all these ad hoc authentication services to interact with SELinux?

#### Authentication for SELinux



- Pluggable Authentication Modules
  - There is a module for SELinux that various authentication services use to create a subject context

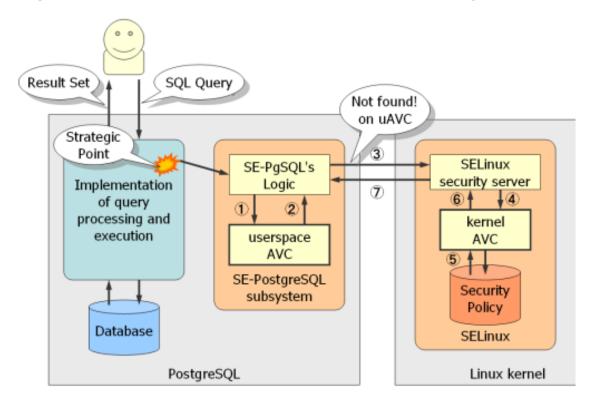




- Access Control
  - Many user-space services are shared among mutually untrusting clients
    - Problem: service may leak one client's secret to another
- If your SELinux policy allows multiple, mutually untrusting clients to talk to the same service, what can SELinux do to prevent exploits?



- Add SELinux support to the service
  - X Windows, postgres, dbus, gconf, telephony server
- E.g., Postgres with the SELinux user-space library



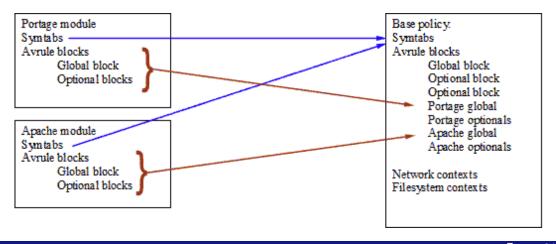


- Configuration
  - You need to get the SELinux policy constructed and loaded into the kernel
    - Without allowing attacker to control the system policy
    - And policy can change dynamically
- How to compose policies?
- How to install policies?

## Compose Policies



- The SELinux policy is modular
  - Although not in a pure, object-oriented sense
    - Too much had been done
- Policy management system composes the policy from modules, linking a module to previous definitions and loads them



## Installing Policies



 How would you enable user-space processes to push data (e.g., MPS configuration) into the kernel?

# sysfs Background



- During the 2.5 development cycle, the Linux driver model was introduced to fix several shortcomings of the 2.4 kernel:
  - No unified method of representing driver-device relationships existed.
  - There was no generic hotplug mechanism.
  - procfs was cluttered with lots of non-process information.
- Main uses
  - Configure drivers
  - Export driver information

# sysfs Example: load\_policy



From kernel: security/selinux/selinuxfs.c

```
enum sel_inos {
   SEL ROOT INO = 2,
   SEL_LOAD, /* load policy */
   SEL ENFORCE, /* get or set enforcing status */
static struct tree descr selinux files[] = {
   [SEL LOAD] = {"load", &sel load ops, S IRUSR|S IWUSR},
   [SEL ENFORCE] = {"enforce", &sel_enforce_ops,
                      S IRUGO|S IWUSR},
static struct file_operations sel_load_ops = {
     .write = sel write load,
```

## sysfs Example: load\_policy



```
From userspace: libselinux/src/load_policy.c
int security load policy(void *data, size t len)
     char path[PATH_MAX];
     int fd, ret;
     snprintf(path, size of path, "%s/load", selinux mnt);
     fd = open(path, O RDWR);
     if (fd < 0)
           return - 1;
     ret = write(fd, data, len);
     close(fd);
```

## sysfs Example: load\_policy



From kernel: security/selinux/selinuxfs.c

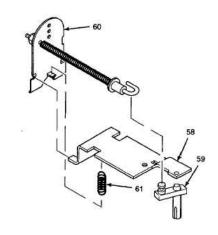
```
static ssize_t sel_write_load(struct file * file, const char user * buf,
                    size t count, loff t *ppos)
length = task_has_security(current, SECURITY LOAD POLICY);
     if (length)
           goto out;
if (copy_from_user(data, buf, count) != 0)
     goto out;
length = security_load_policy(data, count); --- ss/services.c
if (length)
     goto out;
```

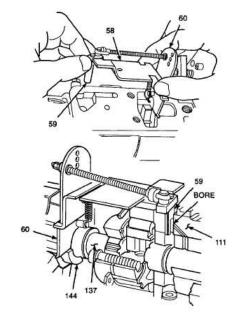
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## When Are We Done?



- There is a significant configuration effort to get the SELinux system deployed
  - Who does this?
  - What happens if I want to change something?
  - Does it prevent the major threats?





## Take Away



- Problem: Turn the SELinux policy into a working, usable reference monitor
  - Work with user-space services
  - Design the policy that you want
- There are many requirements for user-space services to provide authentication, access control, and policy configuration itself
  - PAM, Policy Mgmt, User-space access, Network support
- Design of MPS can only be semi-automated
  - Prevent network threats and design for app integrity