CSE 153
Design of Operating Systems
Fall 2017
Lecture 4: Threads
Announcements

- Homework 1 will be assigned later today.
- Lab 1 is available
  - Read scheduling in the textbook
    » Don’t wait for it to be covered in class

- All set with project groups?
  - Email your TA today if you have not notified your group or if you are looking for a partner
Processes

● Recall that …
  ◆ A process includes:
    » An address space (defining all the code and data pages)
    » OS resources (e.g., open files) and accounting info
    » Execution state (PC, SP, regs, etc.)
    » PCB to keep track of everything
  ◆ Processes are completely isolated from each other

● But…
Some issues with processes

- Creating a new process is costly because of new address space and data structures that must be allocated and initialized
  - Recall struct proc in xv6 or Solaris

- Communicating between processes is costly because most communication goes through the OS
  - Inter Process Communication (IPC) – we will discuss later
  - Overhead of system calls and copying data
Parallel Programs

- Also recall our Web server example that forks off copies of itself to handle multiple simultaneous requests

- To execute these programs we need to
  - Create several processes that execute in parallel
  - Cause each to map to the same address space to share data
    - They are all part of the same computation
  - Have the OS schedule these processes in parallel

- This situation is very inefficient (CoW helps)
  - **Space**: PCB, page tables, etc.
  - **Time**: create data structures, fork and copy addr space, etc.
Rethinking Processes

● What is similar in these cooperating processes?
  ◆ They all share the same code and data (address space)
  ◆ They all share the same privileges
  ◆ They all share the same resources (files, sockets, etc.)

● What don’t they share?
  ◆ Each has its own execution state: PC, SP, and registers

● Key idea: Separate resources from execution state
● Exec state also called thread of control, or thread
Recap: Process Components

- A process is named using its process ID (PID)
- A process contains all of the state for a program in execution

### Per-Process State
- An address space
- The code for the executing program
- The data for the executing program
- A set of operating system resources
  - Open files, network connections, etc.

### Per-Thread State
- An execution stack encapsulating the state of procedure calls
- The program counter (PC) indicating the next instruction
- A set of general-purpose registers with current values
- Current execution state (Ready/Running/Waiting)
Threads

- Separate execution and resource container roles
  - The **thread** defines a sequential execution stream within a process (PC, SP, registers)
  - The **process** defines the address space, resources, and general process attributes (everything but threads)

- Threads become the unit of scheduling
  - Processes are now the **containers** in which threads execute
  - Processes become static, threads are the dynamic entities
Recap: Process Address Space

- Stack
- Heap (Dynamic Memory Alloc)
- Static Data (Data Segment)
- Code (Text Segment)
Threads in a Process

- Stack (T1)
- Stack (T2)
- Stack (T3)
- Heap
- Static Data
- Code

Thread 1
- PC (T1)

Thread 2
- PC (T2)

Thread 3
- PC (T3)
Thread Design Space

<table>
<thead>
<tr>
<th>Address Space</th>
<th>Thread</th>
<th>One Thread/Process</th>
<th>Many Threads/Process</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td>One Address Space</td>
<td>One Address Space</td>
</tr>
<tr>
<td></td>
<td></td>
<td>(MSDOS)</td>
<td>(Pilot, Java)</td>
</tr>
<tr>
<td></td>
<td></td>
<td>(Early Unix)</td>
<td>(Mac OS, Unix, Windows)</td>
</tr>
</tbody>
</table>
Separating threads and processes makes it easier to support multithreaded applications
- Concurrency does not require creating new processes

Concurrency (multithreading) can be very useful
- Improving program structure
- Handling concurrent events (e.g., Web requests)
- Writing parallel programs

So multithreading is even useful on a uniprocessor
Threads: Concurrent Servers

- Using fork() to create new processes to handle requests in parallel is overkill for such a simple task
- Recall our forking Web server:

```c
while (1) {
    int sock = accept();
    if ((child_pid = fork()) == 0) {
        Handle client request
        Close socket and exit
    } else {
        Close socket
    }
}
```
Threads: Concurrent Servers

- Instead, we can create a new thread for each request

```c
web_server() {
    while (1) {
        int sock = accept();
        thread_fork(handle_request, sock);
    }
}

handle_request(int sock) {
    Process request
    close(sock);
}
```
Implementing threads

- One option: OS now manages threads *and* processes
  - All thread operations are implemented in the kernel
  - The OS schedules all of the threads in the system

- OS-managed threads are called *kernel-level threads* or *lightweight processes*
  - Windows: threads
  - Solaris: lightweight processes (LWP)
  - POSIX Threads (pthreads): PTHREAD_SCOPE_SYSTEM
Kernel Thread (KLT) Limitations

- KLTs make concurrency cheaper than processes
  - Much less state to allocate and initialize

- However, there are a couple of issues
  - Issue 1: KLT overhead still high
    - Thread operations still require system calls
    - Ideally, want thread operations to be as fast as a procedure call
  - Issue 2: KLTs are general; unaware of application needs

