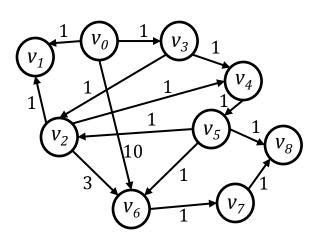
# Efficient Processing of Large Graphs via Input Reduction

Amlan Kusum, **Keval Vora**, Rajiv Gupta, Iulian Neamtiu

> HPDC'16 – Kyoto, Japan 04 June, 2016



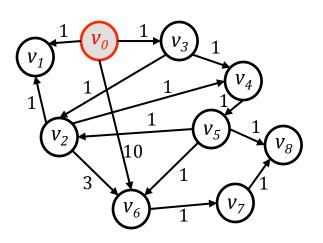
- Iterative graph algorithms
  - Vertices are processed over continuously
  - Highly parallel execution



	ı								
t	$v_o$	<i>v</i> <sub>1</sub>	$v_2$	$v_3$	$v_4$	$v_5$	$v_6$	$v_7$	$v_8$
$t_0$	0	$\infty$	$\infty$	∞	∞	∞	∞	∞	$\infty$
$t_1$	0	1	$\infty$	1	$\infty$	$\infty$	10	$\infty$	$\infty$
$t_2$	0	1	2	1	2	$\infty$	10	11	$\infty$
$t_3$	0	1	2	1	2	3	5	11	12
$t_4$	0	1	2	1	2	3	4	6	4
$t_5$	0	1	2	1	2	3	4	5	4
$t_6$	0	1	2	1	2	3	4	5	4



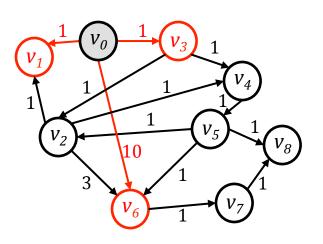
- Iterative graph algorithms
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t.	$ v_0 $	V.	V <sub>o</sub>	V <sub>o</sub>	ν.	v.	V.	v-	V <sub>o</sub>
$t_0$	0	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$
$t_1$	0	1	$\infty$	1	$\infty$	$\infty$	10	$\infty$	$\infty$
<i>t</i> _	0	1	2	1/	2	00	10	11	00
$t_3$	0	1	2	1	2	3	5	11	12
$t_4$	U	1	Z	1	Z	3	4	6	4
$t_5$	0	1	2	1	2	3	4	5	4
$t_6$	0	1	2	1	2	3	4	5	4



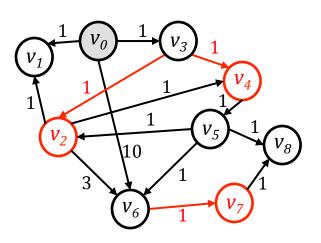
- Iterative graph algorithms
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_									
t	$v_{o}$	<i>v</i> <sub>1</sub>	$v_2$	$v_3$	$v_4$	<i>v</i> <sub>5</sub>	<i>v</i> <sub>6</sub>	$v_7$	<i>v</i> <sub>8</sub>
$t_0$	0	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$
$t_1$	0	1	$\infty$	1	$\infty$	$\infty$	10	$\infty$	$\infty$
$t_2$	0	1	2	1	2	$\infty$	10	11	$\infty$
$t_3$	0	1	2	1	2	3	5	11	12
$t_4$	0	1	2	1	2	3	4	6	4
$t_5$	0	1	2	1	2	3	4	5	4
$t_6$	0	1	2	1	2	3	4	5	4



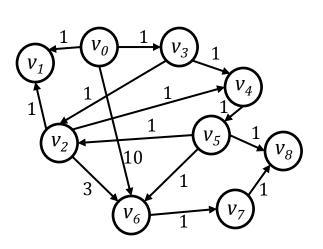
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	I								
	$v_{o}$								
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	0								
$t_2$	0	1	2	1	2	$\infty$	10	11	$\infty$
	0								
	0								
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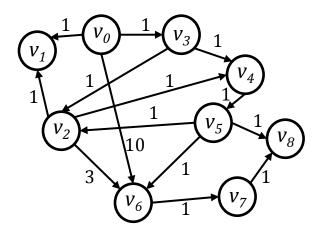
- Iterative graph algorithms
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  - Highly parallel execution
- > Challenging due to ever-growing graph sizes



	ı								
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$t_0$	0	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$
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$t_2$	0	1	2	1	2	$\infty$	10	11	$\infty$
•	0	1	2	1	2	3	5	11	12
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$t_5$	0	1	2	1	2			5	
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- Iterative graph algorithms
  - Vertices are processed over continuously
  - Highly parallel execution
- Challenging due to ever-growing graph sizes
- Convergence speed is dependent on initializations



t	$v_0$	$v_1$	$v_2$	$v_3$	$v_4$	$v_5$	$v_6$	$v_7$	$v_8$
$t_0$	0	$\infty$	∞	∞	$\infty$	3 3 3 3	$\infty$	∞	$\infty$
$t_1$	0	1	$\infty$	1	$\infty$	3	4	$\infty$	4
$t_2$	0	1	2	1	2	3	4	6	4
$t_3$	0	1	2	1	2	3	4	6	4



How to find better initializations?

