

**EXAMPLE 7** In Section 4.1, we will show that  $n < 2^n$  whenever  $n$  is a positive integer. Show that this inequality implies that  $n$  is  $O(2^n)$ , and use this inequality to show that  $\log n$  is  $O(n)$ .

*Solution:* Using the inequality  $n < 2^n$ , we quickly can conclude that  $n$  is  $O(2^n)$  by taking  $k = C = 1$  as witnesses. Note that because the logarithm function is increasing, taking logarithms (base 2) of both sides of this inequality shows that

$$\log n < n.$$

It follows that

$$\log n \text{ is } O(n).$$

(Again we take  $C = k = 1$  as witnesses.)

If we have logarithms to a base  $b$ , where  $b$  is different from 2, we still have  $\log_b n$  is  $O(n)$  because

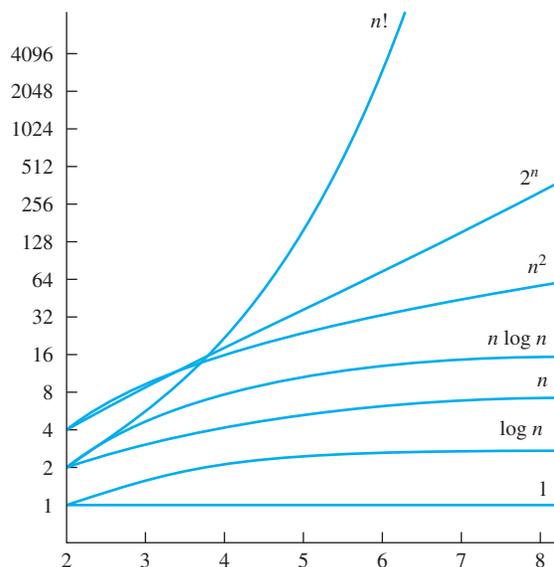
$$\log_b n = \frac{\log n}{\log b} < \frac{n}{\log b}$$

whenever  $n$  is a positive integer. We take  $C = 1/\log b$  and  $k = 1$  as witnesses. (We have used Theorem 3 in Appendix 2 to see that  $\log_b n = \log n / \log b$ .)

As mentioned before, big- $O$  notation is used to estimate the number of operations needed to solve a problem using a specified procedure or algorithm. The functions used in these estimates often include the following:

$$1, \log n, n, n \log n, n^2, 2^n, n!$$

Using calculus it can be shown that each function in the list is smaller than the succeeding function, in the sense that the ratio of a function and the succeeding function tends to zero as  $n$  grows without bound. Figure 3 displays the graphs of these functions, using a scale for the values of the functions that doubles for each successive marking on the graph. That is, the vertical scale in this graph is logarithmic.



**FIGURE 3** A Display of the Growth of Functions Commonly Used in Big- $O$  Estimates.