- Alternative: User-level threads (ULT)
User-Level Threads

- Implement threads using user-level library

- ULTs are small and fast
  - A thread is simply represented by a PC, registers, stack, and small thread control block (TCB)
  - Creating a new thread, switching between threads, and synchronizing threads are done via procedure call
    » No kernel involvement
  - User-level thread operations 100x faster than kernel threads
  - pthreads: PTHREAD_SCOPE_PROCESS
User and Kernel Threads

Multiplexing user-level threads on a single kernel thread for each process

Multiplexing user-level threads on multiple kernel threads for each process
ULT Limitations

- But, user-level threads are not a perfect solution
  - As with everything else, they are a tradeoff
- ULTs are invisible to the OS

- As a result, the OS can make poor decisions
  - Scheduling a process with idle threads
  - Blocking a process whose thread initiated an I/O, even though the process has other threads that can execute
  - Unscheduling a process with a thread holding a lock

- Solving this requires communication between the kernel and the user-level thread manager
Summary KLT vs. ULT

- **Kernel-level threads**
  - Integrated with OS (informed scheduling)
  - Slow to create, manipulate, synchronize

- **User-level threads**
  - Fast to create, manipulate, synchronize
  - Not integrated with OS (uninformed scheduling)

- Understanding the differences between kernel and user-level threads is important
  - For programming (correctness, performance)
  - For test-taking 😊
Kernel and User Threads

- Or use both kernel and user-level threads
  - Can associate a user-level thread with a kernel-level thread
  - Or, multiplex user-level threads on top of kernel-level threads

- Java Virtual Machine (JVM) (also pthreads)
  - Java threads are user-level threads
  - On older Unix, only one “kernel thread” per process
    - Multiplex all Java threads on this one kernel thread
  - On NT, modern Unix
    - Can multiplex Java threads on multiple kernel threads
    - Can have more Java threads than kernel threads
    - Why?
Implementing Threads

- Implementing threads has a number of issues
  - Interface
  - Context switch
  - Preemptive vs. non-preemptive
  - Scheduling
  - Synchronization (next lecture)

- Focus on ULT
  - Kernel-level threads are similar to original process management and implementation in the OS
  - What you will be dealing with in xv6
  - Not only will you be using threads in xv6, you will be implementing more thread functionality
Sample Thread Interface

- **thread_fork(procedure_t)**
  - Create a new thread of control
  - Also thread_create(), thread_setstate()

- **thread_stop()**
  - Stop the calling thread; also thread_block

- **thread_start(thread_t)**
  - Start the given thread

- **thread_yield()**
  - Voluntarily give up the processor

- **thread_exit()**
  - Terminate the calling thread; also thread_destroy
Thread Scheduling

- The thread scheduler determines when a thread runs
- It uses queues to keep track of what threads are doing
  - Just like the OS and processes
  - But it is implemented at user-level in a library
- Run queue: Threads currently running (usually one)
- Ready queue: Threads ready to run
- Are there wait queues?
  - How would you implement thread_sleep(time)?
Non-Preemptive Scheduling

- Threads voluntarily give up the CPU with `thread_yield`

### Ping Thread
```c
while (1) {
    printf("ping\n");
    thread_yield();
}
```

### Pong Thread
```c
while (1) {
    printf("pong\n");
    thread_yield();
}
```

- What is the output of running these two threads?
The semantics of `thread_yield` are that it gives up the CPU to another thread.

- In other words, it *context switches* to another thread.

So what does it mean for `thread_yield` to return?

**Execution trace of ping/pong**

- `printf("ping\n");`
- `thread_yield();`
- `printf("pong\n");`
- `thread_yield();`
- `...`
Implementing `thread_yield()`

```c
thread_yield() {
    thread_t old_thread = current_thread;
    current_thread = get_next_thread();
    append_to_queue(ready_queue, old_thread);
    context_switch(old_thread, current_thread);
    return;
}
```

- The magic step is invoking `context_switch()`
- **Why do we need to call `append_to_queue()`?**
The context switch routine does all of the magic

- Saves context of the currently running thread (old_thread)
  - Push all machine state onto its stack (*not* its TCB)
- Restores context of the next thread
  - Pop all machine state from the next thread’s stack
- The next thread becomes the current thread
- Return to caller as new thread

This is all done in assembly language

- It works at the level of the procedure calling convention, so it cannot be implemented using procedure calls
Preemptive Scheduling

- Non-preemptive threads have to voluntarily give up CPU
  - A long-running thread will take over the machine
  - Only voluntary calls to thread_yield(), thread_stop(), or thread_exit() causes a context switch

- Preemptive scheduling causes an involuntary context switch
  - Need to regain control of processor asynchronously
  - Use timer interrupt (How do you do this?)
  - Timer interrupt handler forces current thread to “call” thread_yield
Threads Summary

- Processes are too heavyweight for multiprocessing
  - Time and space overhead
- Solution is to separate threads from processes
  - Kernel-level threads much better, but still significant overhead
  - User-level threads even better, but not well integrated with OS
- Scheduling of threads can be either preemptive or non-preemptive
- Now, how do we get our threads to correctly cooperate with each other?
  - Synchronization…