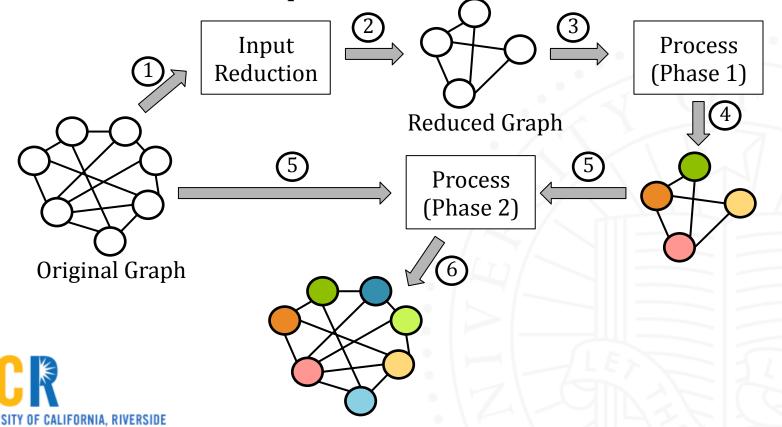
# **Key Idea**

- Compute initial values using a smaller signature of the original graph
  - Generate smaller graph using light-weight input reduction techniques



### **Key Idea**

- Compute initial values using a smaller signature of the original graph
  - Generate smaller graph using light-weight input reduction techniques

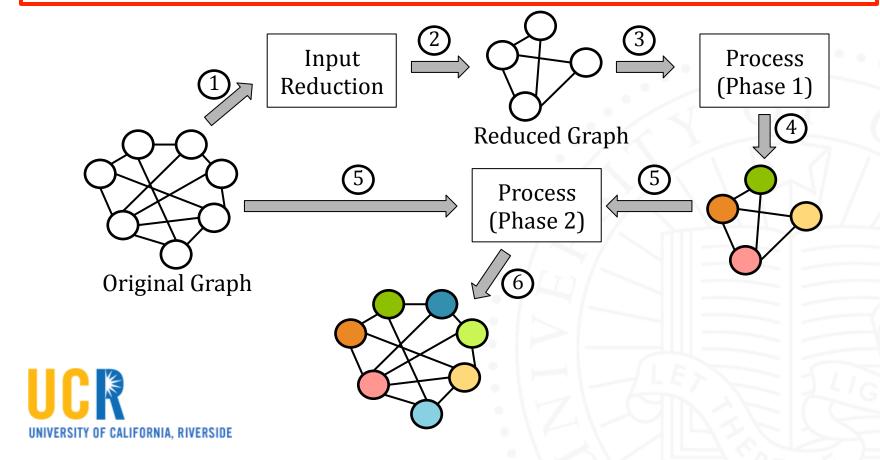


### **Key Idea**

#### time(Input Reduction)

- + time(Phase 1)
- + time(Phase 2)

< time(Original)



### **Outline**

- Input Reduction
- Vertex Transformations
- Correctness of Results
- > Evaluation
- Conclusion

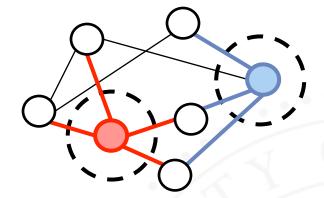


- Must be light-weight & general
  - Multilevel graph partitioning [SC'95, SC'01]
    - Matching based contraction [ICPP'95, JPDC'98]
    - Pruning based on edge costs affecting paths [ICDM'10]
  - Gate graph for shortest paths problem [ICDM'11]

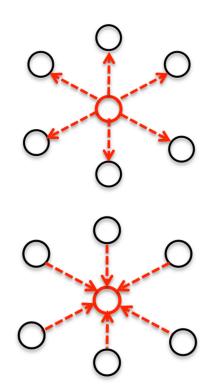
- Develop vertex level transformations
  - Easily parallelizable using the vertex centric graph processing systems

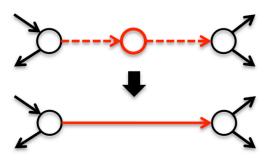


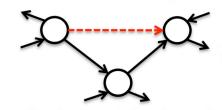
- Maintain structural integrity of the graph
  - Preserve the overall connectivity
- > Light-weight
  - Local
  - Non-interfering



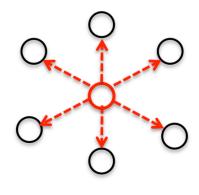


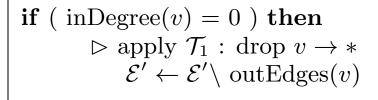








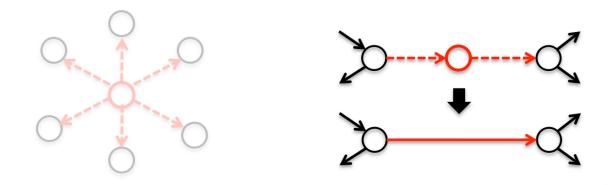






if (outDegree(
$$v$$
) = 0) then  
 $\Rightarrow$  apply  $\mathcal{T}_2$ : drop  $* \rightarrow v$   
 $\mathcal{E}' \leftarrow \mathcal{E}' \setminus \text{inEdges}(v)$ 



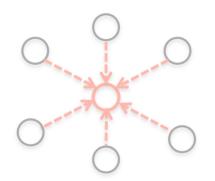


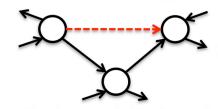
if (inDegree(v) = outDegree(v) = 1) then  

$$\Rightarrow$$
 apply  $\mathcal{T}_3$ : bypass v  
 $\mathcal{E}' \leftarrow (\mathcal{E}' \setminus \{u \rightarrow v, v \rightarrow w\}) \cup \{u \rightarrow w\}$   
where  $\{u \rightarrow v, v \rightarrow w\} \subseteq \mathcal{E}'$ 

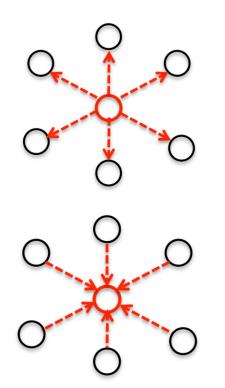


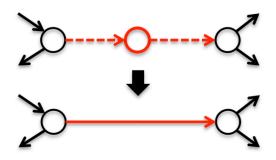
if  $(w \in \text{outNeighbors}(v) \text{ s.t. } w \text{ is unchanged and } \text{outNeighbors}(v) \cap \text{inNeighbors}(w) \neq \phi)$  then  $\triangleright \text{apply } \mathcal{T}_5 : \text{drop } v \rightarrow w$   $\mathcal{E}' \leftarrow \mathcal{E}' \setminus \{(v \rightarrow w)\}$ 

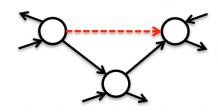












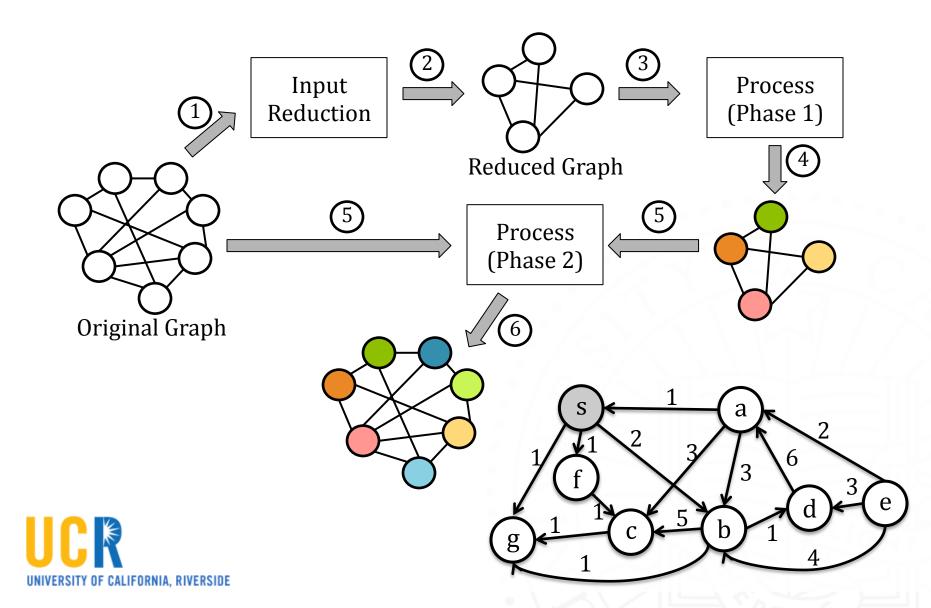


#### **Other Details**

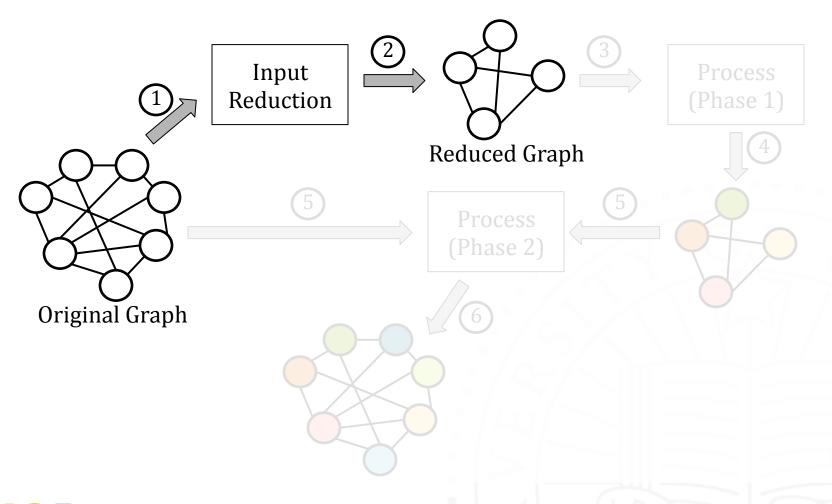
- More vertex transformations
  - Some relax structural integrity
- Order of transformations
  - Unified graph reduction algorithm



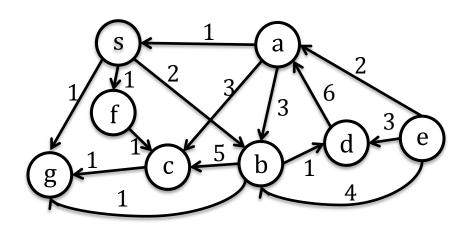
# **Processing workflow**

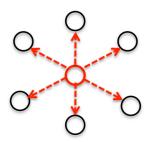


# **Processing workflow**

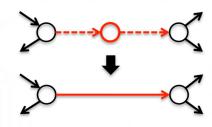


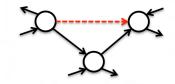




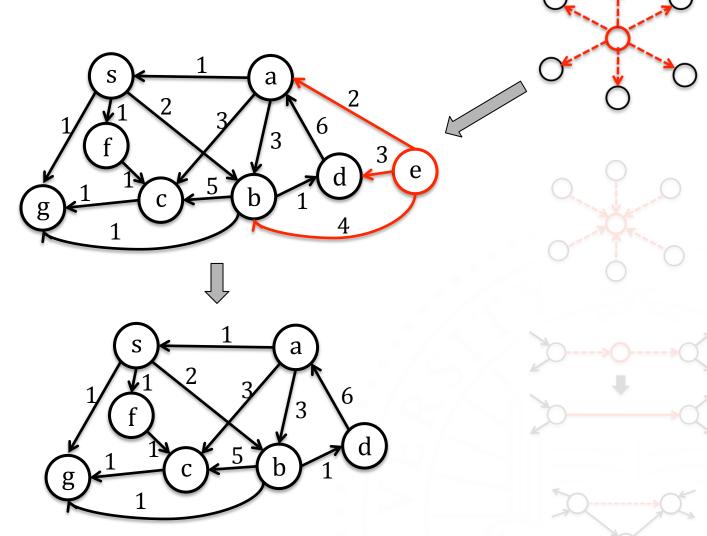




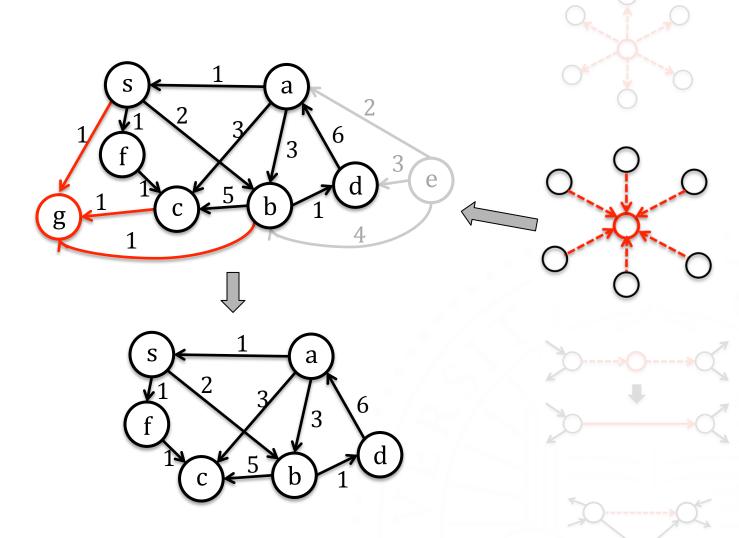




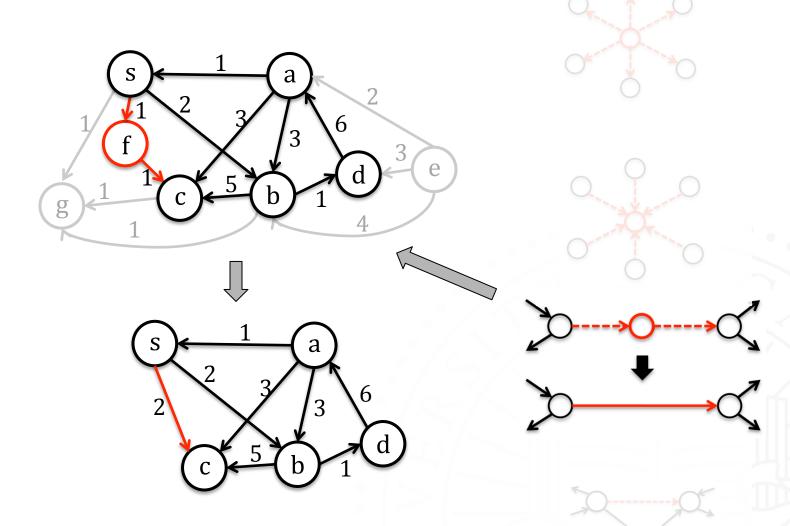




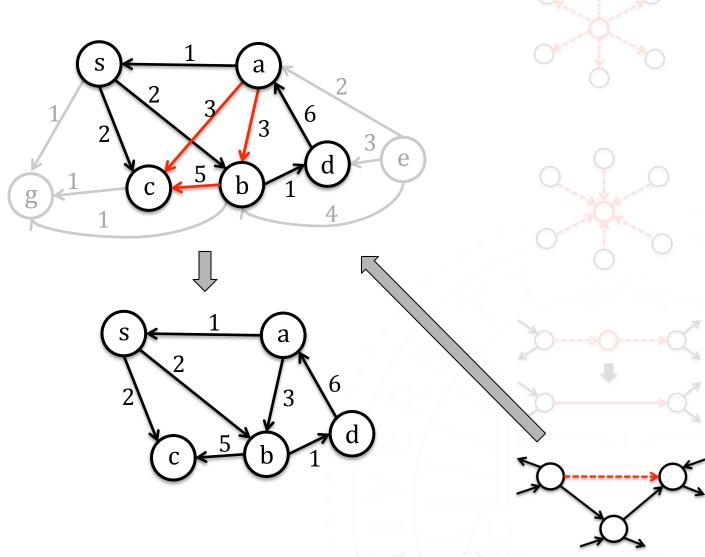




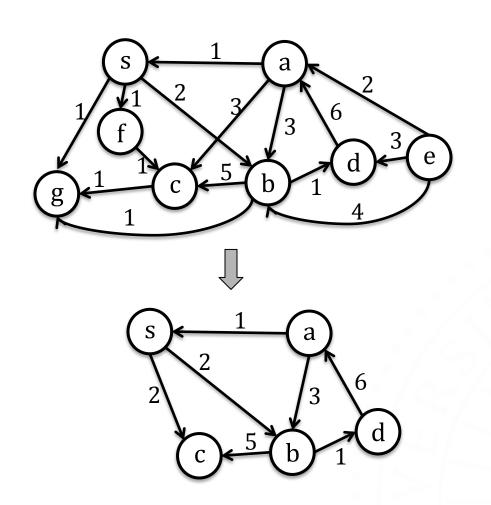


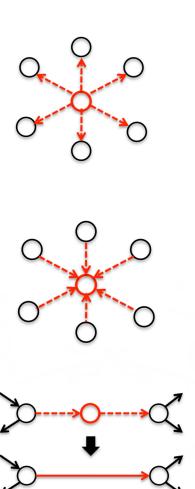






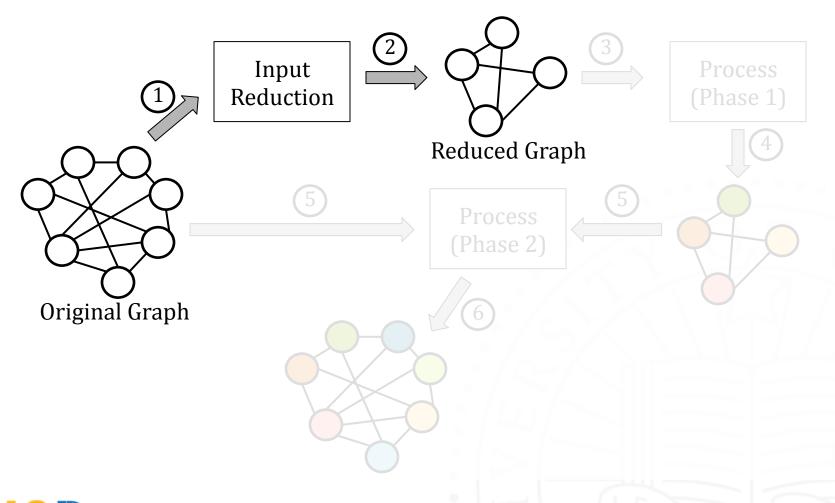






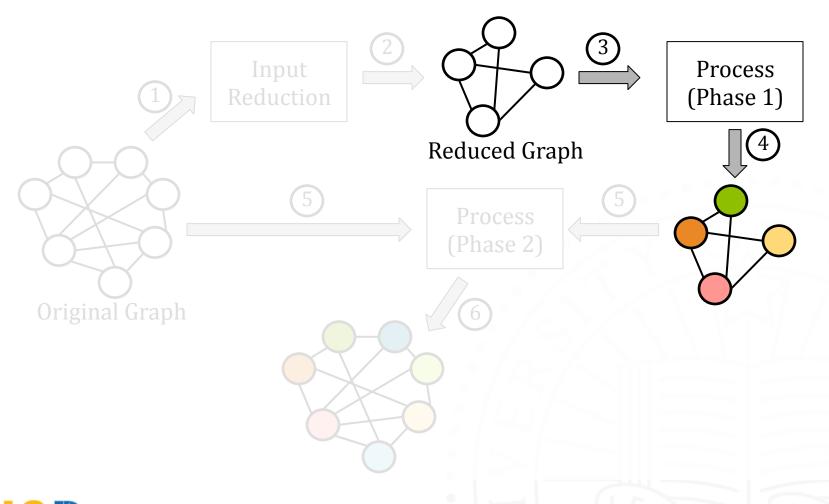


### Workflow





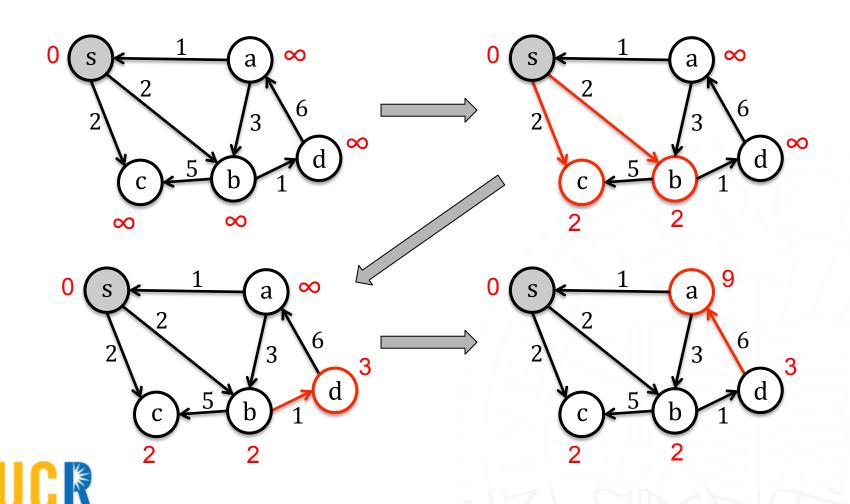
### Workflow



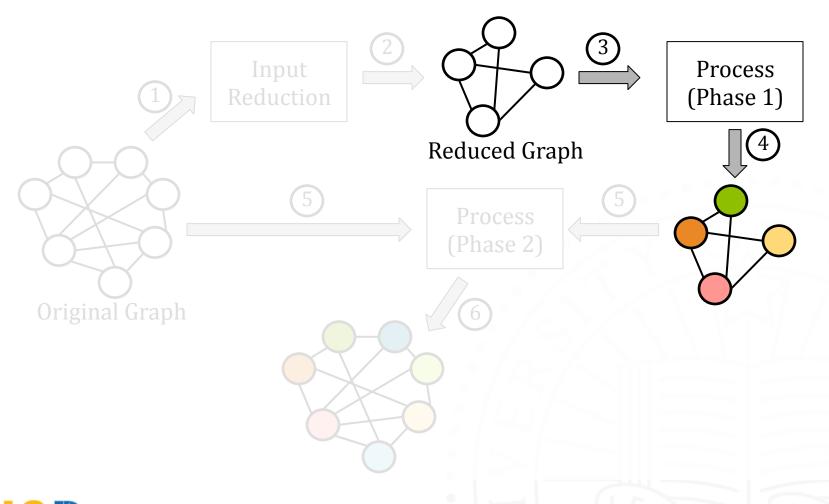


# **Processing Reduced Graph**

> Use the original iterative algorithm

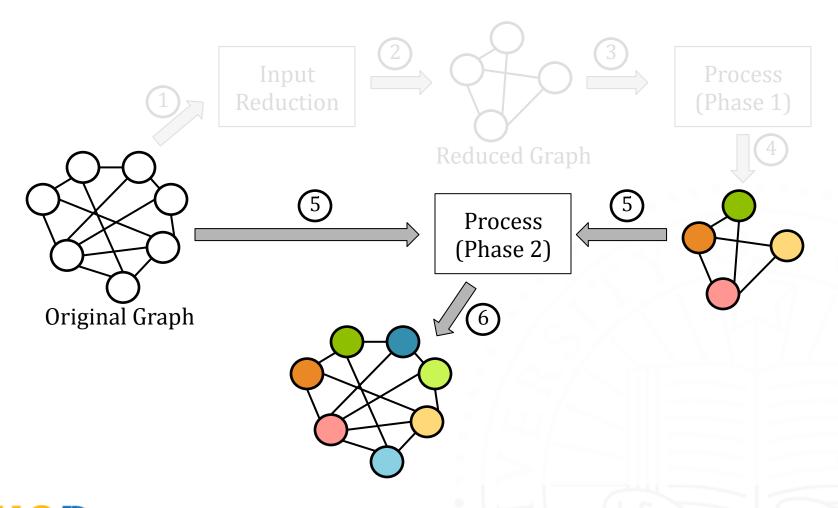


### Workflow





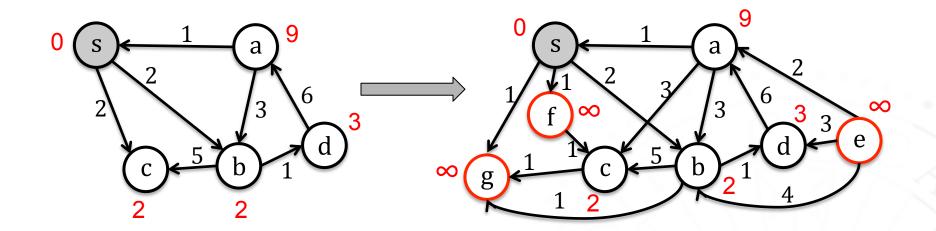
### Workflow





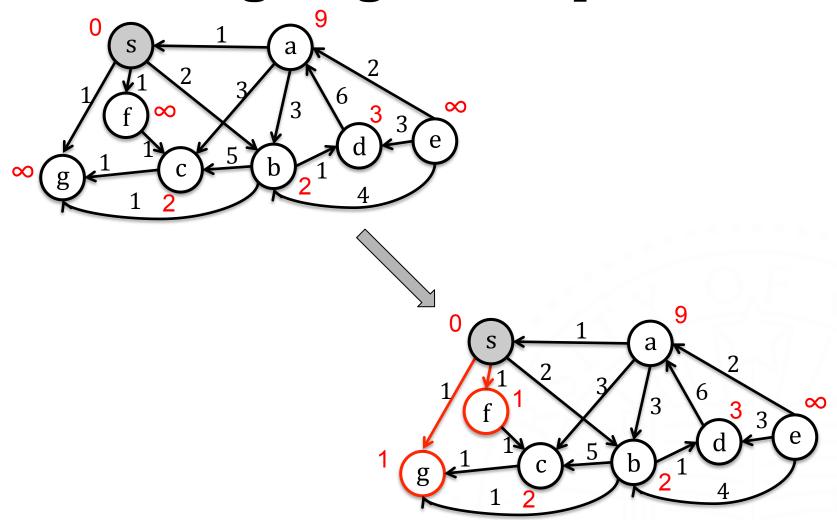
# **Mapping Results**

Use default values for missing vertices



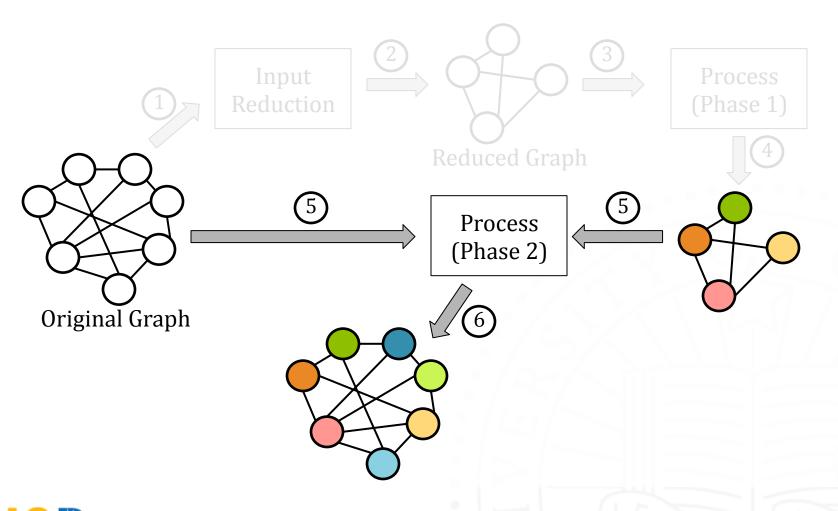


# **Processing Original Graph**



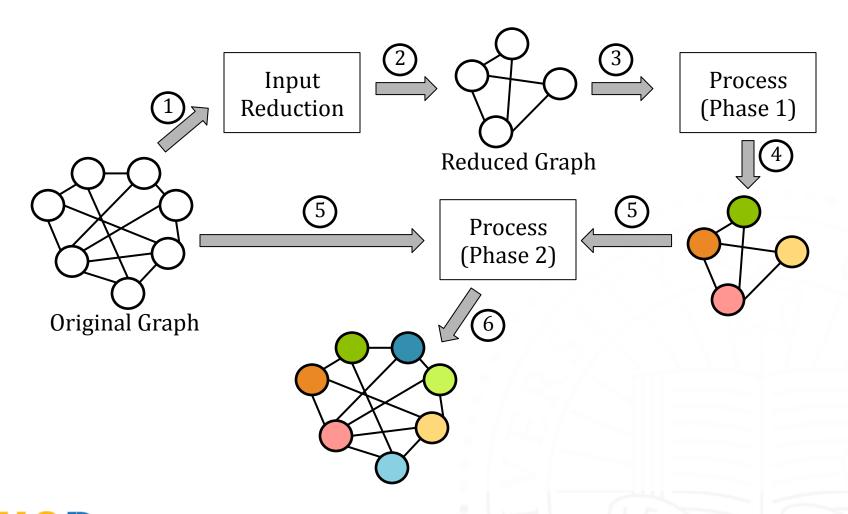


### Workflow

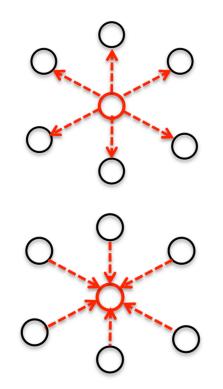


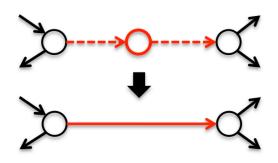


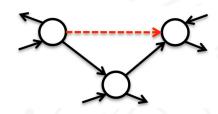
### Workflow



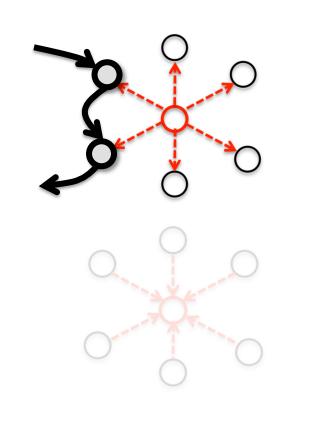






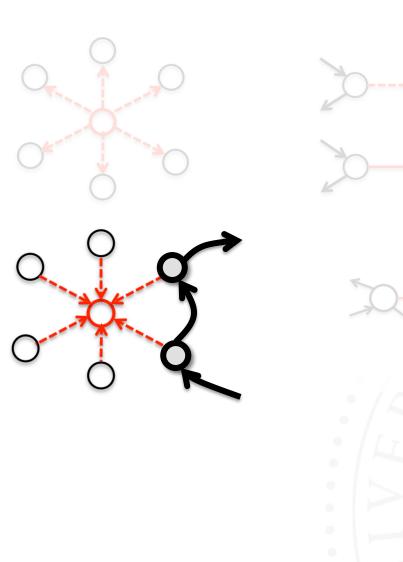






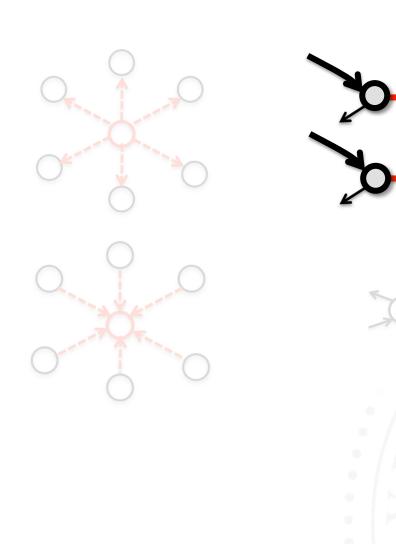




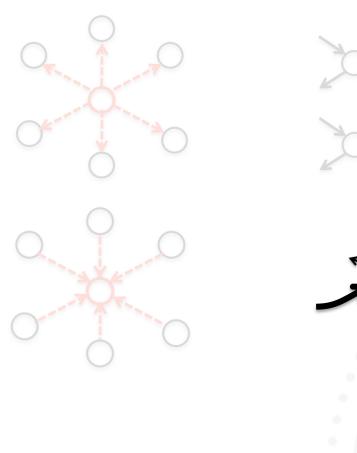


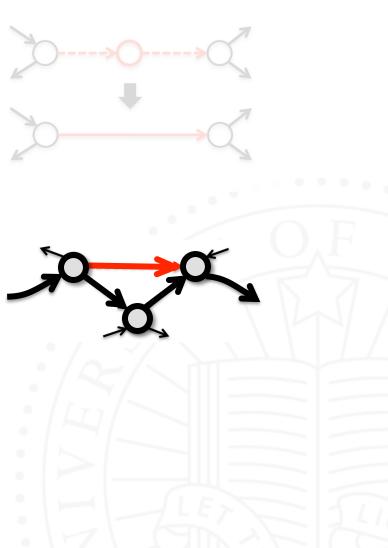














#### **Correctness**

- Transformation properties
  - Level of vertices, edges and components
  - Allow developing & reasoning for new transformations
- Algorithm behavior can be reasoned
  - Phase 2 initializations
  - Properties of aggregation function
- Correctness argued for algorithms used
  - 5 accurate and 1 approximate

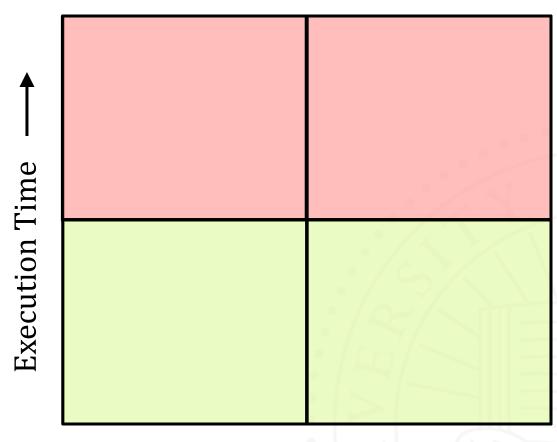


#### **Evaluation**

- Techniques independent of frameworks & processing environment
  - Incorporated in Galois [PLDI'11]
  - Single machine: 24-core, 32GB RAM
- > 6 benchmarks
  - PR, SSSP, SSWP, CC, GC, CD
- 4 input graphs
  - Friendster (|E| = 2.6B), Twitter (|E| = 1.5B),
     UKDomain (|E| = 936M), RMAT-24 (|E| = 268M)



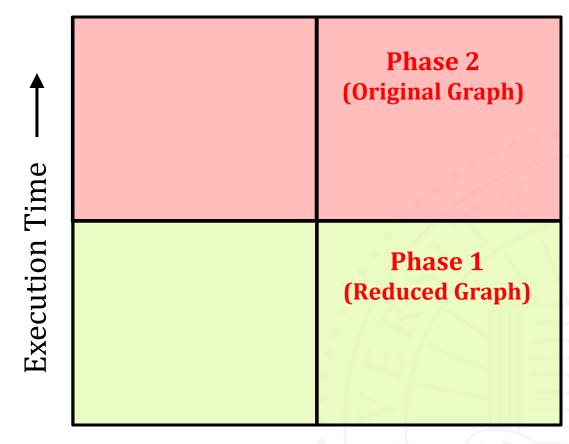
 $ERP = \frac{|E_{REDUCED}|}{|E_{ORIGINAL}|} \times 100$ 





Reduction --->

 $ERP = \frac{|E_{REDUCED}|}{|E_{ORIGINAL}|} \times 100$ 





Reduction --->

$$ERP = \frac{|E_{REDUCED}|}{|E_{ORIGINAL}|} \times 100$$

Execution Time —

Phase 1 (Reduced Graph)

Phase 2 (Original Graph)

Phase 2 (Original Graph)

Phase 1 (Reduced Graph)



Reduction —

$$> ERP = \frac{|E_{REDUCED}|}{|E_{ORIGINAL}|} \times 100$$

Execution Time —

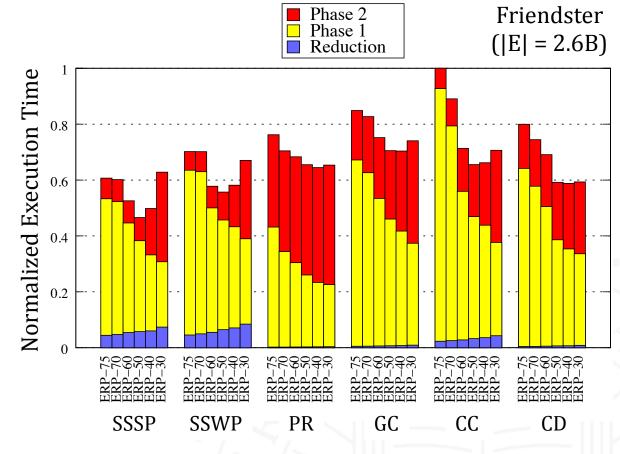
Phase 1	Phase 2
(Reduced Graph)	(Original Graph)
	Reduction
Phase 2	Phase 1
(Original Graph)	(Reduced Graph)
Reduction	



Reduction ---

#### **Execution Time**

- Speedups over parallel versions
- Speedups
   increase as ERP
   decreases up to
   an extent
- 1.3x 1.7x for75% 50%
- Structural dissimilarity for very low ERP

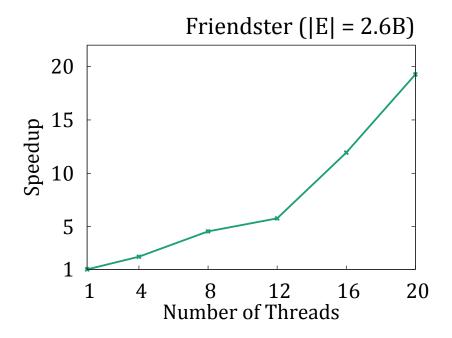


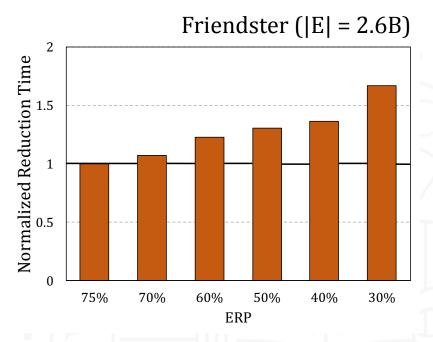
$$ERP = \frac{|E_{REDUCED}|}{|E_{ORIGINAL}|} \times 100$$



### **Input Reduction**

- Transformations are local, i.e., parallelizable
- > Higher reduction requires more work

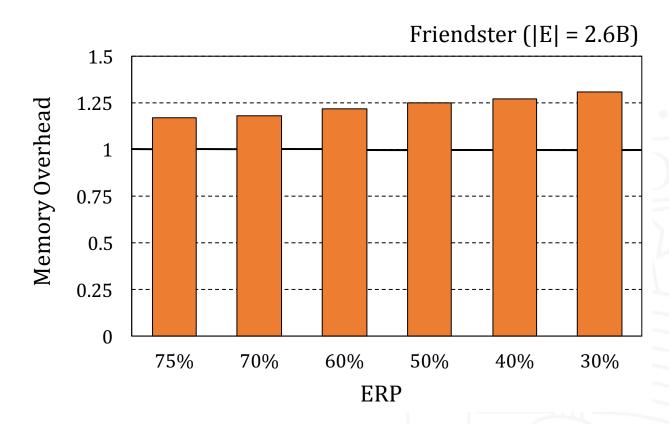






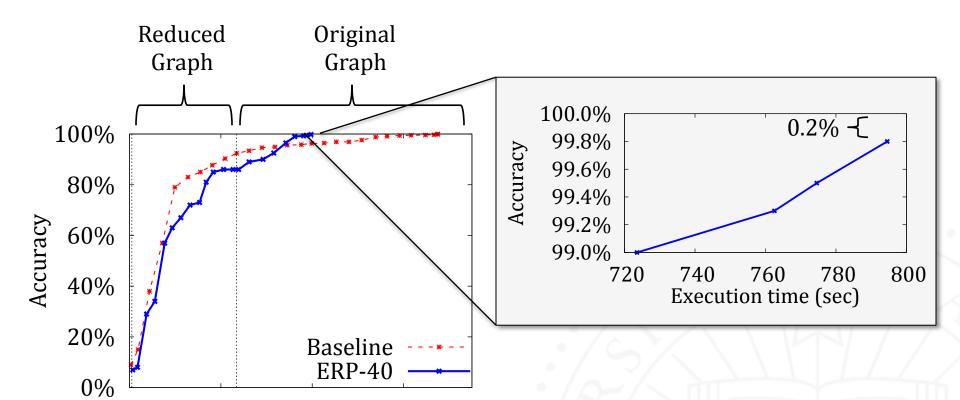
## **Memory Overhead**

- Tracking dissimilar elements
  - Newly added vertices & edges





## **Community Detection**



Friendster (|E| = 2.6B)

Execution time (sec)

800

1200



400

#### **More Results**

- Contribution of individual transformations
  - Some transformations more useful than others
  - Different graphs benefit from different transformations
- Improvement in scalability
- Results for all inputs



#### Conclusion

- Input reduction using transformations that are
  - Light-weight
  - Parallelizable
  - General
- Correctness reasoned using fine-grained transformation properties
- Achieve 1.25-2.14x speedups



### **Thanks**

- GRASP
  - http://grasp.cs.ucr.edu/

